

Postprocessing of quantum instruments

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Studying sequential measurements is of the utmost importance to both the foundational aspects of quantum theory and the practical implementations of quantum technologies, with both of these applications being abstractly described by the concatenation of quantum instruments into a sequence of a certain length. In general, the choice of instrument at any given step in the sequence can be conditionally chosen based on the classical results of all preceding instruments. For two instruments in a sequence we consider the conditional second instrument as an effective way of postprocessing the first instrument into a new one. This is similar to how a measurement described by a positive operator-valued measure (POVM) can be postprocessed into another by way of classical randomization of its outcomes using a stochastic matrix. In this work we study the postprocessing relation of instruments and the partial order it induces on their equivalence classes. We characterize the greatest and the least element of this order, give examples of postprocessings between different types of instruments, and draw connections between postprocessings of some of these instruments and their induced POVMs.

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I. INTRODUCTION

Measurements are the most essential part of every physical theory as they are, by definition, the only way to extract information from the objects that the theory describes. Detailed understanding and characterization of measurements may reveal crucial features of the theory. This is especially true in quantum theory where the discoveries of phenomena such as inevitable trade-off between information and disturbance or intrinsic randomness of measurement outcomes has made our (classical) intuition imprecise and renders quantum theory puzzling to most physicist even after a hundred years.

With the bloom of the field of quantum computation and information processing measurements are viewed not only as a mere tool for making predictions and verification of the theory but are also used as the key element assuring the implementation of the given task. For example, in quantum teleportation [1] they enable transmission of quantum information, but they can act as a step in quantum computation as well [2], or as the decoder of classical information in its transmission via noisy quantum channels [3,4].

Most often quantum measurements are considered in two ways. If we are interested only in the classical outcome produced by the quantum measurement then it is conveniently described by a positive operator-valued measure (POVM) [5]. However, if the measured quantum system will be part of any

further interaction then we shall use a quantum instrument [5] to describe probabilities of outcomes and the state change induced by the measurement. Especially if we want to consider evolutions, where measurements of the same system happen at different times, description of sequential measurements by quantum instruments becomes really practical.

Sequences of quantum measurements are at the same time very fundamental object to study and a fruitful tool for many applications. They can be used, for example, for quantum state tomography [6,7] and estimation, property testing [8], computation [9], quantum sequential decoding [10], joint measurability [11], or sequential state discrimination [12].

Let us illustrate the practical relevance of sequences of measurements with two of the previously mentioned applications. First, suppose we would be repeating the same instrument in the sequence. If the instrument corresponds to a least disturbing realization of a nondegenerate projective measurement (its Lüder's instrument) then such a repetition would not be useful at all, because the sequence of outcomes would be just a repetition of the first measurement outcome. However, in practical implementations, measurements are not ideal and will always be noisy. Recently, in Ref. [13], the authors discovered that repetition of such noisy measurement can lead to an effective measurement with suppressed noise level. Second, the authors of Ref. [14] showed that any jointly measurable pair of observables can be jointly measured by a Lüder's instrument of the first observable and a subsequent measurement of its output by a suitably chosen measurement. Hence, we can also say that joint measurability can be operationally realized via sequential measurements.

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In the current work we study a (concatenated) sequence of two quantum instruments, i.e., two quantum measurements having both classical and quantum output such that the classical outcome of the first instrument determines the second quantum instrument that will further process the quantum output of the first instrument. Thus, the concatenation of the first quantum instrument with the conditionally selected second instrument leads to a new compound quantum instrument. This can be viewed as a postprocessing of the first instrument into the resulting compound quantum instrument. If one fixes the first quantum instrument and considers all possible postprocessing quantum instruments, one can understand how an action of the first already-performed instrument can be altered or modified. In particular, one may also try to answer when it can be completely reversed.

The postprocessing relation can be used to define a partial order between (equivalence classes of) quantum instruments: if an (equivalence class of an) instrument can be postprocessed from some other (equivalence class of an) instrument, then the former must be greater than the latter. From the resource theory perspective this can be interpreted as a statement that the resulting instrument is less useful or less of a resource than the instrument that is being postprocessed. The special case of single-outcome quantum instruments corresponds to postprocessing of quantum channels, which was previously studied in Ref. [15–17].

The postprocessing that we define is analogous to the case of postprocessing of POVMs, where deterministic postprocessing of the classical outcomes of a POVM lead to a new POVM. In the case of quantum instruments, instead of just postprocessing the classical outcome we should also consider the postprocessing of the postmeasurement state, and in general these two aspects do not need to be independent. Thus, we are not restricted to changing the classical outcome by a classical postprocessing matrix and modifying the postmeasurement state by a quantum channel. In general, we can choose a different postprocessing quantum instrument for every outcome of the original instrument.

Study of the partial order induced by the postprocessing relation is one of the main goals of this paper. For general quantum instruments we characterize the least and the greatest element of the partial order and characterize their equivalence classes. We give examples of postprocessings of various types of instruments and draw connections between postprocessing of quantum instruments and postprocessing of their induced POVMs. Finally, we apply the postprocessing as a part of the simulation scheme for instruments, where in addition to postprocessing a given set of instruments we are also allowed to classically mix them. We note that quantum instruments that map to one-dimensional output space are isomorphic to POVMs. Thus, we must recover POVM simulability [18–21] as a special case of our results. In this sense our work can be also seen as a generalization of POVM simulability.

The rest of the paper is organized as follows: In Sec. II we introduce the relevant concepts and notation. In Sec. III we first consider the postprocessing of POVMs before generalizing this concept to instruments in Sec. IV. In this section we also analyze the structure given by the postprocessing relation. In Sec. V we show connections between the postprocessings of instruments and their induced POVMs. Finally, in Sec. VI

we focus on the concept of simulation of instruments using the postprocessing relation defined in Sec. IV. Our findings are summarized in Sec. VII.

II. PRELIMINARIES AND NOTATION

Let \mathcal{H} be a finite-dimensional complex Hilbert space. We denote by $\mathcal{L}(\mathcal{H})$ the set of bounded operators on \mathcal{H} and by $\mathcal{L}_s(\mathcal{H})$ the set of self-adjoint operators in $\mathcal{L}(\mathcal{H})$. States of a quantum system are described by positive semidefinite operators on \mathcal{H} with unit trace. We denote the set of quantum states on \mathcal{H} by $\mathcal{S}(\mathcal{H})$ so that

$$\mathcal{S}(\mathcal{H}) = \{\varrho \in \mathcal{L}_s(\mathcal{H}) | \varrho \geq 0, \text{tr}[\varrho] = 1\},$$

where 0 is the zero operator.

The set of effects on \mathcal{H} , denoted by $\mathcal{E}(\mathcal{H})$, consists of self-adjoint operators on \mathcal{H} bounded by 0 and I , where I (or $I_{\mathcal{H}}$ if we want to be more specific) is the identity operator on \mathcal{H} , i.e.,

$$\mathcal{E}(\mathcal{H}) = \{E \in \mathcal{L}_s(\mathcal{H}) | 0 \leq E \leq I\}.$$

An observable with a finite number of outcomes is described by a *positive operator-valued measure* (POVM), i.e., by mapping $\mathbf{A} : x \mapsto \mathbf{A}(x)$ from a finite outcome set Ω to $\mathcal{E}(\mathcal{H})$ such that $\sum_{x \in \Omega} \mathbf{A}(x) = I$. The set of observables on \mathcal{H} with outcome set Ω is denoted by $\mathcal{O}(\Omega, \mathcal{H})$.

Let \mathcal{H} and \mathcal{K} be Hilbert spaces. Transformations of states on \mathcal{H} to states on \mathcal{K} are described by *quantum channels*, i.e., completely positive trace-preserving maps from $\mathcal{L}(\mathcal{H})$ to $\mathcal{L}(\mathcal{K})$. Probabilistic transformations are described by *quantum operations*, i.e., completely positive trace-nonincreasing maps.

Quantum channels and operations have a well-known representation in an operator-sum form: a linear map $\mathcal{N} : \mathcal{L}(\mathcal{H}) \rightarrow \mathcal{L}(\mathcal{K})$ is a quantum operation if and only if there exists bounded operators $K_i : \mathcal{H} \rightarrow \mathcal{K}$ for all $i = 1, 2, \dots$ such that $\mathcal{N}(\varrho) = \sum_i K_i \varrho K_i^*$ for all $\varrho \in \mathcal{L}(\mathcal{H})$ and $\sum_i K_i^* K_i \leq I$. The operators K_i are called *Kraus operators* of \mathcal{N} , and in the finite-dimensional case it is possible to choose $\dim(\mathcal{H}) \dim(\mathcal{K})$ or fewer Kraus operators. The minimal number of Kraus operators for a given operation is called the *Kraus rank* of the operation.

A *quantum instrument* describes a device that takes a quantum input state and gives a quantum output state conditioned on a classical measurement outcome. Formally, a quantum instrument \mathcal{I} is a mapping $\mathcal{I} : x \mapsto \mathcal{I}_x$ from a finite outcome set Ω to the set of operations such that $\sum_{x \in \Omega} \mathcal{I}_x$ is a quantum channel. Given an input $\varrho \in \mathcal{S}(\mathcal{H})$, the (un-normalized) conditional output state is then described by $\mathcal{I}_x(\varrho)$ when we get the outcome x in the measurement of the *induced POVM* $\mathbf{A}^{\mathcal{I}} \in \mathcal{O}(\Omega, \mathcal{H})$ described by the probabilities $\text{tr}[\mathbf{A}^{\mathcal{I}}(x)\varrho] = \text{tr}[\mathcal{I}_x(\varrho)]$. In general it is clear that different instruments can have the same induced POVM but that, for any given instrument, the induced POVM is unique. The set of instruments from $\mathcal{L}(\mathcal{H})$ to $\mathcal{L}(\mathcal{K})$ with outcome set Ω is denoted by $\text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$. When the input and the output spaces are the same, $\mathcal{K} = \mathcal{H}$, we denote the set simply $\text{Ins}(\Omega, \mathcal{H})$.

Example 1. As an important class of instruments, we consider the *measure-and-prepare instruments* that perform a (demolishing) measurement on the input state and then

prepare a new state according to the measurement outcome. Thus, if $\mathbf{A} \in \mathcal{O}(\Omega, \mathcal{H})$ is a POVM, we can define a measure-and-prepare instrument $\mathcal{P}^{\mathbf{A}} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ for some set of states $\{\xi_x\}_{x \in \Omega} \subset \mathcal{S}(\mathcal{K})$ as $\mathcal{P}_x^{\mathbf{A}}(\varrho) = \text{tr}[\mathbf{A}(x)\varrho]\xi_x$ for all $\varrho \in \mathcal{L}(\mathcal{H})$. Note that, by changing the set $\{\xi_x\}_{x \in \Omega}$, one can use the same POVM \mathbf{A} to define countless measure-and-prepare instruments. As a special case of measure-and-prepare instruments we have the *trash-and-prepare instruments* where one just fixes the measured POVM to be trivial, i.e., $\mathbf{A}(x) = p_x I$ for all $x \in \Omega$ for some probability distribution $(p_x)_x$ over Ω so that $\mathcal{P}_x^{\mathbf{A}}(\varrho) = \text{tr}[\varrho]p_x\xi_x$ for all $\varrho \in \mathcal{L}(\mathcal{H})$ and $x \in \Omega$.

III. POSTPROCESSING OF POSITIVE OPERATOR-VALUED MEASURES

Before generalizing the concept of postprocessing to instruments, we recall some important results for POVMs. After obtaining the outcome statistics of a measurement of an observable, one may want to process the information obtained. One can, for instance, see if it is possible to reveal some other property of the system by manipulating the data and obtain the outcome statistics of some other observable. This is what is usually called the postprocessing of observables.

A. The postprocessing partial order

We can formalize the previous paragraph with the following definition:

Definition 1. Let $\mathbf{A} \in \mathcal{O}(\Omega_A, \mathcal{H})$ and $\mathbf{B} \in \mathcal{O}(\Omega_B, \mathcal{H})$ be observables. If there exists a stochastic matrix $\nu = (\nu_{xy})_{x \in \Omega_A, y \in \Omega_B}$, i.e., $\nu_{xy} \geq 0$ for all $x \in \Omega_A, y \in \Omega_B$, and $\sum_{y \in \Omega_B} \nu_{xy} = 1$ for all $x \in \Omega_A$, such that

$$\mathbf{B}(y) = \sum_{x \in \Omega_A} \nu_{xy} \mathbf{A}(x)$$

for all $y \in \Omega_B$, we say that \mathbf{B} is a *postprocessing* of \mathbf{A} and denote it $\mathbf{A} \rightarrow \mathbf{B}$. Furthermore, we say that observables \mathbf{A} and \mathbf{B} are *postprocessing equivalent*, denoted by $\mathbf{A} \leftrightarrow \mathbf{B}$, if $\mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{B} \rightarrow \mathbf{A}$.

Example 2. A special kind of postprocessing, called *re-labeling*, is one where all the elements of the stochastic postprocessing matrix are either 0 or 1. Following Ref. [22], this can be formalized by the existence of a function $f: \Omega_A \rightarrow \Omega_B$ such that $\nu_{xy} = \delta_{f(x), y}$, where $\delta_{x, x'}$ is the Kronecker delta, so that

$$\mathbf{B}(y) = \sum_{x \in f^{-1}(y)} \mathbf{A}(x)$$

for all $y \in \Omega_B$. In this case, we say that \mathbf{B} is a *relabeling* of \mathbf{A} and that \mathbf{A} is a *refinement* of \mathbf{B} .

Postprocessing captures the idea that the outcome statistics of \mathbf{B} can be deterministically obtained from the statistics of \mathbf{A} by some classical process represented by the stochastic matrix. It is easy to see that postprocessing induces a pre-order on the set of all observables on \mathcal{H} , and by extending it to the equivalence classes of postprocessing equivalent observables it becomes a partial order. A natural thing to consider is whether there exists a least or greatest element with respect to this order.

It is easy to see that the trivial observables, i.e., observables $\mathbf{T}^p \in \mathcal{O}(\Omega_{\mathbf{T}^p}, \mathcal{H})$ of the form $\mathbf{T}^p(x) = p_x I$ for all $x \in \Omega_{\mathbf{T}^p}$ for some probability distribution $(p_x)_x$ on $\Omega_{\mathbf{T}^p}$, can be postprocessed from any other observable $\mathbf{A} \in \mathcal{O}(\Omega_A, \mathcal{H})$ by using the postprocessing matrix ν with $\nu_{yx} = p_x$ for all $x \in \Omega_{\mathbf{T}^p}$ and $y \in \Omega_A$. Thus, $\mathbf{A} \rightarrow \mathbf{T}^p$ for any observable \mathbf{A} and any probability distribution p . Furthermore, if \mathbf{B} is an observable such that $\mathbf{T}^p \rightarrow \mathbf{B}$ with a postprocessing μ , then also \mathbf{B} is a trivial observable, $\mathbf{B} = \mathbf{T}^q$, where $q_z = \sum_{x \in \Omega_{\mathbf{T}^p}} \mu_{xz} p_x$ for all $z \in \Omega_B$. Thus, the equivalence class of trivial observables is the least element with respect to the partial order.

How about the greatest element? It turns out that there is no greatest element [15,23]. Instead, we get a class of maximal elements that we call postprocessing clean observables.

Definition 2. An observable \mathbf{A} is *postprocessing clean* if for any observable \mathbf{B} such that $\mathbf{B} \rightarrow \mathbf{A}$ we also have $\mathbf{A} \rightarrow \mathbf{B}$.

The postprocessing clean observables were characterized in Refs. [23,24]: an observable is postprocessing clean if and only if it is rank 1, i.e., each of its effects is a rank-1 operator. For a POVM $\mathbf{A} \in \mathcal{O}(\Omega_A, \mathcal{H})$, being rank 1 is equivalent to being *indecomposable* [25], i.e., if any of its nonzero effect $\mathbf{A}(x)$ is decomposed as a sum of some two effects on \mathcal{H} so that $\mathbf{A}(x) = E_x + F_x$ for some $E_x, F_x \in \mathcal{E}(\mathcal{H})$, then there exist positive numbers $e_x, f_x > 0$ such that $\mathbf{A}(x) = e_x E_x = f_x F_x$. It holds that any observable can be postprocessed from a postprocessing clean observable. Thus, rank-1 POVMs are in fact the maximal elements with respect to the postprocessing partial order as everything else can be postprocessed from them.

B. Minimally sufficient positive operator-valued measures

Postprocessing can thus be seen as a way to construct new observables out of existing ones by a classical process. Another way to look at postprocessing is to say that if we have $\mathbf{A} \rightarrow \mathbf{B}$ for two observables \mathbf{A} and \mathbf{B} , then \mathbf{A} must be more informative because \mathbf{B} can be deduced from \mathbf{A} . But as was pointed out earlier, when we talk about the partial order induced by the postprocessing relation, we are actually comparing equivalence classes of observables. Then especially postprocessing equivalent observables would be just as informative. However, even though they can be seen as having the same information, the following notion introduced in Ref. [26] captures the idea that, even in the same equivalence class, there are observables with minimum informational redundancy:

Definition 3. An observable \mathbf{A} is *minimally sufficient* if, whenever $\mathbf{A} \leftrightarrow \mathbf{B}$ with some observable \mathbf{B} , then \mathbf{B} is a refinement of \mathbf{A} .

It was shown in Ref. [26] that a (discrete) POVM $\mathbf{A} \in \mathcal{O}(\Omega_A, \mathcal{H})$ is minimally sufficient if and only if it is *non-vanishing*, i.e., $\mathbf{A}(x) \neq 0$ for all $x \in \Omega_A$, and it is *pairwise linearly independent*, i.e., $\mathbf{A}(x) \neq c\mathbf{A}(y)$ for any $c > 0$ for all $x \neq y, x, y \in \Omega_A$. Furthermore, for any POVM \mathbf{A} , there exists a minimally sufficient POVM $\tilde{\mathbf{A}}$ such that $\mathbf{A} \leftrightarrow \tilde{\mathbf{A}}$, and $\tilde{\mathbf{A}}$ is unique up to a bijective relabeling of its outcomes.

The minimally sufficient representative of the equivalence class of a POVM \mathbf{A} can be constructed as follows: define an equivalence relation \sim in Ω_A so that $x \sim y$ if and only if there exists $c > 0$ such that $\mathbf{A}(x) = c\mathbf{A}(y)$. We denote the set of

equivalence classes Ω_A/\sim by $\tilde{\Omega}_A$ and define a minimally sufficient POVM $\tilde{\mathbf{A}} \in \mathcal{O}(\tilde{\Omega}_A, \mathcal{H})$ that is postprocessing equivalent to \mathbf{A} by

$$\tilde{\mathbf{A}}([y]) = \sum_{x \in [y]} \mathbf{A}(x), \quad [y] \in \tilde{\Omega}_A. \quad (1)$$

The uniqueness of the pairwise linearly independent minimally sufficient representative can be used to characterize the whole postprocessing equivalence class: two POVMs \mathbf{A} and \mathbf{B} are postprocessing equivalent if and only if the pairwise linearly independent POVMs $\tilde{\mathbf{A}}$ and $\tilde{\mathbf{B}}$ are bijective relabelings of each other. This shows that the effects of two postprocessing equivalent POVMs must be proportional to each other. This result can be extracted from Refs. [23,26] but we still provide a proof for completeness because we will use it later.

Proposition 1. Let $\mathbf{A} \in \mathcal{O}(\Omega_A, \mathcal{H})$ and $\mathbf{B} \in \mathcal{O}(\Omega_B, \mathcal{H})$ be two postprocessing equivalent nonvanishing POVMs. Then for all $x \in \Omega_A$ there exists $y_x \in \Omega_B$ and $c_{xy_x} > 0$ such that $\mathbf{A}(x) = c_{xy_x} \mathbf{B}(y_x)$. Furthermore, there exist postprocessings ν for $\mathbf{B} \rightarrow \mathbf{A}$ and μ for $\mathbf{A} \rightarrow \mathbf{B}$ such that $\nu_{yx}, \mu_{xy} \neq 0$ only if $\mathbf{B}(y)$ is proportional to $\mathbf{A}(x)$.

Proof. We define the pairwise linearly independent POVMs $\tilde{\mathbf{A}} \in \mathcal{O}(\tilde{\Omega}_A, \mathcal{H})$ and $\tilde{\mathbf{B}} \in \mathcal{O}(\tilde{\Omega}_B, \mathcal{H})$ as in Eq. (1), so that

$$\tilde{\mathbf{A}}([x]) = \sum_{x' \in [x]} \mathbf{A}(x'), \quad \tilde{\mathbf{B}}([y]) = \sum_{y' \in [y]} \mathbf{B}(y'),$$

for all $[x] \in \tilde{\Omega}_A$ and $[y] \in \tilde{\Omega}_B$. Thus, for all $x' \in [x]$, we have that $\mathbf{A}(x') = c_{x'} \tilde{\mathbf{A}}([x])$ for some $c_{x'} \in (0, 1]$ such that $\sum_{x' \in [x]} c_{x'} = 1$. Similarly, for all $y' \in [y]$, we have that $\mathbf{B}(y') = d_{y'} \tilde{\mathbf{B}}([y])$ for some $d_{y'} \in (0, 1]$ such that $\sum_{y' \in [y]} d_{y'} = 1$.

As was mentioned earlier, the minimally sufficient representative is essentially unique in each equivalence class, so that since $\tilde{\mathbf{A}} \leftrightarrow \mathbf{A} \leftrightarrow \mathbf{B} \leftrightarrow \tilde{\mathbf{B}}$ and since $\tilde{\mathbf{A}}$ and $\tilde{\mathbf{B}}$ are both minimally sufficient, there exists a bijective map $f: \tilde{\Omega}_A \rightarrow \tilde{\Omega}_B$ such that $\tilde{\mathbf{A}}([x]) = \tilde{\mathbf{B}}(f([x]))$ for all $[x] \in \tilde{\Omega}_A$. Thus, for each $x \in \Omega_A$ and $y_x \in f([x]) \subset \tilde{\Omega}_B$ we have that

$$\mathbf{A}(x) = c_x \tilde{\mathbf{A}}([x]) = c_x \tilde{\mathbf{B}}(f([x])) = c_x \tilde{\mathbf{B}}([y_x]) = c_{xy_x} \mathbf{B}(y_x),$$

where we have denoted $c_{xy_x} = c_x/d_{y_x} > 0$.

For the second part of the claim, let us define $\nu_{yx} = c_x \delta_{f([x]), [y]}$ for all $x \in \Omega_A$ and $y \in \Omega_B$. Clearly for all $y \in \Omega_B$ we have that

$$\sum_{x \in \Omega_A} \nu_{yx} = \sum_{x \in f^{-1}([y])} c_x = 1$$

since $f^{-1}([y]) = [x']$ for some $[x'] \in \tilde{\Omega}_A$ for all $y \in \Omega_B$. Furthermore,

$$\sum_{y \in \Omega_B} \nu_{yx} \mathbf{B}(y) = \sum_{y \in f([x])} c_x \mathbf{B}(y) = c_x \tilde{\mathbf{B}}(f([x])) = \mathbf{A}(x)$$

for all $x \in \Omega_A$. Thus, ν is a postprocessing for $\mathbf{B} \rightarrow \mathbf{A}$ that has $\nu_{yx} \neq 0$ only if $\mathbf{B}(y)$ is proportional to $\mathbf{A}(x)$. The postprocessing μ for $\mathbf{A} \rightarrow \mathbf{B}$ can be defined analogously. ■

Example 3. We note that not all POVMs whose effects are proportional to each other are postprocessing equivalent. For

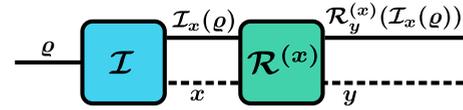


FIG. 1. The postprocessing of instrument \mathcal{I} by the instruments $\mathcal{R}^{(x)}$. Instrument \mathcal{I} takes quantum state ρ as an input and after measuring it and obtaining a measurement outcome x the instrument $\mathcal{R}^{(x)}$ is chosen. The conditional output state of \mathcal{I} [described by $\mathcal{I}_x(\rho)$] serves as an input to $\mathcal{R}^{(x)}$, which measures it and produces an outcome y after which the output is in the (un-normalized) state $\mathcal{R}_y^{(x)}(\mathcal{I}_x(\rho))$.

example, let us define two four-outcome qubit POVMs \mathbf{A} and \mathbf{B} as

$$\mathbf{A}(1) = \frac{1}{2}|\varphi_1\rangle\langle\varphi_1|, \quad \mathbf{A}(2) = \frac{1}{2}|\varphi_2\rangle\langle\varphi_2|, \quad (2)$$

$$\mathbf{A}(3) = \frac{1}{2}|\psi_1\rangle\langle\psi_1|, \quad \mathbf{A}(4) = \frac{1}{2}|\psi_2\rangle\langle\psi_2|, \quad (3)$$

$$\mathbf{B}(1) = \frac{1}{3}|\varphi_1\rangle\langle\varphi_1|, \quad \mathbf{B}(2) = \frac{1}{3}|\varphi_2\rangle\langle\varphi_2|, \quad (4)$$

$$\mathbf{B}(3) = \frac{2}{3}|\psi_1\rangle\langle\psi_1|, \quad \mathbf{B}(4) = \frac{2}{3}|\psi_2\rangle\langle\psi_2|, \quad (5)$$

where $\{\varphi_1, \varphi_2\}$ and $\{\psi_1, \psi_2\}$ are two orthonormal bases in \mathbb{C}^2 . We see that, although the effects of \mathbf{A} and \mathbf{B} are proportional, they cannot be postprocessed from each other since, for instance, $\mathbf{A}(1) = 3/2 \mathbf{B}(1)$ and $\mathbf{B}(3) = 4/3 \mathbf{A}(3)$, where obviously the postprocessing elements $3/2$ and $4/3$ would be larger than one. Instead, as required by the bijective relabeling of two minimally sufficient representatives, two POVMs whose effects are proportional to each other are postprocessing equivalent if their pairwise linearly dependent effects sum up to the same effect for both observables. For the POVMs defined in Eqs. (2)–(5) this would mean that the coefficients of the rank-1 projectors would have to be the same for \mathbf{A} and \mathbf{B} .

IV. POSTPROCESSING OF INSTRUMENTS

We can now define postprocessing of instruments analogously to the postprocessing of POVMs but unlike in the case of POVMs we are not only processing classical information but we must also process the output state of the instrument.

Definition 4. Let $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ and $\mathcal{J} \in \text{Ins}(\Lambda, \mathcal{H}, \mathcal{V})$ be quantum instruments. If there exists a set of instruments $\{\mathcal{R}^{(x)}\}_{x \in \Omega} \subset \text{Ins}(\Lambda, \mathcal{K}, \mathcal{V})$ such that

$$\mathcal{J}_y(\rho) = \sum_{x \in \Omega} \mathcal{R}_y^{(x)}(\mathcal{I}_x(\rho)) \quad (6)$$

for all $\rho \in \mathcal{L}(\mathcal{H})$ and $y \in \Lambda$, then we denote $\mathcal{I} \rightarrow \mathcal{J}$ and say that \mathcal{J} is a *postprocessing* of \mathcal{I} . Furthermore, we say that \mathcal{I} and \mathcal{J} are *postprocessing equivalent*, denoted by $\mathcal{I} \leftrightarrow \mathcal{J}$, if $\mathcal{I} \rightarrow \mathcal{J}$ and $\mathcal{J} \rightarrow \mathcal{I}$.

The postprocessing Eq. (6) is depicted in Fig. 1. As an illustrative example of postprocessing of instruments, we consider what type of instruments can be postprocessed from measure-and-prepare instruments.

Example 4. Let $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ be a measure-and-prepare instrument with an induced POVM \mathbf{A} , i.e., it is of the

form

$$\mathcal{I}_x(\varrho) = \text{tr}[\mathbf{A}(x)\varrho]\sigma_x$$

for all $\varrho \in \mathcal{S}(\mathcal{H})$ and $x \in \Omega$ for some set of states $\{\sigma_x\}_{x \in \Omega} \subset \mathcal{S}(\mathcal{K})$. Let $\mathcal{J} \in \text{Ins}(\Lambda, \mathcal{H}, \mathcal{V})$ be an instrument such that $\mathcal{I} \rightarrow \mathcal{J}$ so that there exist instruments $\mathcal{R}^{(x)} \in \text{Ins}(\Lambda, \mathcal{K}, \mathcal{V})$ such that $\mathcal{J}_y(\varrho) = \sum_x \mathcal{R}_y^{(x)}(\mathcal{I}_x(\varrho))$ for all $\varrho \in \mathcal{S}(\mathcal{H})$ and $y \in \Lambda$. By expanding the previous expression, we see that

$$\mathcal{J}_y(\varrho) = \sum_{x \in \Omega} \text{tr}[\mathbf{A}(x)\varrho]\mathcal{R}_y^{(x)}(\sigma_x)$$

for all $y \in \Lambda$. By denoting $v_{xy} = \text{tr}[\mathcal{R}_y^{(x)}(\sigma_x)] \in [0, 1]$ for all $x \in \Omega$ and $y \in \Lambda$, we can define $\xi_{xy} = \mathcal{R}_y^{(x)}(\sigma_x)/v_{xy} \in \mathcal{S}(\mathcal{V})$ when $v_{xy} \neq 0$ and $\xi_{xy} = \xi$ for some fixed $\xi \in \mathcal{S}(\mathcal{V})$ when $v_{xy} = 0$. We see that $\sum_y v_{xy} = 1$ for all $x \in \Omega$ so that v is actually a valid postprocessing. Thus, we have that

$$\mathcal{J}_y(\varrho) = \sum_{x \in \Omega} \text{tr}[v_{xy}\mathbf{A}(x)\varrho]\xi_{xy}$$

for all $y \in \Lambda$ and $\varrho \in \mathcal{S}(\mathcal{H})$. We note that the induced POVM $\mathbf{A}^{\mathcal{J}}$ of \mathcal{J} is then a postprocessing of \mathbf{A} since by taking the trace on the last expression we see that $\mathbf{A}^{\mathcal{J}}(y) = \sum_{x \in \Omega} v_{xy}\mathbf{A}(x)$ for all $y \in \Lambda$ so that $\mathbf{A} = \mathbf{A}^{\mathcal{I}} \rightarrow \mathbf{A}^{\mathcal{J}}$.

We note that, in the case of channels, Definition 4 reduces to the pre-order given in Ref. [15]: If $\mathcal{C} : \mathcal{L}(\mathcal{H}) \rightarrow \mathcal{L}(\mathcal{K})$ and $\mathcal{D} : \mathcal{L}(\mathcal{H}) \rightarrow \mathcal{L}(\mathcal{V})$ are two quantum channels such that there exists a channel $\mathcal{E} : \mathcal{L}(\mathcal{K}) \rightarrow \mathcal{L}(\mathcal{V})$ such that $\mathcal{D} = \mathcal{E} \circ \mathcal{C}$, where \circ denotes the composition of maps, then $\mathcal{C} \rightarrow \mathcal{D}$. Just as in the case of channels, the defined postprocessing gives a pre-order in the set of instruments so that it can be used to define a partial order for the equivalence classes of instruments.

A. Characterizing the greatest and the least element

As in the case of POVMs, one of the first questions is to determine the greatest and the least element (if they exist) with respect to the postprocessing partial order. Let us start with the greatest element, and as before let us first consider the maximal elements, namely, the postprocessing clean instruments.

Definition 5. An instrument \mathcal{I} is *postprocessing clean* if for any instrument \mathcal{J} such that $\mathcal{J} \rightarrow \mathcal{I}$ we also have $\mathcal{I} \rightarrow \mathcal{J}$.

Unlike in the case of POVMs, there is only one maximal element in the equivalence classes of instruments so that it must also be the greatest element. As expected, this is similar to the case of channels that was studied in Ref. [15] and the greatest element is the same.

Proposition 2. An instrument is postprocessing clean if and only if it is postprocessing equivalent with the identity channel id .

Proof. First of all, it is clear that one can postprocess every instrument with an input Hilbert space \mathcal{H} from the identity channel $id : \mathcal{L}(\mathcal{H}) \rightarrow \mathcal{L}(\mathcal{H})$ defined as $id(\varrho) = \varrho$ for all $\varrho \in \mathcal{L}(\mathcal{H})$. Namely, if $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ is any instrument, then $\mathcal{I}_x(\varrho) = \mathcal{I}_x(id(\varrho))$ for all $x \in \Omega$. By noting that the identity channel can be viewed as an instrument with only a single outcome, we see that the previous equation is of the same form as Eq. (6). Thus, $id \rightarrow \mathcal{I}$ for any instrument \mathcal{I} . Similarly, if \mathcal{J} is an instrument that is postprocessing equivalent with id so that in particular $\mathcal{J} \rightarrow id$, then by the transitivity of the postprocessing it follows that also $\mathcal{J} \rightarrow \mathcal{I}$ for any instrument

\mathcal{I} . Moreover, if \mathcal{I} is an instrument such that $\mathcal{I} \rightarrow \mathcal{J}$, then \mathcal{I} must actually be postprocessing equivalent with \mathcal{J} (and with id) so that by definition \mathcal{J} is postprocessing clean.

Let then an instrument $\mathcal{R} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ be postprocessing clean. By the previous consideration, we have that $id \rightarrow \mathcal{R}$, so that because \mathcal{R} is postprocessing clean it follows that also $\mathcal{R} \rightarrow id$. Thus, any postprocessing clean instrument is postprocessing equivalent with id . ■

From the proof of Proposition 2 we see that every instrument can be postprocessed from any instrument that is postprocessing equivalent with the identity channel. Furthermore, we can characterize the equivalence class of id .

Proposition 3. An instrument $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ is equivalent with the identity channel on \mathcal{H} if and only if for all $x \in \Omega$ and $\varrho \in \mathcal{S}(\mathcal{H})$ we have that

$$\mathcal{I}_x(\varrho) = \sum_{i=1}^{n_x} p_{xi} V_{xi} \varrho V_{xi}^* \tag{7}$$

for some probability distribution $(p_{xi})_{x,i}$ over $x \in \Omega$, $i \in \{1, \dots, n_x\}$ and some isometries $V_{xi} : \mathcal{H} \rightarrow \mathcal{K}$ such that $V_{xi}^* V_{xj} = 0$ for all $i \neq j$ for all $x \in \Omega$.

Proof. First, let \mathcal{I} be of the form of Eq. (7) for some probability distribution p_{xi} and isometries V_{xi} . Since every instrument can be postprocessed from the identity channel, we have that $id \rightarrow \mathcal{I}$. We define channels (1-outcome instruments) $\mathcal{R}^{(x)} \in \text{Ins}(\{\Omega\}, \mathcal{K}, \mathcal{H})$ by $\mathcal{R}^{(x)}(\varrho) = \sum_{i=1}^{n_x} V_{xi}^* \varrho V_{xi} + \text{tr}[\Pi_x \varrho] |\psi\rangle\langle\psi|$ for all $x \in \Omega$ and $\varrho \in \mathcal{S}(\mathcal{K})$, where $|\psi\rangle$ is an arbitrary unit vector in \mathcal{H} and $\Pi_x = I_{\mathcal{K}} - \sum_{i=1}^{n_x} V_{xi} V_{xi}^*$. We remind the reader that $\Pi_{xi} = V_{xi} V_{xi}^*$ are orthogonal projectors since V_{xi} are isometries and $V_{xj}^* V_{xi} = 0$ for $i \neq j$. If we denote the orthonormal vectors spanning the subspace on which Π_x projects by $|e_{xk}\rangle$, for $k = 1, \dots, \text{tr}[\Pi_x]$, and furthermore if we set $K_{xk} = |\psi\rangle\langle e_{xk}|$, then $\text{tr}[\Pi_x \varrho] |\psi\rangle\langle\psi| = \sum_k K_{xk} \varrho K_{xk}^*$ and also $\{V_{xi}^*\}_{i=1}^{n_x} \cup \{K_{xk}\}_{k=1}^{\text{tr}[\Pi_x]}$ are Kraus operators of the channel $\mathcal{R}^{(x)}$. Indeed, the following calculation shows that $\mathcal{R}^{(x)}$ is trace-preserving:

$$\sum_i V_{xi} V_{xi}^* + \sum_k K_{xk}^* K_{xk} = \sum_i \Pi_{xi} + \Pi_x = I_{\mathcal{K}}.$$

Moreover, $V_{xi}^* |e_{xk}\rangle = 0$ for all $i \in \{1, \dots, n_x\}$, $k \in \{1, \dots, \text{tr}[\Pi_x]\}$, and $x \in \Omega$, which follows from the definition of Π_x . We see that

$$\begin{aligned} \sum_{x \in \Omega} \mathcal{R}^{(x)}(\mathcal{I}_x(\varrho)) &= \sum_{x \in \Omega} \sum_{j=1}^{n_x} V_{xj}^* \left(\sum_{i=1}^{n_x} p_{xi} V_{xi} \varrho V_{xi}^* \right) V_{xj} \\ &= \sum_{x,i} p_{xi} \varrho = id(\varrho). \end{aligned}$$

Hence, $\mathcal{I} \rightarrow id$ so that $\mathcal{I} \leftrightarrow id$.

The second part of the proof is an adaptation of Theorem 2.1 of Ref. [27] for quantum operations forming a quantum instrument. Let $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ be postprocessing equivalent with the identity channel so that there exist channels $\mathcal{J}^{(x)} \in \text{Ins}(\{\Omega\}, \mathcal{K}, \mathcal{H})$ such that $\sum_{x \in \Omega} \mathcal{J}^{(x)} \circ \mathcal{I}_x = id$ on \mathcal{H} . Let \mathcal{I}_x and $\mathcal{J}^{(x)}$ have minimal Kraus operators $\{A_{ix}\}_i$ and $\{B_j^{(x)}\}_j$, respectively, so that

$$\varrho = id(\varrho) = \sum_{x,i,j} B_j^{(x)} A_{ix} \varrho A_{ix}^* (B_j^{(x)})^*$$

for all $\varrho \in \mathcal{S}(\mathcal{H})$.

By the unitary equivalence of the Kraus operators it follows that there exists a set of complex numbers $\{u_{ijx}\}_{ijx} \subset \mathbb{C}$ such that

$$B_j^{(x)} A_{ix} = u_{ijx} I_{\mathcal{H}} \quad \forall i, j, x, \quad (8)$$

and $\sum_{i,j,x} |u_{ijx}|^2 = 1$. Thus, by multiplying the Eq. (8) by its adjoint on the left, summing over j and noting that $\sum_j (B_j^{(x)})^* B_j^{(x)} = I$, we see that $A_{i'x}^* A_{ix} = \beta_{i'i}^x I_{\mathcal{H}}$ for all i', i, x , where we denoted $\beta_{i'i}^x := \sum_j u_{i'jx}^* u_{ijx}$. For each $x \in \Omega$ we see that $\beta_{i'i}^x$ is a positive-semidefinite matrix which can be diagonalized. Let us denote its eigenvalues and eigenvectors by γ_k^x and v_{ki}^x , respectively. Thus, we have $\beta_{i'i}^x = \sum_k \gamma_k^x v_{ki}^x (v_{ki}^x)^*$. Let us define a new set of Kraus operators for each quantum operation \mathcal{I}_x via the relations

$$C_{kx} = \sum_i v_{ki}^x A_{ix}.$$

Due to the unitarity of the matrix $\{v_{ki}^x\}_{ik}$, the Kraus operators $\{C_{kx}\}_k$ also represent the quantum operation \mathcal{I}_x . The important property of operators C_{kx} is that their range spaces are orthogonal, as is proved via the following calculation:

$$\begin{aligned} C_{k'x}^* C_{kx} &= \left(\sum_{i'} (v_{k'i'}^x)^* A_{i'x}^* \right) \left(\sum_i v_{ki}^x A_{ix} \right) \\ &= \sum_{i',i} (v_{k'i'}^x)^* v_{ki}^x (A_{i'x}^* A_{ix}) \\ &= \sum_{i',i} (v_{k'i'}^x)^* v_{ki}^x \beta_{i'i}^x I_{\mathcal{H}} \\ &= \delta_{k'k} \gamma_k^x I_{\mathcal{H}}, \end{aligned} \quad (9)$$

where δ is the Kronecker symbol. Equation (9) implies that the singular value decomposition of C_{kx} has the form $C_{kx} = (\gamma_k^x)^{1/2} \sum_l |y_{kl}^x\rangle \langle u_{kl}^x|$, where $|u_{kl}^x\rangle$ is an orthonormal basis of \mathcal{H} and $|y_{kl}^x\rangle$ are orthonormal vectors in \mathcal{K} . This means that C_{kx} equals $(\gamma_k^x)^{1/2}$ times unitary embedding $V_{xk} \equiv \sum_l |y_{kl}^x\rangle \langle u_{kl}^x|$ of \mathcal{H} into \mathcal{K} . Since operations \mathcal{I}_x form an instrument we have that

$$I_{\mathcal{H}} = \sum_x \sum_k C_{kx}^* C_{kx} = \left(\sum_x \sum_k \gamma_k^x \right) I_{\mathcal{H}}.$$

This allow us to define probability distribution $p_{xi} = \gamma_i^x$ for all $x \in \Omega$ and $i = 1, \dots, n_x$ so that altogether we have

$$\mathcal{I}_x(\varrho) = \sum_{k=1}^{n_x} C_{kx} \varrho C_{kx}^* = \sum_{k=1}^{n_x} \gamma_k^x V_{xk} \varrho V_{xk}^* = \sum_{k=1}^{n_x} p_{xk} V_{xk} \varrho V_{xk}^*,$$

which concludes the proof. \blacksquare

As we have seen from Proposition 2, the equivalence class of the identity channel is the unique greatest element with respect to the postprocessing partial order. What about the least element? We can show the following:

Proposition 4. Any trash-and-prepare instrument can be postprocessed from any instrument. The equivalence class of any trash-and-prepare instrument only consists of trash-and-prepare instruments.

Proof. Let $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ be any instrument and let $\mathcal{T} \in \text{Ins}(\Lambda, \mathcal{H}, \mathcal{V})$ be a trash-and-prepare instrument defined

as $\mathcal{T}_y(\varrho) = \text{tr}[\varrho] p_y \xi_y$ for all $y \in \Lambda$ for some probability distribution $(p_y)_y$ over Λ and some set of states $\{\xi_y\}_y$. To show $\mathcal{I} \rightarrow \mathcal{T}$, we see that we can use nearly the same trash-and-prepare instrument as a postprocessing, i.e., we set $\mathcal{R}^{(x)} \in \text{Ins}(\Lambda, \mathcal{K}, \mathcal{V})$ such that $\mathcal{R}_y^{(x)}(\varrho) = \text{tr}[\varrho] p_y \xi_y$ for all $x \in \Omega$. It follows that

$$\begin{aligned} \sum_{x \in \Omega} \mathcal{R}_y^{(x)}(\mathcal{I}_x(\varrho)) &= \sum_{x \in \Omega} \text{tr}[\mathcal{I}_x(\varrho)] p_y \xi_y \\ &= \text{tr} \left[\left(\sum_{x \in \Omega} \mathcal{I}_x \right) (\varrho) \right] p_y \xi_y \\ &= \text{tr}[\varrho] p_y \xi_y \end{aligned}$$

for all $y \in \Lambda$ and $\varrho \in \mathcal{L}(\mathcal{H})$. This proves the first statement.

For the latter part we note that the argument above applies also for other trash-and-prepare instrument so that any trash-and-prepare instrument can be postprocessed from any other trash-and-prepare instrument. To see that this is the whole equivalence class, let $\mathcal{Q}^{(y)} \in \text{Ins}(\Gamma, \mathcal{V}, \mathcal{K})$ be any set of postprocessing instruments for the formerly defined trash-and-prepare instrument \mathcal{T} . Then

$$\sum_{y \in \Lambda} \mathcal{Q}_z^{(y)}(\mathcal{T}_y(\varrho)) = \text{tr}[\varrho] \sum_{y \in \Lambda} p_y \mathcal{Q}_z^{(y)}(\xi_y) = \text{tr}[\varrho] q_z \sigma_z$$

for all $z \in \Gamma$, where we have defined $q_z = \text{tr}[\sum_{y \in \Lambda} p_y \mathcal{Q}_z^{(y)}(\xi_y)]$ and $\sigma_z = \sum_{y \in \Lambda} p_y / q_z \mathcal{Q}_z^{(y)}(\xi_y) \in \mathcal{S}(\mathcal{H})$ when $q_z \neq 0$ and $\sigma_z = \sigma$ for some fixed $\sigma \in \mathcal{S}(\mathcal{H})$ when $q_z = 0$. Hence, postprocessing a trash-and-prepare instrument just leads to another trash-and-prepare instrument. \blacksquare

We can reformulate the previous result as follows: \mathcal{J} is an instrument such that $\mathcal{I} \rightarrow \mathcal{J}$ for all other instruments \mathcal{I} if and only if \mathcal{J} is a trash-and-prepare instrument. Indeed, as in the proof of the previous proposition, if \mathcal{J} is trash-and-prepare, it can be postprocessed from any other instrument \mathcal{I} . Conversely, if \mathcal{J} can be postprocessed from any instrument then it can be postprocessed from some trash-and-prepare instrument, from which it follows by Proposition 4 that \mathcal{J} also must be trash-and-prepare.

B. Indecomposable instruments

In the case of POVMs the indecomposable (rank-1) POVMs formed the set of maximal elements of the equivalence classes in the postprocessing order. Although we already characterized the single maximal element for instruments, we will see that considering indecomposability in the case of instruments gives us some resemblance to the POVM case.

Definition 6. A (nonzero) quantum operation \mathcal{M} is *indecomposable* if $\mathcal{M} = \mathcal{N} + \mathcal{N}'$ for some other quantum operations $\mathcal{N}, \mathcal{N}'$ only when $\mathcal{N} = \mu \mathcal{M}$ and $\mathcal{N}' = \mu' \mathcal{M}$ for some $\mu, \mu' > 0$. We call a quantum instrument indecomposable if all of its nonzero operations are indecomposable.

We can show the following characterization of indecomposable instruments:

Proposition 5. A quantum operation is indecomposable if and only if it has (Kraus) rank equal to one.

Proof. First let $\mathcal{N} : \mathcal{L}(\mathcal{H}) \rightarrow \mathcal{L}(\mathcal{K})$ be an indecomposable operation with a minimal Kraus decomposition $\mathcal{N}(\varrho) = \sum_{i=1}^r K_i \varrho K_i^*$ for all $\varrho \in \mathcal{L}(\mathcal{H})$ with Kraus rank r . We can de-

fine operations \mathcal{N}_i by setting $\mathcal{N}_i(\varrho) = K_i \varrho K_i^*$ for all $\varrho \in \mathcal{L}(\mathcal{H})$ and $i = 1, \dots, r$, and we see that $\mathcal{N} = \sum_{i=1}^r \mathcal{N}_i$. Because \mathcal{N} is indecomposable, there exists $v_i > 0$ such that $\mathcal{N}_i = v_i \mathcal{N}$ for all $i = 1, \dots, r$ so that the Kraus rank of \mathcal{N} must be one.

Let then $\mathcal{N} : \mathcal{L}(\mathcal{H}) \rightarrow \mathcal{L}(\mathcal{K})$ be an operation with only one Kraus operator, i.e., $\mathcal{N}(\varrho) = K \varrho K^*$ for all $\varrho \in \mathcal{L}(\mathcal{H})$. Let then \mathcal{Q} and \mathcal{R} be nonzero operations such that $\mathcal{N} = \mathcal{Q} + \mathcal{R}$ with Kraus decompositions $\mathcal{Q}(\varrho) = \sum_i A_i \varrho A_i^*$ and $\mathcal{R}(\varrho) = \sum_j B_j \varrho B_j^*$ for all $\varrho \in \mathcal{L}(\mathcal{H})$. By the unitary equivalence of the Kraus operators there exists complex numbers $\{u_i\}_i, \{v_j\}_j \subset \mathbb{C}$ such that $A_i = u_i K$ and $B_j = v_j K$ with $\sum_i |u_i|^2 + \sum_j |v_j|^2 = 1$. Hence,

$$\begin{aligned} \mathcal{Q}(\varrho) &= \sum_i A_i \varrho A_i^* = \left(\sum_i |u_i|^2 \right) K \varrho K^* = u \mathcal{N}(\varrho), \\ \mathcal{R}(\varrho) &= \sum_j B_j \varrho B_j^* = \left(\sum_j |v_j|^2 \right) K \varrho K^* = v \mathcal{N}(\varrho), \end{aligned}$$

where we have denoted $u := \sum_i |u_i|^2 > 0$ and $v := \sum_j |v_j|^2 > 0$. Thus, \mathcal{N} is indecomposable. ■

Example 5. An important class of indecomposable instruments are the *Lüders instruments*: if $A \in \mathcal{O}(\Omega, \mathcal{H})$ is a measurement on \mathcal{H} , then the corresponding Lüders instrument $\mathcal{I}^A \in \text{Ins}(\Omega, \mathcal{H})$ with induced POVM A is defined as $\mathcal{I}_x^A(\varrho) = \sqrt{A(x)} \varrho \sqrt{A(x)}$ for all $x \in \Omega$ and $\varrho \in \mathcal{S}(\mathcal{H})$. Since $\sqrt{A(x)}$ is the only Kraus operator of \mathcal{I}_x^A , by Proposition 5 we see that Lüders instruments are indecomposable.

As we saw, just as with POVMs, the indecomposable elements are those that have (Kraus) rank equal to one. Although from the previous characterization it is obvious that the indecomposable instruments in general are not maximal elements, we will see that they can be used to produce every instrument as a postprocessing of them.

Namely, if $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ has a Kraus decomposition $\mathcal{I}_x(\varrho) = \sum_{i=1}^{n_x} K_{ix} \varrho K_{ix}^*$ for all $x \in \Omega$ and $\varrho \in \mathcal{S}(\mathcal{H})$ for some $n_x \in \mathbb{N}$, then it can be (classically) postprocessed from the instrument $\hat{\mathcal{I}}$ that is constructed from the single Kraus operators of \mathcal{I} , i.e., $\hat{\mathcal{I}}_{(i,x)}(\varrho) = K_{ix} \varrho K_{ix}^*$ for all $i \in \{1, \dots, n_x\}, x \in \Omega$ and $\varrho \in \mathcal{S}(\mathcal{H})$. By using the (classical) postprocessing instruments $\mathcal{R}_x^{(i,x)}$ defined as $\mathcal{R}_x^{(i,x)} = \delta_{xx'} id_{\mathcal{K}}$ for all $i \in \{1, \dots, n_x\}$ and $x, x' \in \Omega$, we see that

$$\begin{aligned} \sum_{x \in \Omega} \sum_{i=1}^{n_x} \mathcal{R}_x^{(i,x)}(\hat{\mathcal{I}}_{(i,x)}(\varrho)) &= \sum_{x \in \Omega} \sum_{i=1}^{n_x} \delta_{xx'} K_{ix} \varrho K_{ix}^* \\ &= \sum_{i=1}^{n_x} K_{ix'} \varrho K_{ix'}^* \\ &= \mathcal{I}_x(\varrho) \end{aligned} \tag{10}$$

for all $\varrho \in \mathcal{S}(\mathcal{H})$ and $x' \in \Omega$ so that $\hat{\mathcal{I}} \rightarrow \mathcal{I}$.

We call $\hat{\mathcal{I}}$ the *detailed instrument* of \mathcal{I} and note that any instrument has many detailed instruments depending on their Kraus decomposition, but all the detailed instruments are indecomposable.

We note that it is also known that any instrument \mathcal{I} can be postprocessed from the Lüders instrument that has the same induced POVM as \mathcal{I} [28]. Since also Lüders instruments are

indecomposable this is another way to see that every instrument can be postprocessed from indecomposable instruments.

To conclude, every instrument can be postprocessed from its (indecomposable) detailed instrument. Next, we give a sufficient condition when the inverse statement also holds so that an instrument under this condition is postprocessing equivalent to its detailed instrument.

Proposition 6. An instrument $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ with Kraus decomposition $\mathcal{I}_x(\varrho) = \sum_{i=1}^{n_x} K_{ix} \varrho K_{ix}^*$ is postprocessing equivalent with its detailed instrument if $K_{ix}^* K_{jx} = 0$ for all $i \neq j$ and $x \in \Omega$.

Proof. We have shown in Eq. (10) that $\hat{\mathcal{I}} \rightarrow \mathcal{I}$ for any instrument \mathcal{I} . Next we focus on proving the relation $\mathcal{I} \rightarrow \hat{\mathcal{I}}$ under the orthogonality condition of the Kraus operators of \mathcal{I}_x for all $x \in \Omega$.

Suppose $K_{ix}^* K_{jx} = 0$ for all $i \neq j$ for all $x \in \Omega$. If π_{ix} is the projector onto the image of K_{ix} we see that together they are mutually orthogonal for all $x \in \Omega$. Define projectors $\{\Pi_{ix}\}$ as $\Pi_{ix} = I - \sum_{i \neq 1} \pi_{ix}$, $\Pi_{ix} = \pi_{ix}$ for $i \neq 1$ and all $x \in \Omega$. Define instruments $\mathcal{R}^{(x)} \in \text{Ins}(\{1, \dots, n_x\} \times \Omega, \mathcal{H}, \mathcal{K})$ as $\mathcal{R}_{(i,y)}^{(x)}(\varrho) = \delta_{xy} \Pi_{iy} \varrho \Pi_{iy}$. Now we see that for $i \neq 1$, we have that

$$\sum_{x \in \Omega} \mathcal{R}_{(i,y)}^{(x)}(\mathcal{I}_x(\varrho)) = \sum_{x \in \Omega} \sum_{j=1}^{n_x} \delta_{xy} \pi_{iy} K_{jx} \varrho K_{jx}^* \pi_{iy} = K_{iy} \varrho K_{iy},$$

and for $i = 1$ we have

$$\begin{aligned} \sum_{x \in \Omega} \mathcal{R}_{(1,y)}^{(x)}(\mathcal{I}_x(\varrho)) &= \sum_{x \in \Omega} \sum_{j=1}^{n_x} \left[\delta_{xy} \left(I - \sum_{i \neq 1} \pi_{iy} \right) K_{jx} \varrho K_{jx}^* \left(I - \sum_{i' \neq 1} \pi_{i'y} \right) \right] \\ &= \sum_{j=1}^{n_y} K_{jy} \varrho K_{jy}^* - \sum_{j=1}^{n_y} \sum_{i \neq 1} \pi_{iy} K_{jy} \varrho K_{jy}^* \\ &\quad - \sum_{j=1}^{n_y} \sum_{i' \neq 1} K_{jy} \varrho K_{jy}^* \pi_{i'y} + \sum_{j=1}^{n_y} \sum_{i, i' \neq 1} \pi_{iy} K_{jy} \varrho K_{jy}^* \pi_{i'y} \\ &= \sum_j K_{jy} \varrho K_{jy}^* - \sum_{j \neq 1} K_{jy} \varrho K_{jy}^* \\ &\quad - \sum_{j \neq 1} K_{jy} \varrho K_{jy}^* + \sum_{j \neq 1} K_{jy} \varrho K_{jy}^* \\ &= K_{1y} \varrho K_{1y}^*, \end{aligned}$$

so that $\sum_{x \in \Omega} \mathcal{R}_{(i,y)}^{(x)}(\mathcal{I}_x(\varrho)) = K_{iy} \varrho K_{iy}^* = \hat{\mathcal{I}}_{(i,y)}(\varrho)$ for all $i \in \{1, \dots, n_y\}$ and $y \in \Omega$ for all $\varrho \in \mathcal{S}(\mathcal{H})$. Hence, $\mathcal{I} \rightarrow \hat{\mathcal{I}}$. ■

For measure-and-prepare instruments we can use the previous result to show a necessary and sufficient condition for being equivalent with an indecomposable instrument.

Proposition 7. A measure-and-prepare instrument $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ is postprocessing equivalent with an indecomposable instrument if and only if $A^{\mathcal{I}} \in \mathcal{O}(\Omega, \mathcal{H})$ is indecomposable.

Proof. Let us consider a measure-and-prepare instrument $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ that is of the form $\mathcal{I}_x(\varrho) = \text{tr}[A(x)\varrho] \xi_x$ for all $x \in \Omega$ and $\varrho \in \mathcal{S}(\mathcal{H})$ for some POVM $A \in \mathcal{O}(\Omega, \mathcal{H})$ and some states $\{\xi_x\}_{x \in \Omega} \subset \mathcal{S}(\mathcal{K})$. For each $x \in \Omega$, let

us consider the spectral decomposition of the state $\xi_x = \sum_{i=1}^{n_x} p_{ix} |\varphi_{ix}\rangle\langle\varphi_{ix}|$, where $p_{ix} \geq 0$, $\sum_i p_{ix} = 1$ and $\{\varphi_{ix}\}_i$ is a set of orthonormal vectors in \mathcal{K} for all $x \in \Omega$. Similarly for each $x \in \Omega$, we can write $\mathbf{A}(x)$ as $\mathbf{A}(x) = \sum_{j=1}^{m_x} q_{jx} |\psi_{jx}\rangle\langle\psi_{jx}|$ for some orthogonal set of vectors $\{\psi_{jx}\}_j$ in \mathcal{H} and some positive numbers $\{q_{jx}\}_j$. Let us define operators $K_{ijx} := \sqrt{p_{ix} q_{jx}} |\varphi_{ix}\rangle\langle\psi_{jx}|$ for all $i \in \{1, \dots, n_x\}$ and $j \in \{1, \dots, m_x\}$ for each $x \in \Omega$. One can confirm that $\sum_{i,j} K_{ijx}^* K_{ijx} = \mathbf{A}(x)$ for all $x \in \Omega$ so that $\sum_{i,j,x} K_{ijx}^* K_{ijx} = \mathcal{I}_{\mathcal{H}}$, and that $\sum_{i=1}^{n_x} \sum_{j=1}^{m_x} K_{ijx} \varrho K_{ijx}^* = \mathcal{I}_x(\varrho)$ for all $x \in \Omega$ so that $\{K_{ijx}\}_{i,j}$ is a set of Kraus operators for \mathcal{I}_x .

Let first \mathbf{A} be indecomposable, i.e., rank 1, so that $m_x = 1$ for all $x \in \Omega$ and we can omit the index j in the previous consideration and thus $K_{ix} := \sqrt{p_{ix} q_{1x}} |\varphi_{ix}\rangle\langle\psi_{1x}|$ form the set of Kraus operators for \mathcal{I}_x for all $x \in \Omega$. We see that $K_{ix}^* K_{i'x} = 0$ for all $i \neq i'$ for all $x \in \Omega$ and thus by Proposition 6 the instrument \mathcal{I} is equivalent with its detailed instrument related to that Kraus decomposition. From the proof of Proposition 6 we see that the simulator instruments $\mathcal{R}^{(x)} \in \text{Ins}(\{1, \dots, n_x\} \times \Omega, \mathcal{K})$ can be written as $\mathcal{R}_{(i,y)}^{(x)}(\varrho) = \delta_{xy} \Pi_{ix} \varrho \Pi_{ix}$, where the projectors $\{\Pi_{ix}\}_i$ are defined as $\Pi_{ix} = |\varphi_{ix}\rangle\langle\varphi_{ix}|$ for all $i \neq 1$ and $\Pi_{1x} = I_{\mathcal{K}} - \sum_{i \neq 1} |\varphi_{ix}\rangle\langle\varphi_{ix}|$ for all $x \in \Omega$.

Let then \mathbf{A} not be rank 1, i.e., there exists $x' \in \Omega$ such that $m_{x'} \geq 2$. Suppose that \mathcal{I} is equivalent to some indecomposable instrument $\mathcal{J} \in \text{Ins}(\Lambda, \mathcal{H}, \mathcal{V})$ with Kraus decomposition $\mathcal{J}_y(\varrho) = L_y \varrho L_y^*$ for all $y \in \Lambda$ and $\varrho \in \mathcal{S}(\mathcal{H})$. Thus, there exist instruments $\mathcal{R}^{(x)} \in \text{Ins}(\Lambda, \mathcal{K}, \mathcal{V})$ with Kraus operators $R_{ky}^{(x)}$ such that

$$\begin{aligned} L_y \varrho L_y^* &= \mathcal{J}_y(\varrho) = \sum_{x \in \Omega} \mathcal{R}_y^{(x)}(\mathcal{I}_x(\varrho)) \\ &= \sum_{i,j,k,x} R_{ky}^{(x)} K_{ijx} \varrho K_{ijx}^* (R_{ky}^{(x)})^* \end{aligned}$$

for all $y \in \Lambda$ and $\varrho \in \mathcal{S}(\mathcal{H})$. From the unitary equivalence of the Kraus operators it follows that there exist complex numbers $u_{ijkxy} \in \mathbb{C}$ such that

$$R_{ky}^{(x)} K_{ijx} = u_{ijkxy} L_y \quad (11)$$

for all i, j, k, x, y and $\sum_{i,j,k,x} |u_{ijkxy}|^2 = 1$ for all $y \in \Lambda$. If we denote the induced POVM of $\mathcal{R}^{(x)}$ by $\mathbf{R}^{(x)}$, i.e., $\mathbf{R}^{(x)}(y) = \sum_k (R_{ky}^{(x)})^* R_{ky}^{(x)}$ for all $y \in \Lambda$ for all $x \in \Omega$, by multiplying Eq. (11) by its adjoint from the left and by summing over indices i, k , we see that

$$\begin{aligned} \left(\sum_{i,k} |u_{ijkxy}|^2 \right) L_y^* L_y &= \sum_{i,k} K_{ijx}^* (R_{ky}^{(x)})^* R_{ky}^{(x)} K_{ijx} \\ &= \sum_i p_{ix} \left\langle \varphi_{ix} \left| \sum_k (R_{ky}^{(x)})^* R_{ky}^{(x)} \varphi_{ix} \right. \right\rangle \\ &\quad \times q_{jx} |\psi_{jx}\rangle\langle\psi_{jx}| \\ &= \text{tr}[\mathbf{R}^{(x)}(y) \xi_x] q_{jx} |\psi_{jx}\rangle\langle\psi_{jx}| \end{aligned}$$

for all $j \in \{1, \dots, m_x\}$, $x \in \Omega$ and $y \in \Lambda$.

As we mentioned, since \mathbf{A} is not rank 1, there exists $x' \in \Omega$ such that $m_{x'} \geq 2$ so that there are indices $j', j'' \in \{1, \dots, m_{x'}\}$ such that $q_{j'x'}, q_{j''x'} > 0$. Also, since $\mathbf{R}^{(x')}$ is a POVM, there exists $y' \in \Lambda$ such that $\text{tr}[\mathbf{R}^{(x')}(y') \xi_{x'}] \neq 0$. Thus,

we have that

$$\begin{aligned} \left(\sum_{i,k} |u_{ij'kx'y'}|^2 \right) L_{y'}^* L_{y'} &= \text{tr}[\mathbf{R}^{(x')}(y') \xi_{x'}] q_{j'x'} |\psi_{j'x'}\rangle\langle\psi_{j'x'}|, \\ \left(\sum_{i,k} |u_{ij''kx'y'}|^2 \right) L_{y'}^* L_{y'} &= \text{tr}[\mathbf{R}^{(x')}(y') \xi_{x'}] q_{j''x'} |\psi_{j''x'}\rangle\langle\psi_{j''x'}|, \end{aligned}$$

which leads to a contradiction since $|\psi_{j'x'}\rangle\langle\psi_{j'x'}|$ is not proportional to $|\psi_{j''x'}\rangle\langle\psi_{j''x'}|$. Hence, if \mathbf{A} is not rank 1, then \mathcal{I} is not postprocessing equivalent to any indecomposable instrument. \blacksquare

This result shows that there are instruments that are not equivalent with an indecomposable instruments and this is the case for all trash-and-prepare instruments. On the other hand, it is known that if the induced POVM $\mathbf{A}^{\mathcal{I}}$ of any instrument \mathcal{I} is rank 1, then \mathcal{I} must be a measure-and-prepare instrument [29]. Hence, we get the following corollary:

Corollary 1. Any instrument with an indecomposable induced POVM is postprocessing equivalent with its detailed (indecomposable) instrument.

V. POSTPROCESSING OF THE INDUCED POSITIVE OPERATOR-VALUED MEASURES

So far we have separately considered the postprocessing of POVMs and instruments. Quantum instruments are also a form of measurement, so we can start finding connections between the two postprocessings. When the resulting instrument is indecomposable, we can show the following:

Proposition 8. Let instruments $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ and $\mathcal{J} \in \text{Ins}(\Lambda, \mathcal{H}, \mathcal{V})$ be such that \mathcal{J} is indecomposable. If $\mathcal{I} \rightarrow \mathcal{J}$, then $\mathbf{A}^{\mathcal{J}} \rightarrow \mathbf{A}^{\mathcal{I}}$.

Proof. If $\mathcal{I} \rightarrow \mathcal{J}$, then there exists postprocessings $\mathcal{R}^{(x)} \in \text{Ins}(\Lambda, \mathcal{K}, \mathcal{V})$ for all $x \in \Omega$ such that $\mathcal{J}_y = \sum_{x \in \Omega} \mathcal{R}_y^{(x)} \circ \mathcal{I}_x$ for all $y \in \Lambda$. Let $\{A_{ix}\}_i$, B_y and $\{R_{ky}^{(x)}\}_k$ be Kraus operators for \mathcal{I}_x , \mathcal{J}_y , and $\mathcal{R}_y^{(x)}$ respectively. Thus,

$$B_y \varrho B_y^* = \sum_{i,k,x} R_{ky}^{(x)} A_{ix} \varrho A_{ix}^* (R_{ky}^{(x)})^*$$

for all $\varrho \in \mathcal{S}(\mathcal{H})$.

From the unitary equivalence of the Kraus operators it follows that there exists complex numbers $\{u_{ikxy}\}_{i,k,x,y} \subset \mathbb{C}$ such that $R_{ky}^{(x)} A_{ix} = u_{ikxy} B_y$ for all i, k, x, y and $\sum_{i,k,x} |u_{ikxy}|^2 = 1$ for all $y \in \Lambda$. By multiplying the left side of the previous expression by its adjoint and summing over the indices i, k, y , we see that

$$\sum_i A_{ix}^* A_{ix} = \sum_{i,k,y} A_{ix}^* (R_{ky}^{(x)})^* R_{ky}^{(x)} A_{ix} = \sum_{i,k,y} |u_{ikxy}|^2 B_y^* B_y$$

for all $x \in \Omega$.

Clearly $\mathbf{A}^{\mathcal{I}}(x) = \sum_i A_{ix}^* A_{ix}$ and $\mathbf{A}^{\mathcal{J}}(y) = B_y^* B_y$ for all $x \in \Omega$ and $y \in \Lambda$. If we denote $\nu_{yx} = \sum_{i,k} |u_{ikxy}|^2 \geq 0$, we see that $\sum_{x \in \Omega} \nu_{yx} = 1$ for all $y \in \Lambda$ so that it defines a postprocessing $(\nu_{yx})_{y \in \Lambda, x \in \Omega}$ such that $\mathbf{A}^{\mathcal{I}}(x) = \sum_{y \in \Lambda} \nu_{yx} \mathbf{A}^{\mathcal{J}}(y)$ for all $x \in \Omega$. Hence, $\mathbf{A}^{\mathcal{J}} \rightarrow \mathbf{A}^{\mathcal{I}}$. \blacksquare

We note that the previous claim is not true when \mathcal{J} is not indecomposable. To see this, let us consider the case when \mathcal{J} is a trash-and-prepare instrument. As we have shown in

Proposition 4, then $\mathcal{I} \rightarrow \mathcal{J}$ for any instrument \mathcal{I} , but since $\mathbf{A}^{\mathcal{J}}$ is now a trivial POVM the relation $\mathbf{A}^{\mathcal{J}} \rightarrow \mathbf{A}^{\mathcal{I}}$ does not hold in general.

When both instruments are indecomposable, we can prove even a stronger result, but for that we need a small lemma first.

Lemma 1. Suppose that operators $K : \mathcal{H} \rightarrow \mathcal{K}, L : \mathcal{H} \rightarrow \mathcal{V}$ satisfy $K^*K = cL^*L$ for some $c > 0$. Then there exists an operator $U : \mathcal{V} \rightarrow \mathcal{K}$ such that $K = \sqrt{c}UL$ with the following properties:

- (1) If $\dim \mathcal{K} \geq \dim \mathcal{V}$, then U is an isometry.
- (2) If $\dim \mathcal{K} < \dim \mathcal{V}$, then U is a partial isometry such that its range is the whole \mathcal{K} and $U^*UL = L$.

Proof. Let us start with a small note about partial isometries. Suppose we have an operator $X : \mathcal{H} \rightarrow \mathcal{K}$ defined as $X = \sum_{m=a}^b |e_m\rangle\langle g_m|$, where $\{|g_m\rangle\}_{m=a}^b$ and $\{|e_m\rangle\}_{m=a}^b$ are two sets of orthonormal vectors from \mathcal{H} and \mathcal{K} , respectively. Thus, operator X isometrically transfers subspace $V_{in} = \text{span}(\{|g_m\rangle\}_{m=a}^b)$ into subspace $V_{out} = \text{span}(\{|e_m\rangle\}_{m=a}^b)$. Suppose $\{|h_m\rangle\}_{m=a}^b$ is another set of orthonormal vectors that span V_{in} . Consequently, a projector onto V_{in} can be written as $P = \sum_{m=a}^b |g_m\rangle\langle g_m| = \sum_{m=a}^b |h_m\rangle\langle h_m|$. Clearly,

$$X = XP = \sum_{m=a}^b |e_m\rangle\langle g_m| \sum_{n=a}^b |h_n\rangle\langle h_n| = \sum_{n=a}^b |\tilde{e}_n\rangle\langle h_n|, \quad (12)$$

where vectors $|\tilde{e}_n\rangle = \sum_{m=a}^b \langle g_m | h_n \rangle |e_m\rangle$ are orthonormal, as one can easily check. Thus, we see that X can be also seen as a (linear) isometric transformation of orthonormal vectors $|h_n\rangle$ onto orthonormal vectors $|\tilde{e}_n\rangle$.

Next, we consider singular value decompositions of operators K and L

$$K = \sum_{m=1}^k \lambda_m |e_m\rangle\langle g_m|, \quad L = \sum_{n=1}^l \mu_n |f_n\rangle\langle h_n|, \quad (13)$$

where the singular values $\lambda_m, \mu_n > 0$ are arranged in the decreasing order and $\{|g_m\rangle\}_{m=1}^k, \{|e_m\rangle\}_{m=1}^k, \{|h_n\rangle\}_{n=1}^l, \{|f_n\rangle\}_{n=1}^l$ are orthonormal vectors in the corresponding Hilbert spaces and we assume they were extended to form an orthonormal basis in each of the spaces. The equality $K^*K = cL^*L$ can be now written as

$$\sum_{m=1}^k \lambda_m^2 |g_m\rangle\langle g_m| = \sum_{n=1}^l c\mu_n^2 |h_n\rangle\langle h_n|.$$

Both the left and right sides have form of a spectral decomposition for the same positive-semidefinite operator K^*K . This has important consequences for the singular value decompositions (13). First of all, $k = l$ and $\lambda_m = \sqrt{c}\mu_m$. If some of the eigenvalues of K^*K (or equivalently singular values of K or L) are degenerate then for every such eigenspace defined by eigenvalue $\lambda_a = \lambda_{a+1} = \dots = \lambda_b$ we have that

$$P_{\lambda_a} = \sum_{m=a}^b |g_m\rangle\langle g_m| = \sum_{m=a}^b |h_m\rangle\langle h_m|,$$

i.e., both $\{|g_m\rangle\}_{m=a}^b$ and $\{|h_m\rangle\}_{m=a}^b$ are an orthonormal basis of this eigenspace. Using the considerations about partial isometries [specifically, Eq. (12)] from the beginning of this proof

we can rewrite operator K as

$$K = \sum_{m=1}^k \sqrt{c}\mu_m |\tilde{e}_m\rangle\langle h_m|. \quad (14)$$

We extend orthonormal vectors $\{|\tilde{e}_m\rangle\}_{m=1}^k$ to form an orthonormal basis of \mathcal{K} . At this point we have to consider separately two cases based on the relation between dimensions of \mathcal{K} and \mathcal{V} . First, let us consider $\dim \mathcal{K} \geq \dim \mathcal{V}$. We note that $k \leq \dim \mathcal{V} \leq \dim \mathcal{K}$ and before we defined $\{|f_n\rangle\}_{n=1}^{\dim \mathcal{V}}$ as a complete orthonormal basis of \mathcal{V} . We can now define operator $U : \mathcal{V} \rightarrow \mathcal{K}$ as

$$U = \sum_{m=1}^{\dim \mathcal{V}} |\tilde{e}_m\rangle\langle f_m|. \quad (15)$$

Clearly, U is an isometry by definition and direct calculation verifies that $\sqrt{c}UL$ equals K as given in Eq. (14). Next, we consider $\dim \mathcal{K} < \dim \mathcal{V}$. Let us now set

$$U = \sum_{m=1}^{\dim \mathcal{K}} |\tilde{e}_m\rangle\langle f_m|.$$

We now see that U is a partial isometry, whose range is the whole \mathcal{K} . Due to $k \leq \dim \mathcal{K}$ we again have that $\sqrt{c}UL = K$ and one can also easily verify that $U^*UL = L$, which concludes the proof. ■

We can now show the following:

Proposition 9. Indecomposable instruments $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ and $\mathcal{J} \in \text{Ins}(\Lambda, \mathcal{H}, \mathcal{V})$ are equivalent ($\mathcal{I} \leftrightarrow \mathcal{J}$) if and only if their induced POVMs are equivalent ($\mathbf{A}^{\mathcal{I}} \leftrightarrow \mathbf{A}^{\mathcal{J}}$).

Proof. Let us assume that the indecomposable instruments \mathcal{I} and \mathcal{J} are equivalent, i.e., both $\mathcal{I} \rightarrow \mathcal{J}$ and $\mathcal{I} \leftarrow \mathcal{J}$ holds. Using Proposition 8 we get that both $\mathbf{A}^{\mathcal{J}} \rightarrow \mathbf{A}^{\mathcal{I}}$ and $\mathbf{A}^{\mathcal{I}} \rightarrow \mathbf{A}^{\mathcal{J}}$ hold, respectively, which means that $\mathbf{A}^{\mathcal{I}} \leftrightarrow \mathbf{A}^{\mathcal{J}}$.

For the opposite direction we assume that $\mathbf{A}^{\mathcal{I}} \leftrightarrow \mathbf{A}^{\mathcal{J}}$. Without loss of generality we may assume $\dim \mathcal{K} \geq \dim \mathcal{V}$ and that $\mathbf{A}^{\mathcal{J}}$ and $\mathbf{A}^{\mathcal{I}}$ are nonvanishing so that they only consist of nonzero effects. Our first goal is to show that, under the assumption $\mathbf{A}^{\mathcal{I}} \leftrightarrow \mathbf{A}^{\mathcal{J}}$, from the relation $\mathbf{A}^{\mathcal{I}} \rightarrow \mathbf{A}^{\mathcal{J}}$ we get $\mathcal{I} \rightarrow \mathcal{J}$ for indecomposable instruments.

From Proposition 1 we have that there is a postprocessing matrix ν such that

$$\mathbf{A}^{\mathcal{J}}(y) = \sum_{x \in \Omega} \nu_{xy} \mathbf{A}^{\mathcal{I}}(x) \quad (16)$$

for all $y \in \Lambda$ such that $\nu_{xy} \neq 0$ only if $\mathbf{A}^{\mathcal{I}}(x)$ is proportional to $\mathbf{A}^{\mathcal{J}}(y)$ for all $x \in \Omega$ and $y \in \Lambda$. Thus, for all $\nu_{xy} > 0$ there exists $c_{xy} > 0$ such that

$$\mathbf{A}^{\mathcal{I}}(x) = c_{xy} \mathbf{A}^{\mathcal{J}}(y). \quad (17)$$

Inserting Eq. (17) into Eq. (16) we get $\mathbf{A}^{\mathcal{J}}(y) = \sum_{x: \nu_{xy} > 0} \nu_{xy} c_{xy} \mathbf{A}^{\mathcal{J}}(y)$ or, equivalently,

$$\sum_{x: \nu_{xy} > 0} \nu_{xy} c_{xy} = 1 \quad (18)$$

for all $y \in \Lambda$.

If we denote the Kraus operators of the instruments as

$$\mathcal{I}_x(\rho) = A_x \rho A_x^*, \quad \mathcal{J}_y(\rho) = B_y \rho B_y^* \quad (19)$$

then Eq. (17) can be rewritten as

$$A_x^* A_x = c_{xy} B_y^* B_y.$$

Using Lemma 1 we obtain

$$A_x = \sqrt{c_{xy}} U_{xy} B_y,$$

where $U_{xy}^* U_{xy} = I_{\mathcal{V}}$. On the other hand, $\Pi_{xy} \equiv U_{xy} U_{xy}^*$ can be a nontrivial projector on \mathcal{K} . We denote its complement as $\bar{\Pi}_{xy} = I_{\mathcal{K}} - \Pi_{xy}$ and we define orthonormal states $\{|e_k^{xy}\rangle\} \subset \mathcal{K}$ such that $\bar{\Pi}_{xy} = \sum_{k=1}^{m_{xy}} |e_k^{xy}\rangle \langle e_k^{xy}|$ with $m_{xy} = \text{tr}[\bar{\Pi}_{xy}]$. We note that, by construction, $U_{xy}^* |e_k^{xy}\rangle = 0$ for all $k \in \{1, \dots, m_{xy}\}$, $x \in \Omega$ and $y \in \Lambda$. For every $x \in \Omega$ we define instrument $\mathcal{R}^{(x)} \in \text{Ins}(\Lambda, \mathcal{K}, \mathcal{V})$ via the following formula:

$$\mathcal{R}_y^{(x)}(\rho) = R_y^{(x)} \rho (R_y^{(x)})^* + \sum_{k=1}^{m_{xy}} Q_{ky}^{(x)} \rho (Q_{ky}^{(x)})^*,$$

where $R_y^{(x)} = \sqrt{\nu_{xy}} U_{xy}^*$ and $Q_{ky}^{(x)} = \sqrt{\nu_{xy}} |\xi\rangle \langle e_k^{xy}|$ for some fixed unit vector $|\xi\rangle \in \mathcal{V}$. Complete positivity of the instrument is obvious from its definition and we check preservation of the trace via the following calculation:

$$\begin{aligned} & \sum_{y \in \Lambda} \left((R_y^{(x)})^* R_y^{(x)} + \sum_{k=1}^{m_{xy}} (Q_{ky}^{(x)})^* Q_{ky}^{(x)} \right) \\ &= \sum_{y \in \Lambda} \nu_{xy} \left(U_{xy} U_{xy}^* + \sum_{k=1}^{m_{xy}} |e_k^{xy}\rangle \langle e_k^{xy}| \right) \\ &= \sum_{y \in \Lambda} \nu_{xy} (\Pi_{xy} + \bar{\Pi}_{xy}) = I_{\mathcal{K}}. \end{aligned}$$

Let us now evaluate the postprocessing of instrument \mathcal{I} via the instruments $\mathcal{R}^{(x)}$. We obtain

$$\begin{aligned} \sum_{x \in \Omega} \mathcal{R}_y^{(x)}(\mathcal{I}_x(\rho)) &= \sum_{x \in \Omega} \left[R_y^{(x)} A_x \rho A_x^* (R_y^{(x)})^* \right. \\ &\quad \left. + \sum_{k=1}^{m_{xy}} Q_{ky}^{(x)} A_x \rho A_x^* (Q_{ky}^{(x)})^* \right] \\ &= \sum_{x \in \Omega} c_{xy} \nu_{xy} \left[U_{xy}^* U_{xy} B_y \rho B_y^* U_{xy}^* U_{xy} \right. \\ &\quad \left. + \sum_{k=1}^{m_{xy}} \langle e_k^{xy} | U_{xy} B_y \rho B_y^* U_{xy}^* | e_k^{xy} \rangle |\xi\rangle \langle \xi| \right] \\ &= \sum_{x: \nu_{xy} > 0} (c_{xy} \nu_{xy}) B_y \rho B_y^* \\ &= B_y \rho B_y^* = \mathcal{J}_y(\rho), \end{aligned}$$

where we used $U_{xy}^* |e_k^{xy}\rangle = 0$, $U_{xy}^* U_{xy} = I_{\mathcal{V}}$ and Eqs. (18) and (19). Thus, when $\mathcal{A}^{\mathcal{J}} \leftrightarrow \mathcal{A}^{\mathcal{I}}$, we have proved that $\mathcal{A}^{\mathcal{I}} \rightarrow \mathcal{A}^{\mathcal{J}}$ implies $\mathcal{I} \rightarrow \mathcal{J}$.

Our next goal is to prove that the equivalence $\mathcal{A}^{\mathcal{J}} \leftrightarrow \mathcal{A}^{\mathcal{I}}$ for indecomposable instruments implies also $\mathcal{J} \rightarrow \mathcal{I}$. We start by explicitly writing out the equations guaranteed by Proposition 1 for $\mathcal{A}^{\mathcal{J}} \rightarrow \mathcal{A}^{\mathcal{I}}$. Thus, there exists a stochastic

matrix μ such that

$$\mathcal{A}^{\mathcal{I}}(x) = \sum_{y \in \Lambda} \mu_{yx} \mathcal{A}^{\mathcal{J}}(y) \quad (20)$$

for all $x \in \Omega$ such that $\mu_{yx} \neq 0$ only if $\mathcal{A}^{\mathcal{J}}(y)$ is proportional to $\mathcal{A}^{\mathcal{I}}(x)$. Thus, for all $\mu_{yx} > 0$ we have that there exists $d_{yx} > 0$ such that

$$\mathcal{A}^{\mathcal{J}}(y) = d_{yx} \mathcal{A}^{\mathcal{I}}(x). \quad (21)$$

Inserting Eq. (21) into Eq. (20) we get $\mathcal{A}^{\mathcal{I}}(x) = \sum_{y: \mu_{yx} > 0} \mu_{yx} d_{yx} \mathcal{A}^{\mathcal{I}}(x)$ or equivalently

$$\sum_{y: \mu_{yx} > 0} \mu_{yx} d_{yx} = 1 \quad (22)$$

for all $x \in \Omega$.

Using Kraus operators A_x, B_y for \mathcal{I}_x and \mathcal{J}_y , respectively, Eq. (21) can be rewritten as

$$B_y^* B_y = d_{yx} A_x^* A_x.$$

Applying Lemma 1 to this equation, we obtain

$$B_y = \sqrt{d_{yx}} V_{yx} A_x,$$

where V_{yx} is a partial isometry since $\dim \mathcal{V} \leq \dim \mathcal{K}$. On the one hand we have that $V_{yx} V_{yx}^* = I_{\mathcal{V}}$, and on the other hand as a consequence of Lemma 1 we have

$$V_{yx}^* V_{yx} A_x = A_x \quad (23)$$

for all $x \in \Omega$ and $y \in \Lambda$. For every $y \in \Lambda$ we define instrument $\mathcal{G}^{(y)} \in \text{Ins}(\Omega, \mathcal{V}, \mathcal{K})$ via the following formula:

$$\mathcal{G}_x^{(y)}(\rho) = G_x^{(y)} \rho (G_x^{(y)})^*,$$

where $G_x^{(y)} = \sqrt{\mu_{yx}} V_{yx}^*$. We check the trace preservation by evaluating

$$\sum_{x \in \Omega} (G_x^{(y)})^* G_x^{(y)} = \sum_{x \in \Omega} \mu_{yx} V_{yx} V_{yx}^* = \sum_{x \in \Omega} \mu_{yx} I_{\mathcal{V}} = I_{\mathcal{V}}.$$

Finally we calculate the postprocessing of instrument \mathcal{J} via the instruments $\mathcal{G}^{(y)}$. We obtain

$$\begin{aligned} \sum_{y \in \Lambda} \mathcal{G}_x^{(y)}(\mathcal{J}_y(\rho)) &= \sum_{y \in \Lambda} G_x^{(y)} B_y \rho B_y^* (G_x^{(y)})^* \\ &= \sum_{y \in \Lambda} \mu_{yx} d_{yx} V_{yx}^* V_{yx} A_x \rho A_x^* V_{yx} V_{yx}^* \\ &= \sum_{y: \mu_{yx} > 0} (\mu_{yx} d_{yx}) A_x \rho A_x^* \\ &= A_x \rho A_x^* = \mathcal{I}_x(\rho), \end{aligned}$$

where we used Eqs. (23), (22), and (19). This concludes the proof, since we showed that $\mathcal{A}^{\mathcal{J}} \leftrightarrow \mathcal{A}^{\mathcal{I}}$ implies $\mathcal{I} \leftrightarrow \mathcal{J}$ for indecomposable instruments. \blacksquare

From the previous result we see that, in the special case when we have two Lüder's instruments \mathcal{I}^A and \mathcal{I}^B for two POVMs \mathbf{A} and \mathbf{B} , then \mathcal{I}^A and \mathcal{I}^B are equivalent if and only if \mathbf{A} and \mathbf{B} are equivalent.

For measure-and-prepare instruments we already saw in Example 4 that if $\mathcal{I} \rightarrow \mathcal{J}$ and \mathcal{I} is a measure-and-prepare instrument, then $\mathcal{A}^{\mathcal{I}} \rightarrow \mathcal{A}^{\mathcal{J}}$. We can now show that this is

actually both a necessary and sufficient condition in the case when also \mathcal{J} is a measure-and-prepare instrument.

Proposition 10. Let $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ and $\mathcal{J} \in \text{Ins}(\Lambda, \mathcal{H}, \mathcal{V})$ be measure-and-prepare instruments. Then $\mathcal{I} \rightarrow \mathcal{J}$ if and only if $\mathbf{A}^{\mathcal{I}} \rightarrow \mathbf{A}^{\mathcal{J}}$.

Proof. Since \mathcal{I} and \mathcal{J} are measure-and-prepare, then there exists POVMs $\mathbf{A} \in \mathcal{O}(\Omega, \mathcal{H})$ and $\mathbf{B} \in \mathcal{O}(\Lambda, \mathcal{H})$ as well as states $\{\sigma_x\}_{x \in \Omega} \subset \mathcal{S}(\mathcal{K})$ and $\{\xi_y\}_{y \in \Lambda} \subset \mathcal{S}(\mathcal{V})$ such that

$$\begin{aligned}\mathcal{I}_x(\varrho) &= \text{tr}[\mathbf{A}(x)\varrho]\sigma_x, \\ \mathcal{J}_y(\varrho) &= \text{tr}[\mathbf{B}(y)\varrho]\xi_y\end{aligned}$$

for all $x \in \Omega$ and $y \in \Lambda$. We see that $\mathbf{A}^{\mathcal{I}} = \mathbf{A}$ and $\mathbf{A}^{\mathcal{J}} = \mathbf{B}$. By following the steps of Example 4, we see that if $\mathcal{I} \rightarrow \mathcal{J}$, then $\mathbf{A} = \mathbf{A}^{\mathcal{I}} \rightarrow \mathbf{A}^{\mathcal{J}} = \mathbf{B}$.

Now let $\mathbf{A}^{\mathcal{I}} \rightarrow \mathbf{A}^{\mathcal{J}}$, i.e., $\mathbf{A} \rightarrow \mathbf{B}$ so that $\mathbf{B}(y) = \sum_{x \in \Omega} \mu_{xy} \mathbf{A}(x)$ for all $y \in \Lambda$ for some postprocessing $\mu = (\mu_{xy})_{x \in \Omega, y \in \Lambda}$. Let us define instruments $\mathcal{R}^{(x)} \in \text{Ins}(\Lambda, \mathcal{K}, \mathcal{V})$ by setting $\mathcal{R}_y^{(x)}(\varrho) = \mu_{xy} \xi_y$ for all $x \in \Omega$, $y \in \Lambda$, and $\varrho \in \mathcal{S}(\mathcal{K})$. We now see that

$$\begin{aligned}\sum_{x \in \Omega} \mathcal{R}_y^{(x)}(\mathcal{I}_x(\varrho)) &= \sum_{x \in \Omega} \text{tr}[\mathbf{A}(x)\varrho] \mathcal{R}_y^{(x)}(\sigma_x) \\ &= \text{tr} \left[\sum_{x \in \Omega} \mu_{xy} \mathbf{A}(x)\varrho \right] \xi_y \\ &= \text{tr}[\mathbf{B}(y)\varrho] \xi_y = \mathcal{J}_y(\varrho)\end{aligned}$$

for all $y \in \Lambda$. Hence, $\mathcal{I} \rightarrow \mathcal{J}$. \blacksquare

VI. SIMULATION OF INSTRUMENTS

The simulation scheme describes a process of obtaining new devices out of some existing devices by the means of operational manipulations. For example, in the case of measurements (see Refs. [18–21]), from a set of measurement devices new observables can be obtained by classical means of mixing and/or postprocessing the classical outcomes. This can be achieved by giving a probability for each device according to which we use it in a measurement and/or by (classically) postprocessing the obtained measurement outcomes.

The previously described procedure can be used to simulate a measurement device that is not directly at hand and which may be hard to implement by itself. One can consider which observables can be obtained from a single POVM via postprocessing, which effectively characterizes joint measurability [11], and in this sense the above concept of measurement simulability can be considered as a generalization of joint measurability. One can also ask which POVMs are needed to simulate all observables (simulation irreducible measurements [20]), or what can one get out of a given observables with specific properties (projective measurements [19,21], effectively dichotomic measurements [30–33]).

Next, we define a simulation of instruments analogously to the simulation of POVMs by using the postprocessing that was defined and studied in the previous sections. We start by briefly recalling mixing of instruments.

A. Mixing of quantum instruments

For a fixed outcome set Ω and Hilbert spaces \mathcal{H} and \mathcal{K} the set of instruments $\text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ is convex. Namely, if we have devices described by instruments $\{\mathcal{I}^{(i)}\}_{i=1}^n \subset \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$, then we can choose to use device $\mathcal{I}^{(i)}$ with probability p_i in every round of the experiment with some probability distribution $(p_i)_{i=1}^n$. The new instrument \mathcal{I} that is formed as a mixture is then defined as

$$\mathcal{I}_x = \sum_{i=1}^n p_i \mathcal{I}_x^{(i)}$$

for all $x \in \Omega$. Note that we can always consider instruments to have the same outcome sets by just adding zero outcomes to instruments if needed, but they still have to have the same input and output spaces $\mathcal{L}(\mathcal{H})$ and $\mathcal{L}(\mathcal{K})$.

We can also consider other type of mixing where we also keep track of the instrument that was used in each round of the experiment. Then we consider the mixed instrument to have to outcomes, the first outcome indicating the instrument that was used and the second giving the outcome that was obtained from the instrument that was used. Thus, in this case we define the new instrument $\tilde{\mathcal{I}}$ to have an outcome set $\{1, \dots, n\} \times \Omega$ so that

$$\tilde{\mathcal{I}}_{(i,x)} = p_i \mathcal{I}_x^{(i)}$$

for all $i \in \{1, \dots, n\}$ and $x \in \Omega$. We note that the traditional mixture \mathcal{I} where we do not keep track of the measured instrument can be obtained as a postprocessing of the instrument $\tilde{\mathcal{I}}$, namely $\mathcal{I}_x = \sum_i \tilde{\mathcal{I}}_{(i,x)}$ for all $x \in \Omega$.

Because of the convex structure of $\text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ we can consider the extreme points of the set.

Definition 7. An instrument $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ is *extreme* if a convex sum decomposition $\mathcal{I} = \lambda \mathcal{J} + (1 - \lambda) \mathcal{K}$ with some other instruments $\mathcal{J}, \mathcal{K} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ and some $\lambda \in (0, 1)$ implies that $\mathcal{I} = \mathcal{J} = \mathcal{K}$.

The extreme instruments were characterized in Ref. [34].

Proposition 11. An instrument $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ with a minimal Kraus decomposition $\mathcal{I}_x(\varrho) = \sum_i K_{ix} \varrho K_{ix}^*$ for all $x \in \Omega$ is extreme if and only if the set $\{K_{ix}^* K_{jx}\}_{i,j,x}$ is linearly independent.

B. Simulation scheme

Similarly to the measurement devices, in the case of quantum instruments, we consider the simulation scheme to be the following: Let \mathfrak{J} be a collection of quantum instruments with outcome set Ω from $\mathcal{L}(\mathcal{H})$ to some other output spaces which can be different for different instruments. For any finite subset $\{\mathcal{J}^{(i)}\}_{i=1}^n \subseteq \mathfrak{J}$ we choose an instrument $\mathcal{J}^{(i)} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K}_i)$ with probability p_i according to some probability distribution $(p_i)_{i=1}^n$, measure with it, and after obtaining an outcome (i, x) by keeping track of the instrument that we used, we send the output state to another instrument $\mathcal{R}^{(i,x)} \in \text{Ins}(\Lambda, \mathcal{K}_i, \mathcal{V})$ according to the classical output x of $\mathcal{J}^{(i)}$. Thus, we obtain a new instrument $\mathcal{I} \in \text{Ins}(\Lambda, \mathcal{H}, \mathcal{V})$ defined by

$$\mathcal{I}_y(\varrho) = \sum_{i=1}^n p_i \sum_{x \in \Omega} \mathcal{R}_y^{(i,x)}(\mathcal{J}_x^{(i)}(\varrho))$$

TABLE I. Summary of the analogous concepts for POVMs, channels, and instruments. For details see the main text.

	<i>POVMs</i>	<i>Channels</i>	<i>Instruments</i>
<i>Maximal elements</i>	Rank-1 POVMs	Random orthogonal	Random orthogonal
<i>Greatest element</i>	None	isometric channels	isometric instruments
<i>Least element</i>	Trivial POVMs	Trash-and-prepare	Trash-and-prepare
		channels	instruments
<i>Indecomposable elements</i>	Rank-1 POVMs	Kraus rank-1	Kraus rank-1
		channels	instruments
<i>Extreme simulation irreducible elements</i>	Extreme rank-1 POVMs	Isometric channels	Isometric channels

for all $y \in \Lambda$. The set of all instruments obtained from \mathfrak{J} by this method with some finite subset of \mathfrak{J} , some probability distribution $(p_i)_i$ and some postprocessing instruments $\mathcal{R}^{(i,x)}$ is denoted by $\text{sim}(\mathfrak{J})$.

In Ref. [20] it was shown that in the case of POVMs (and more generally measurements in general probabilistic theories) there is a collection of POVMs that can be used to simulate all other POVMs. Thus, just as in the case of measurements, we can try to reduce the problem of simulability into a specific class of instruments.

Definition 8. An instrument \mathcal{I} is simulation irreducible if for any set of instruments \mathfrak{J} such that $\mathcal{I} \in \text{sim}(\mathfrak{J})$ there exists an instrument $\mathcal{J} \in \mathfrak{J}$ such that $\mathcal{I} \leftrightarrow \mathcal{J}$.

It is straightforward to verify that the characterization of simulation irreducibility follows the same proof as in Ref. [20].

Proposition 12. An instrument is simulation irreducible if and only if it is postprocessing clean and postprocessing equivalent to an extreme instrument.

Since the identity channel is extremal and postprocessing clean instruments are exactly those that are equivalent with the identity channel, the set of simulation irreducible instruments coincides with the set of postprocessing clean instruments. Thus, we get the following for free from Proposition 2.

Proposition 13. Every instrument can be simulated by any instrument that is equivalent with the identity channel.

Furthermore, it is easy to see that the extreme elements of the equivalence class of the identity channel are just the isometric channels. Namely, if $\mathcal{I} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ is equivalent with the identity channel, then by Proposition 3 we have that $\mathcal{I}_x(\varrho) = \sum_{i=1}^{n_x} p_{xi} V_{xi} \varrho V_{xi}^*$ for all $x \in \Omega$ and $\varrho \in \mathcal{S}(\mathcal{H})$ for some probability distribution $(p_{xi})_{x \in \Omega, i \in \{1, \dots, n_x\}}$ and some isometries $V_{xi} : \mathcal{H} \rightarrow \mathcal{K}$ such that $V_{xj}^* V_{xi} = 0$ for all $i \neq j$ for all $x \in \Omega$. We notice that if we define instruments $\mathcal{V}^{(i,y)} \in \text{Ins}(\Omega, \mathcal{H}, \mathcal{K})$ by setting $\mathcal{V}_x^{(i,y)}(\varrho) = \delta_{xy} V_{yi} \varrho V_{yi}^*$, we see that then $\mathcal{I} = \sum_{y \in \Omega} \sum_{i=1}^{n_y} p_{yi} \mathcal{V}^{(i,y)}$. Thus, \mathcal{I} is extreme if and only if the probability distribution is trivial, i.e., \mathcal{I} has only one outcome so that it is an isometric channel, i.e., $\mathcal{I}(\varrho) = V \varrho V^*$ for all $\varrho \in \mathcal{S}(\mathcal{H})$ for some isometry $V : \mathcal{H} \rightarrow \mathcal{K}$.

VII. SUMMARY

Motivated by the postprocessing (and simulability) of POVMs, the first aim of this paper is to define mathematically the postprocessing of quantum instruments and characterize the partial order that it induces on the set of equivalence

classes of instruments. In particular, we characterize the least and greatest elements and characterize their equivalence classes (see Table I for illustration).

We see that, similarly to the case of quantum channels in Ref. [15], every quantum instrument can be postprocessed to a so-called trash-and-prepare instrument (the least element) that simply ignores the input state and prepares a new state as output. In accordance with this similarity, we also saw that every instrument can be postprocessed from instruments that are equivalent with the identity channel (the greatest element). We find that those instruments consist of randomly chosen isometries with orthogonal output ranges.

Furthermore, we consider two other important classes of instruments. First, the indecomposable instruments have the mathematical advantage that each operation of the instrument consists of only one Kraus operator (for example, Lüders instruments), implying that any instrument can be postprocessed from some indecomposable instruments. For the converse we show that any instrument with Kraus operators that have orthogonal output ranges for each outcome (however many Kraus operators it may have) can be postprocessed to an indecomposable (detailed) instrument.

Our conjecture is that this condition is also necessary for an instrument to be equivalent with an indecomposable (detailed) instrument. However, we have not found a general proof for this claim and we leave this as an open question for further work.

The second class of instruments we study are the measure-and-prepare instruments, which measure some observable on the input and based on the outcome they prepare a new state. We show how an instrument that can be postprocessed from a measure-and-prepare instrument must look like and we show that only measure-and-prepare instruments that have indecomposable (rank 1) induced POVMs are postprocessing equivalent with an indecomposable (detailed) instrument. As a corollary we see that actually any instrument with indecomposable induced POVM falls into this class.

For these two classes of instruments, we draw some connections between postprocessing of quantum instruments and postprocessing of their induced POVMs. In particular, we see that two indecomposable instruments are equivalent if and only if their induced POVMs are equivalent, and that a measure-and-prepare instrument can be postprocessed from another measure-and-prepare instrument if and only if the induced POVM of the former instrument can be postprocessed from the induced POVM of the latter instrument.

Finally, we use the partial order introduced above to study simulability of instruments. We see that every instrument can be simulated by an instrument from the equivalence class of the identity instrument. We find that the extreme simulation irreducible instruments are the isometric channels.

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