

An introduction to

Local Hamiltonians & Quantum Complexity

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S. Mozes, P. Wocjan, O. Regev, P. Love, P. Shor, S. Jordan,
S. Lloyd, A. Landahl, S. Irani, D. Gottesman, S. Bravyi,
E. Farhi, J. Goldstone, D. Aharonov, T. Morimae,
R. Movassagh, F. Brandao, J. Eisert, J. Whitfield, ...

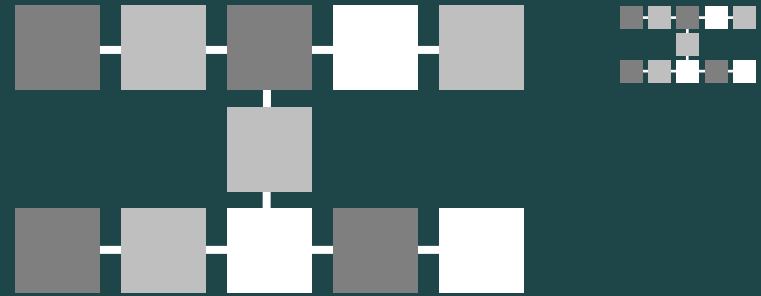
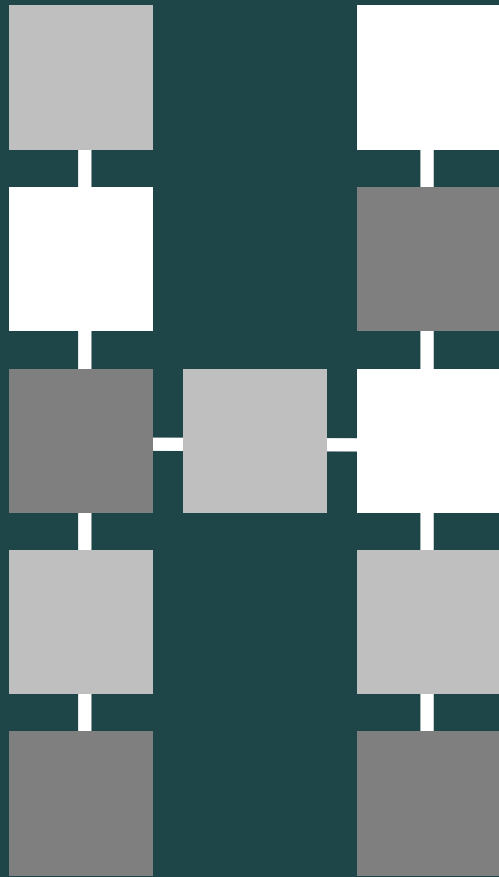
Daniel Nagaj



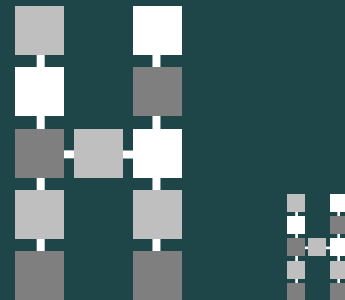


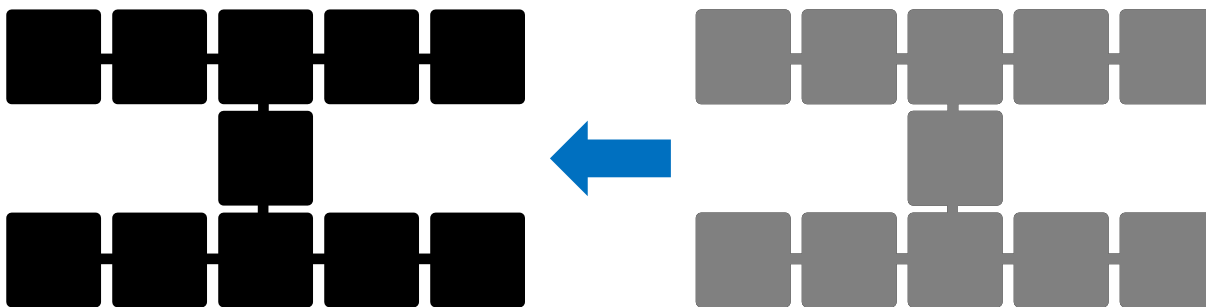
Local Hamiltonian





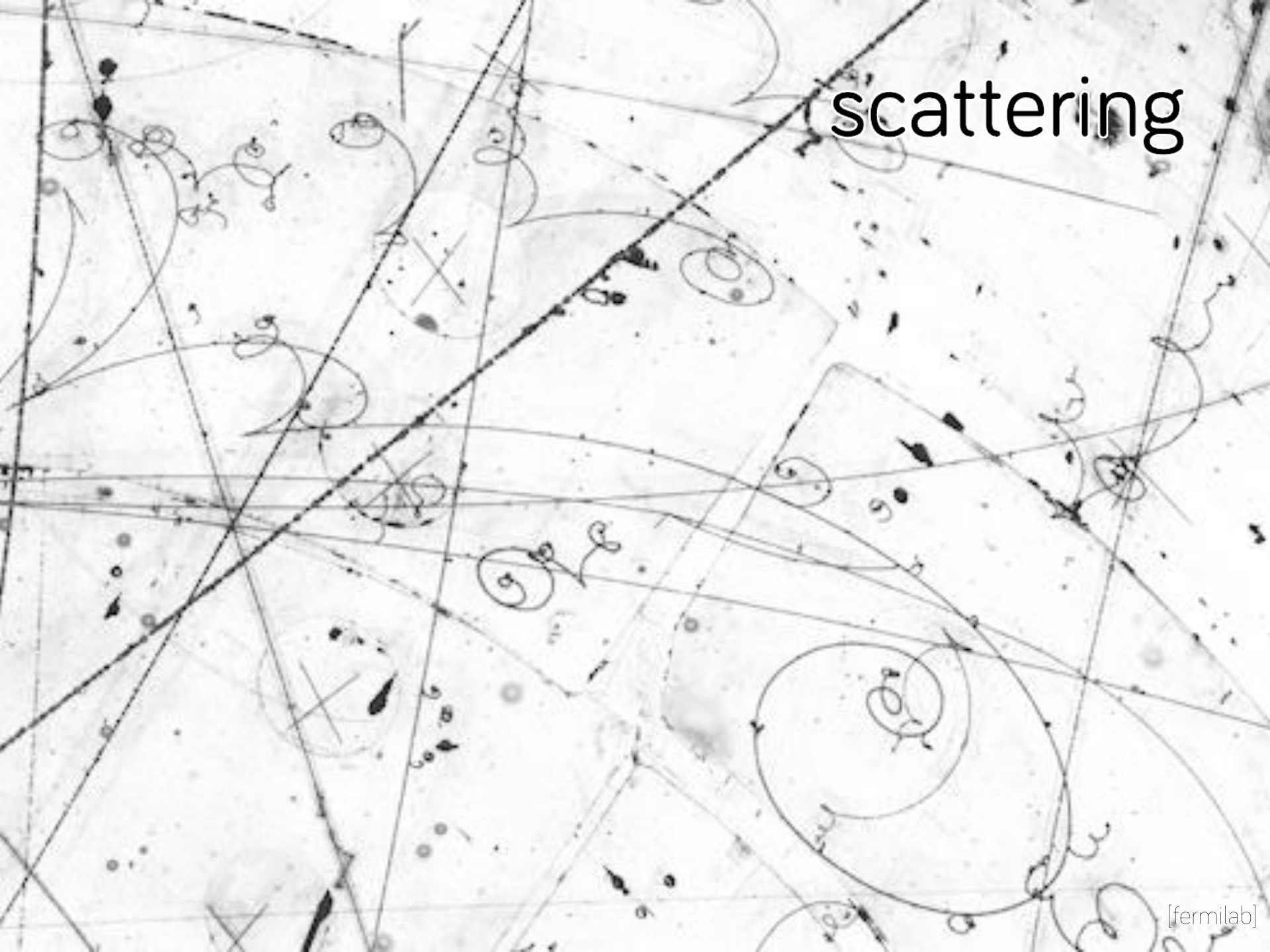
Hamiltonians





dynamics
 $i\hbar \frac{\partial}{\partial t} |\psi\rangle = \hat{H} |\psi\rangle$

scattering



scattering

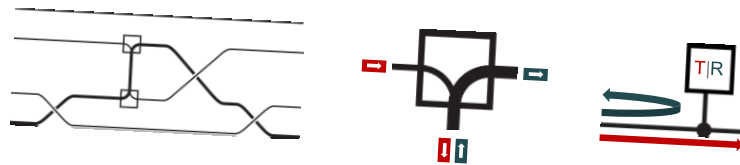
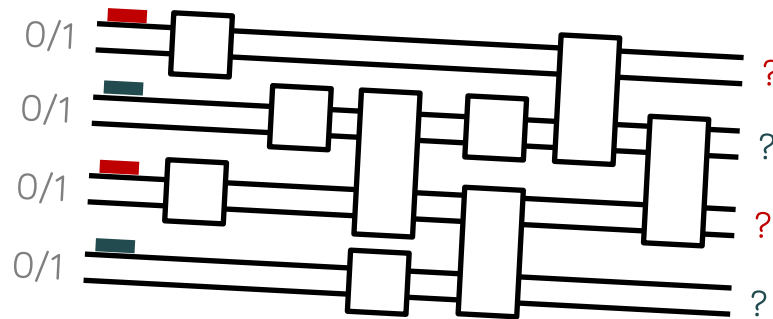
Universal computation by multi-particle quantum walk

- dual-rail encoding
N wavepackets

$$a_j^\dagger a_k + a_k^\dagger a_j$$

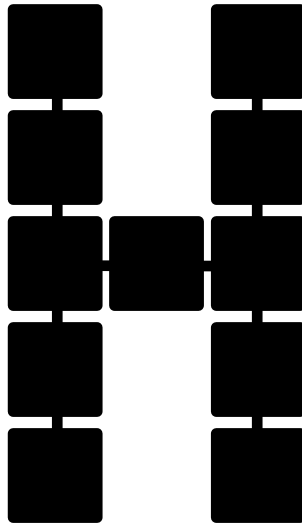
- CPHASE: interaction

$$a_j^\dagger a_k^\dagger a_j a_k$$

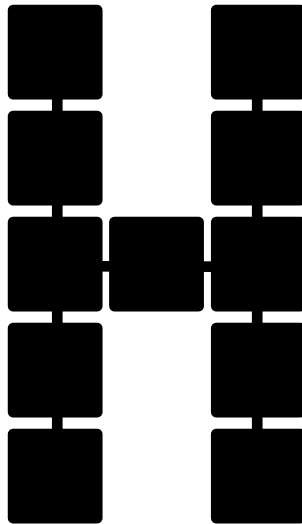


[Childs, Gosset, Webb, Science 339, 791 (2013)]

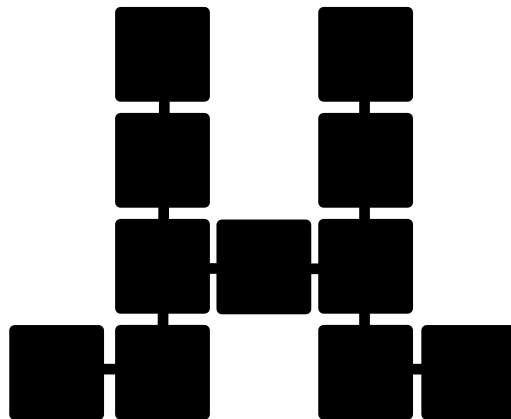
optimization



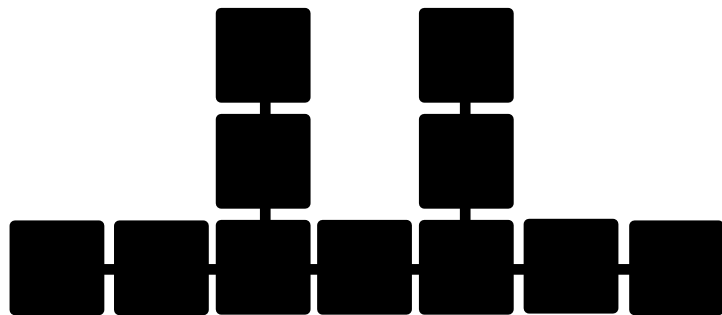
optimization



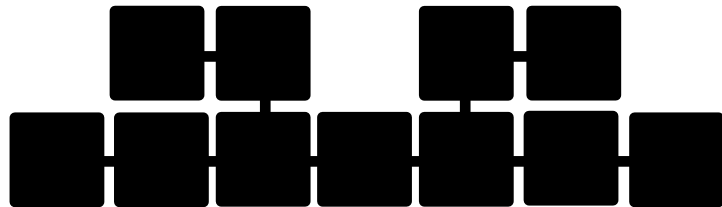
optimization



optimization

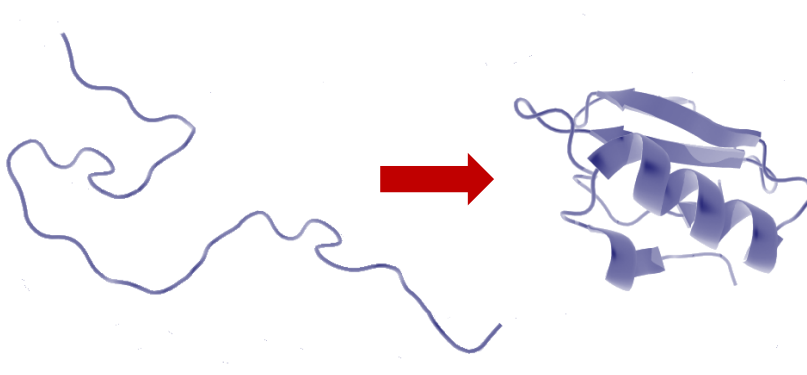


optimization

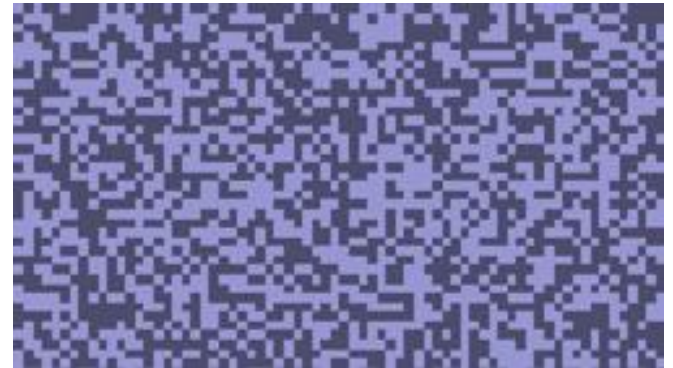


protein folding

spin glasses



[wikipedia]



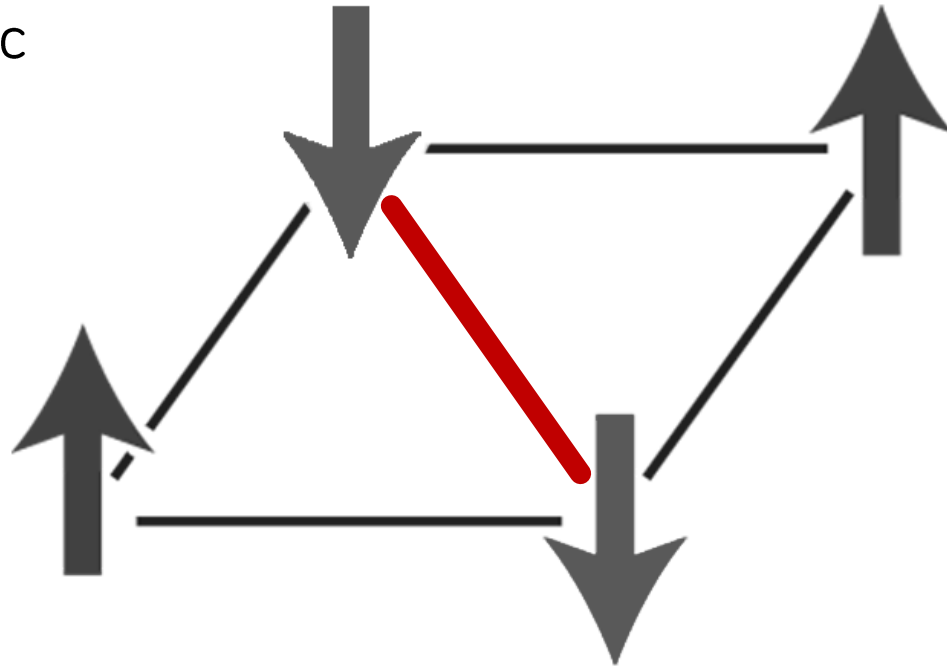
[uni-koeln.de]

local Hamiltonians

0 Frustrated systems

You can't make everybody happy.

antiferromagnetic
spin glass



a global
ground state

HARD?

find & describe it?
is it entangled?

frustrated

FRUST
RATED

0 Local Hamiltonians

- What are they like?

ground state (energy)

QMA-complete problems




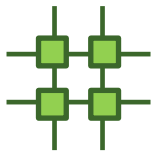


$$H(t) = \sum_j H_j(t)$$

- What are they good for?

quantum computing,
chemistry, control, transport

BQP universality

Many local questions.

- local particle dimension 
- interaction geometry 
- time independence 
- translational invariance 
- promise and eigenvalue gaps 
- energy \times time cost 

how hard
is this
question



We Can Do It!



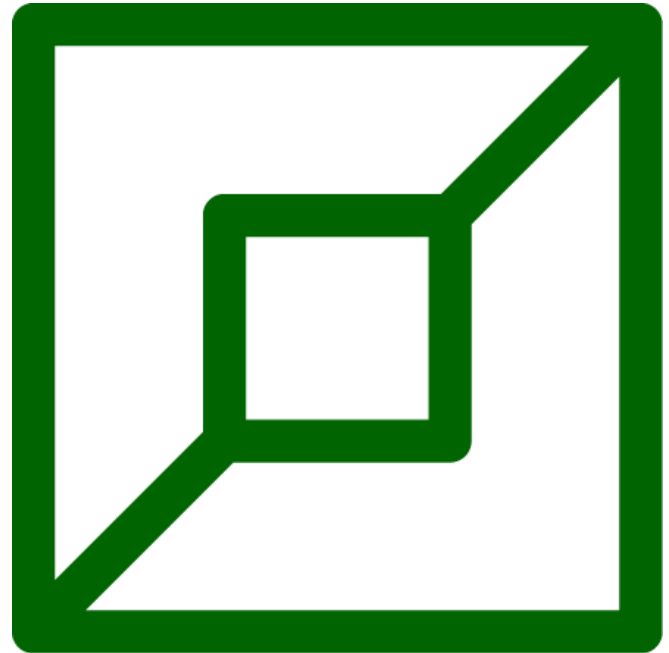
P

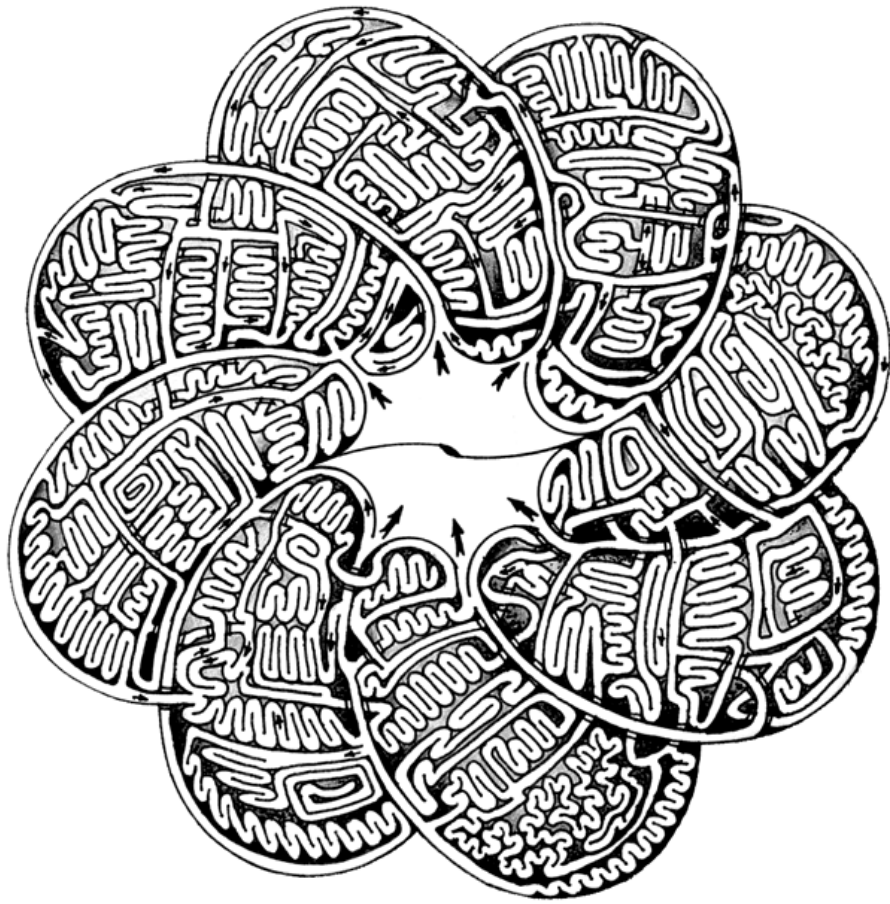
[J. Howard Miller]

POST FEB. 15 TO FEB. 20



WAR PRODUCTION CO-ORDINATING COMMITTEE



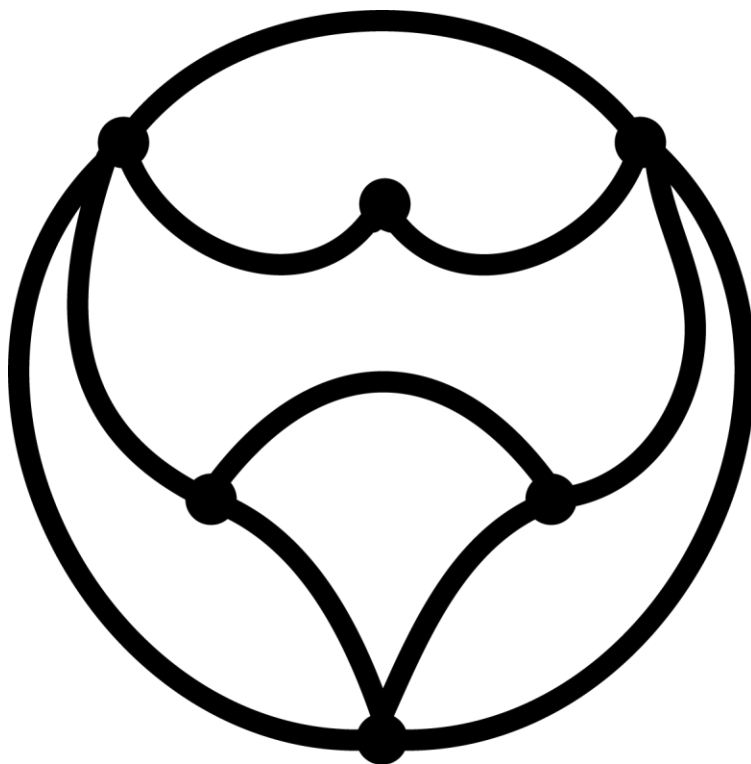
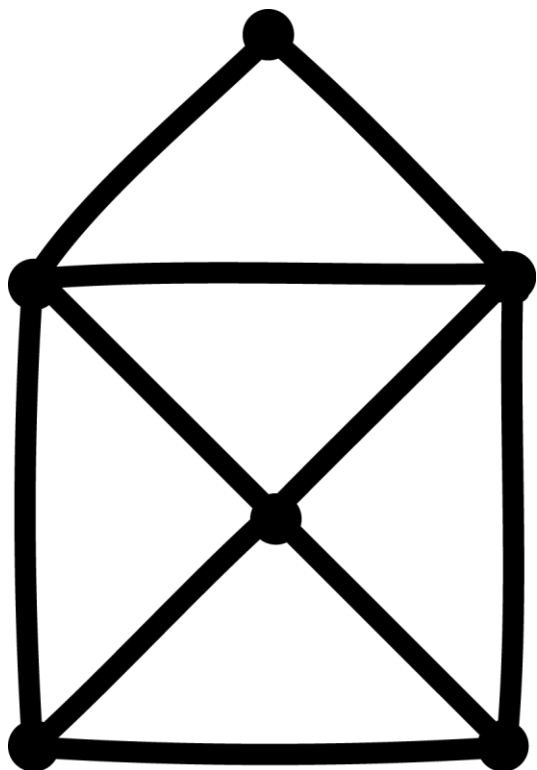


NP

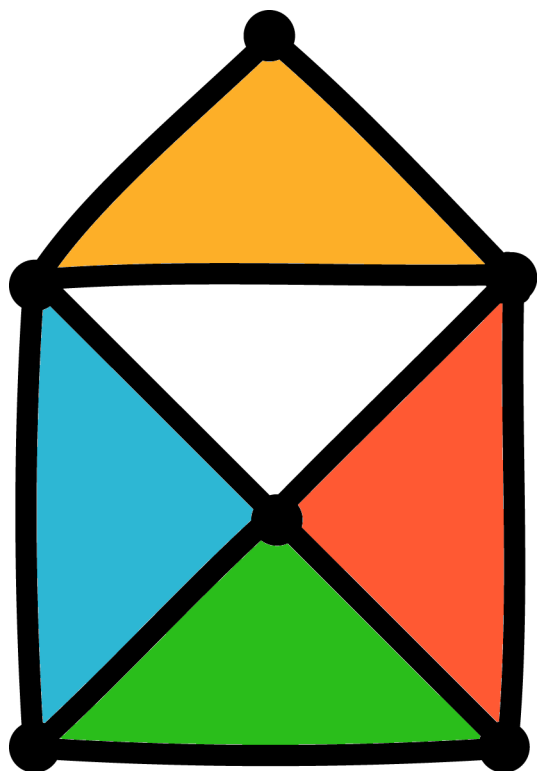


[maze: Andrew Bernhardt]
[A+M: primaryresources.co.uk]

1 A graph isomorphism puzzle



1 A graph isomorphism puzzle



1 A cryptarithmic puzzle

$$\begin{array}{r} \text{DID} \\ + \text{DINOS} \\ \hline \text{CROAK} \end{array}$$



1 A cryptarithmic puzzle

$$\begin{array}{r} 595 \\ + 59842 \\ \hline 60437 \end{array}$$



1 The NP protocol

Did dinosaurs exist?



a proof

1 The NP protocol

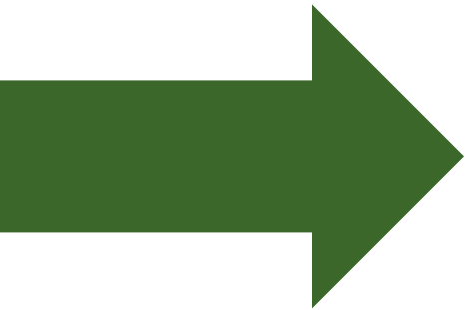
Did dinosaurs exist?



a witness

1 The class NP

Yes/no questions, easy to verify solutions.



a verification
circuit

from a uniform family

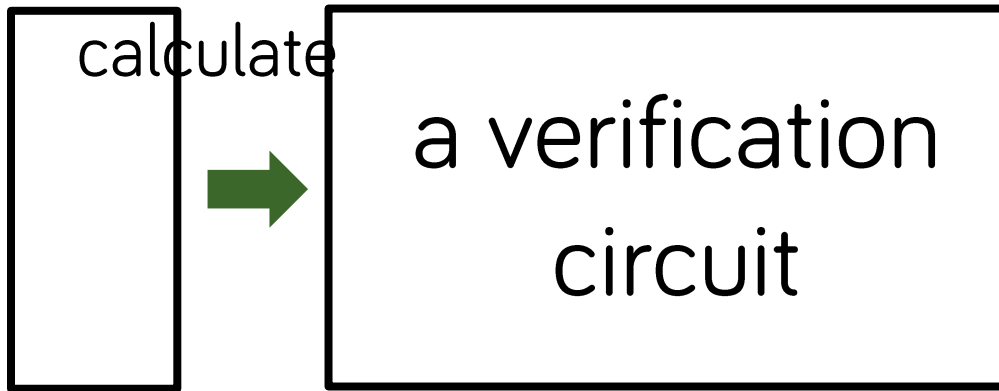


YES? Accept a good proof.

NO? Reject any witness.

1 The class P

Yes/no questions that we can answer.

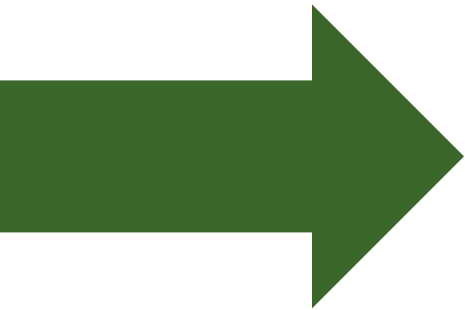


YES? Figure it out by yourself.

NO? Figure it out by yourself.

1 The class NP

Yes/no questions, easy to verify solutions.



a verification
circuit

from a uniform family



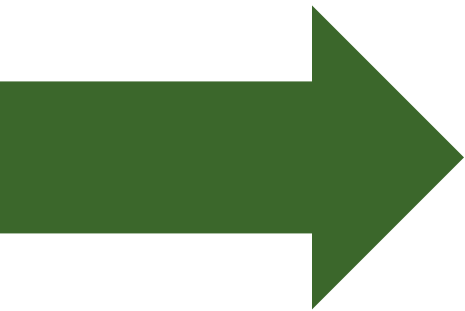
YES? Accept a good proof.

NO? Reject any witness.

1 An NP-hard problem

The mother of them all.

- An NP-hard problem solver solves *anything* in NP.

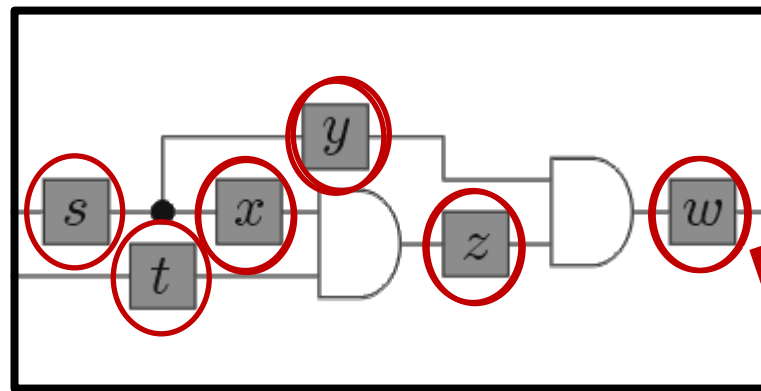


a verification
circuit

1 An NP-hard problem

The mother of them all.

- An NP-hard problem solver solves *anything* in NP.



Could this circuit ever output 1?

3-local conditions

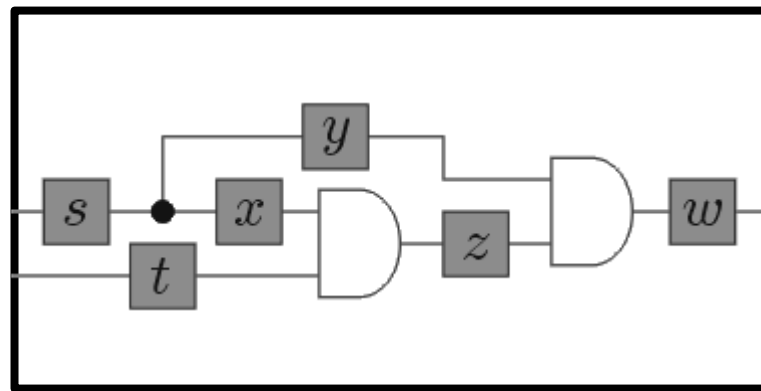
$$(\dots \vee \dots \vee \dots) \wedge (\dots \vee \dots \vee \dots) \wedge \dots$$

- 3-SAT is NP-hard.
- 3-SAT is in NP.

1 An NP-hard problem

Satisfiability.

- An NP-hard problem solver solves anything in NP.

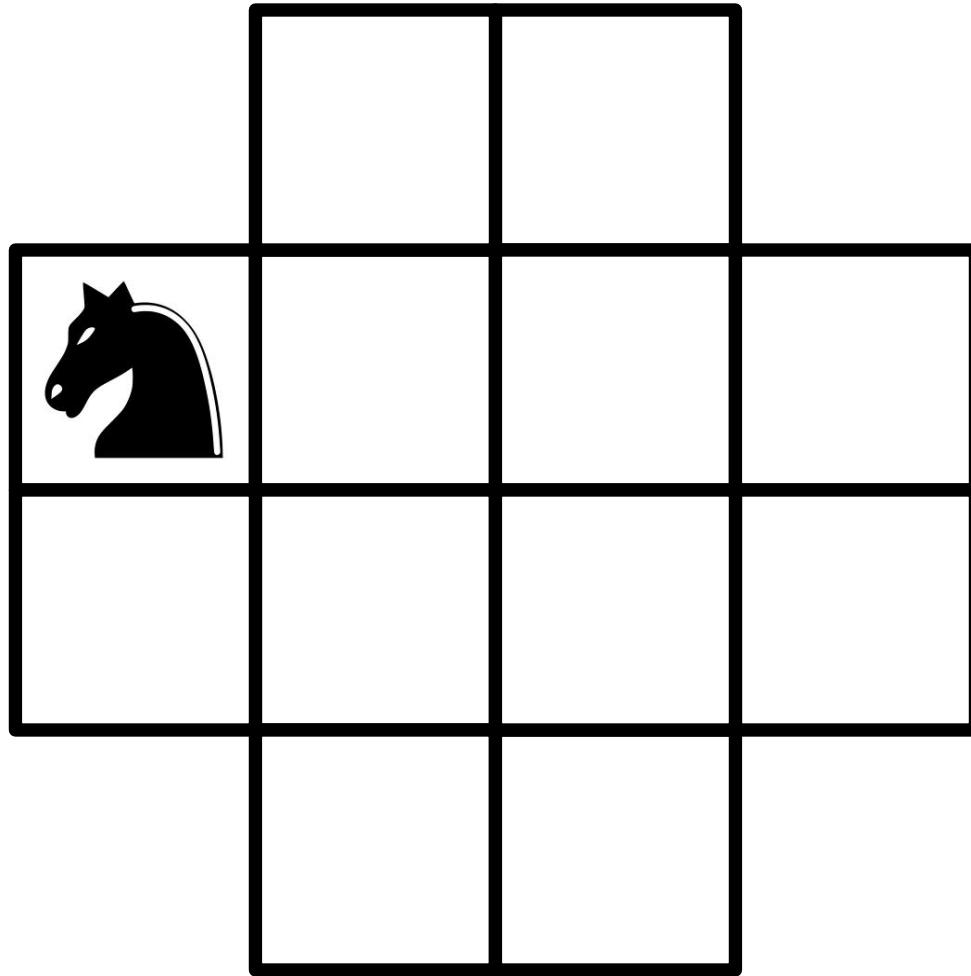


Could this circuit ever output 1?

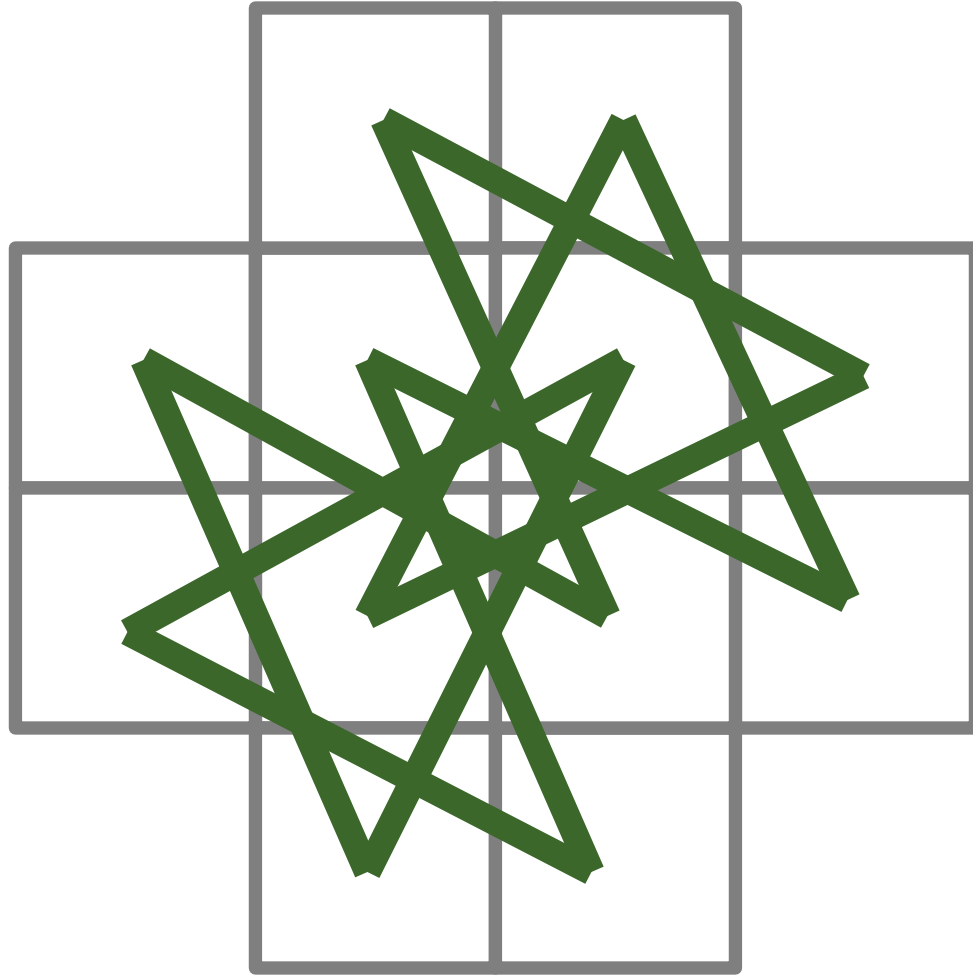
3-local conditions

- 3-SAT is NP-complete. [Cook, Levin]

1 NP-complete problems: Hamiltonian cycle

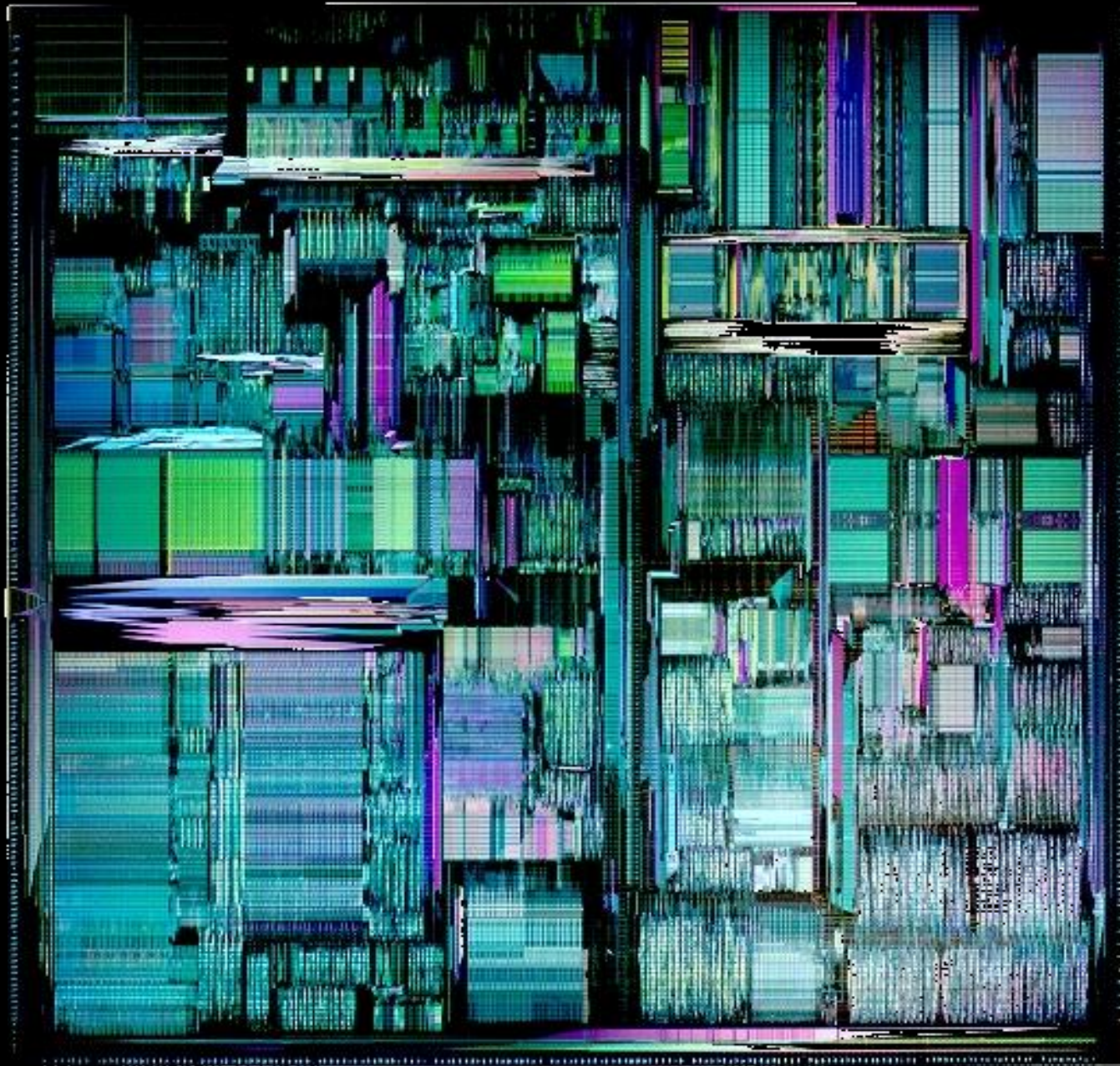


1 NP-complete problems: Hamiltonian cycle





the puzzles of QMA



[1995 Pentium Pro

DAVE

[wire

P



BOUNDED ERROR
QUANTUM
POLYNOMIAL TIME

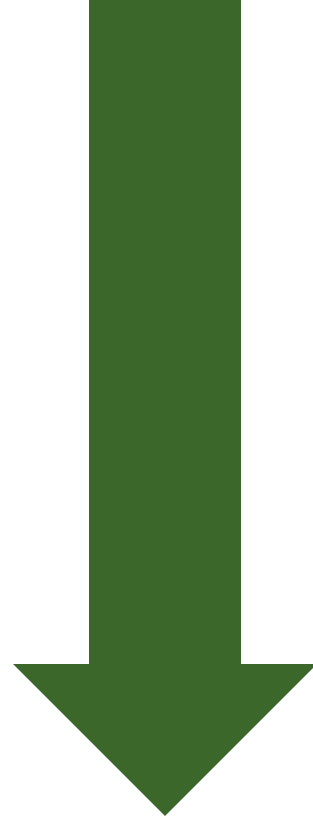
[tha

NP



2 The MA protocol

Did dinosaurs exist?



2 The MA protocol

Did dinosaurs exist?



2 The MA protocol

Did dinosaurs exist?

YES?
Eager to be
convinced.



[magnifying glass: hllllllal]

2 The MA protocol

Recognizing fakes?



2 The MA protocol

Recognizing fakes?

NO?
Don't be
fooled
easily.



2 Probabilistic checks

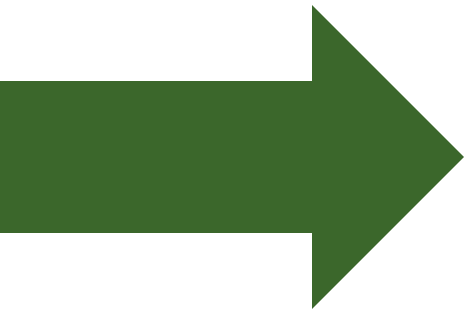
Sometimes reject
a genuine proof?

Accept
a fake?



1 The MA protocol

Probabilistic checks.



probabilistic
verification

from a uniform family



YES? Accept a good proof with $p > a$.

NO? Probability of accepting $p < b$.



2 The QMA protocol

Quantum checks.



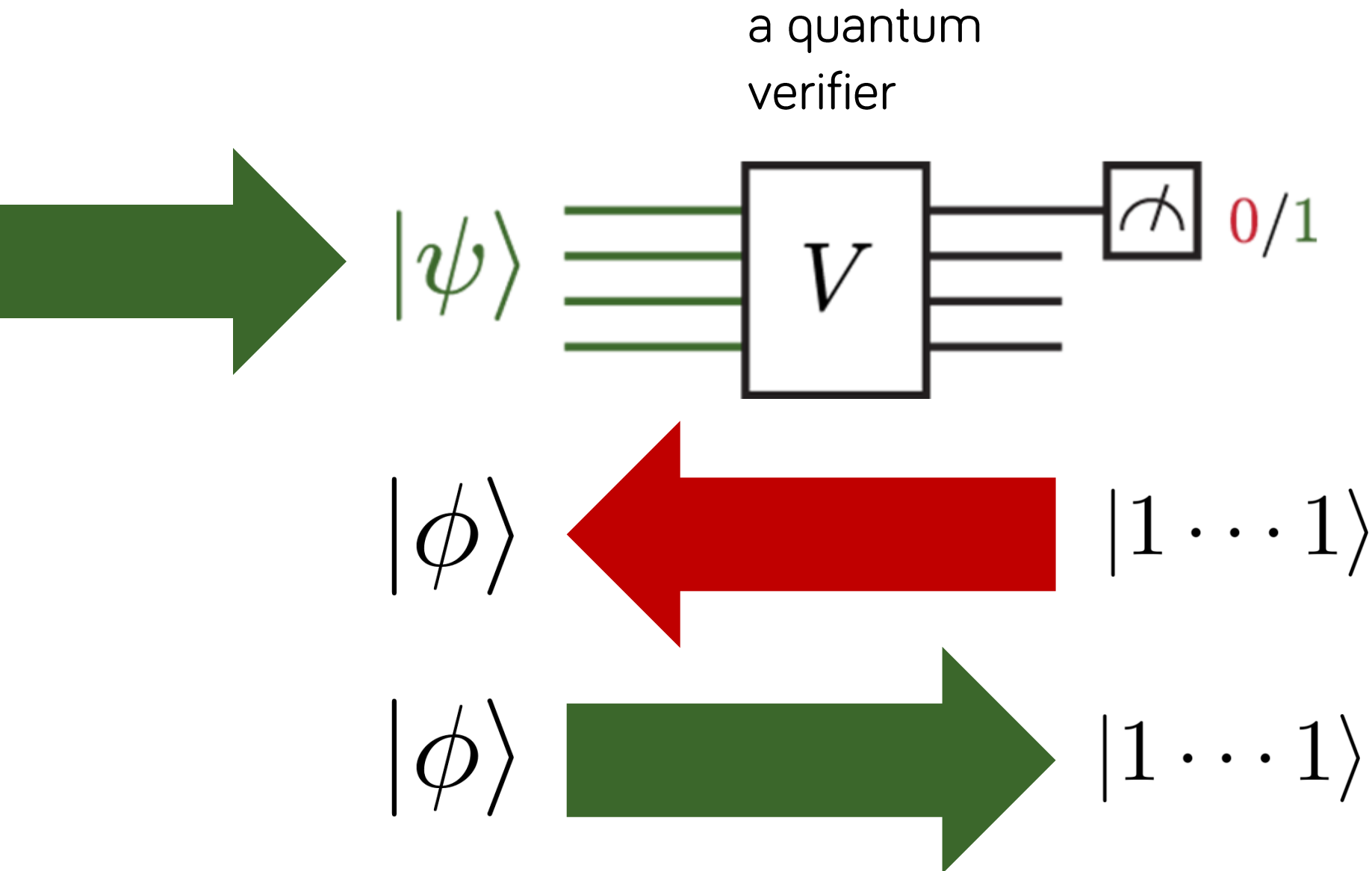
YES? Accept a good proof with $p > a$.

NO? Probability of accepting $p < b$.



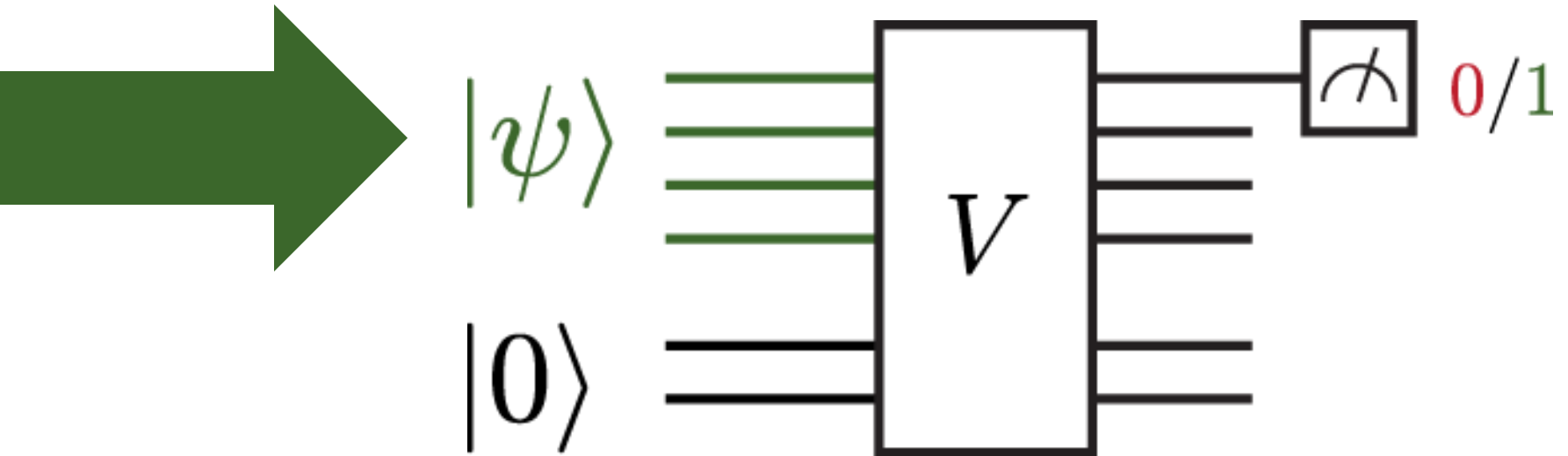
2 The QMA protocol

This is too simple.



2 The QMA protocol

Ancillas are necessary.

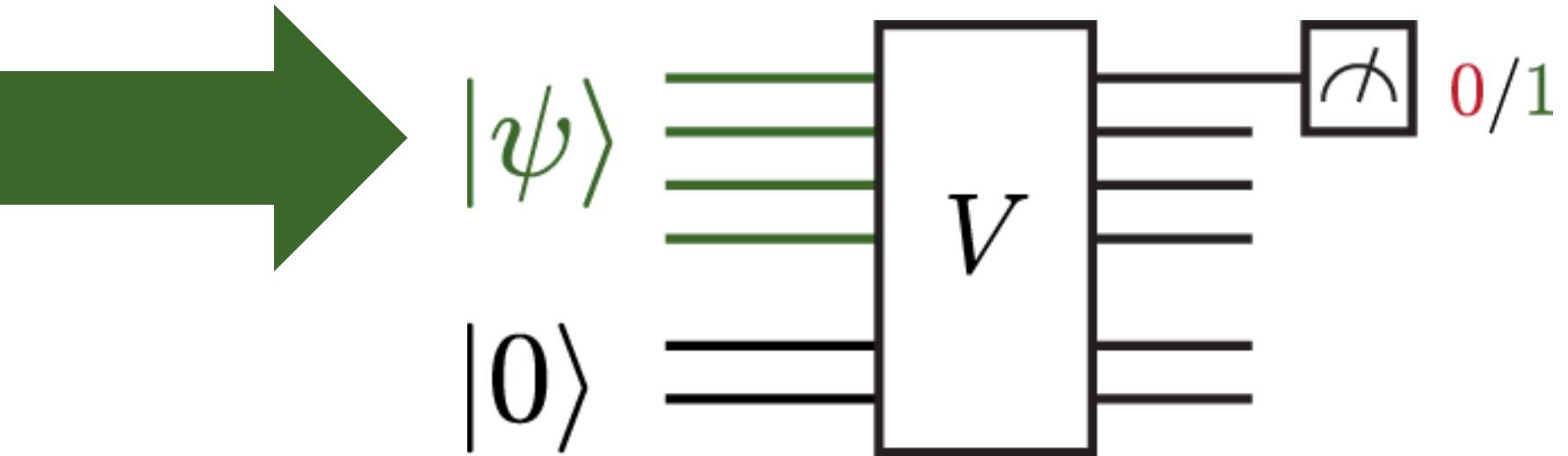


YES? Accept a good proof with $p > a$.

NO? Probability of accepting $p < b$.



2 A QMA-hard question



Could we feed this quantum verifier something that likely outputs 1?

2 Implementing reversible quantum circuits

- The Schrödinger equation

a Hamiltonian H

$$i \frac{d}{dt} |\psi(t)\rangle = H(t) |\psi(t)\rangle$$

made of local terms

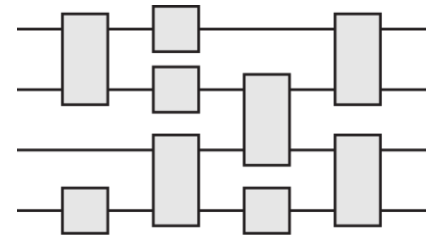
$$H(t) = \sum_j H_j(t)$$

unitary time evolution

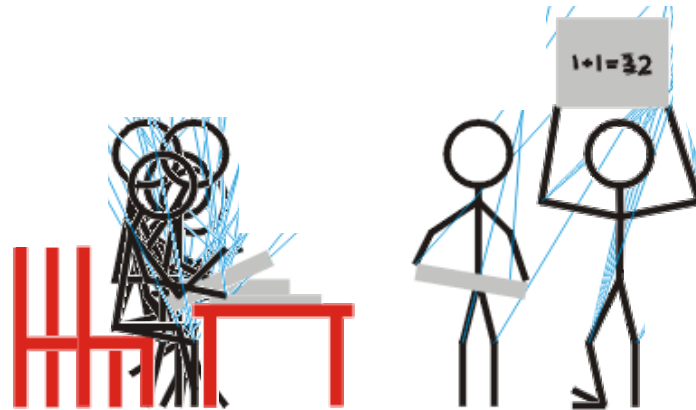
$$|\psi(t)\rangle = U_{t,0} |\psi(0)\rangle$$

- Implementing a circuit U

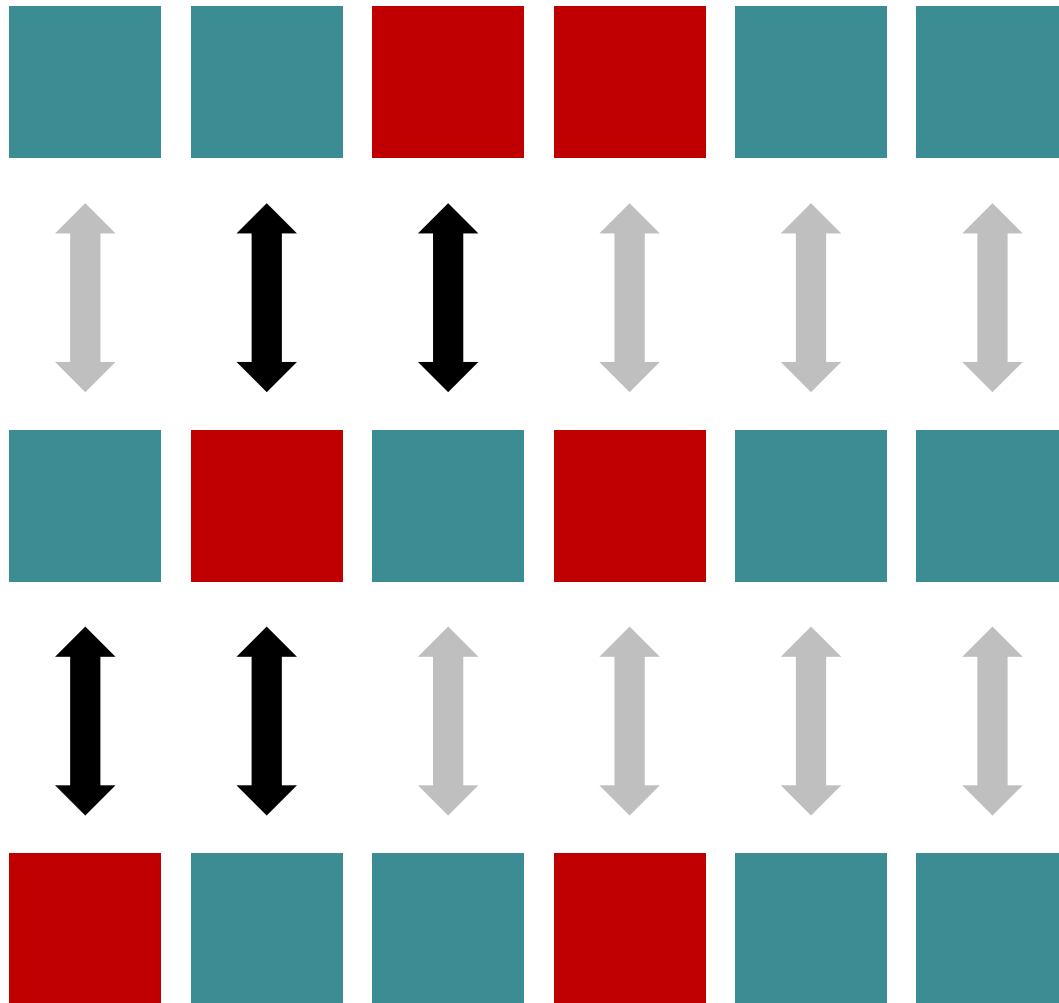
how to make a system compute?



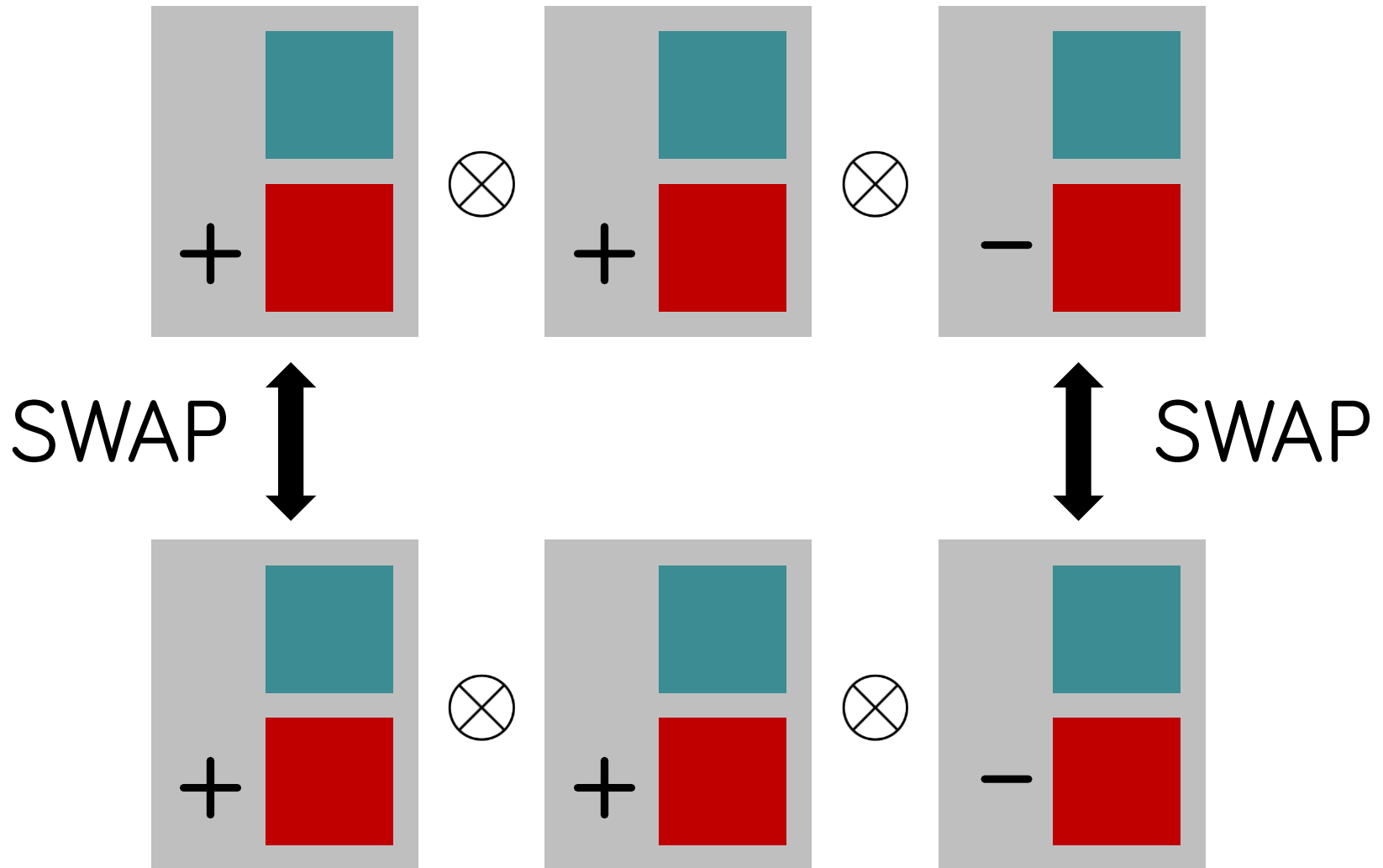
2 Snapshots of a computation



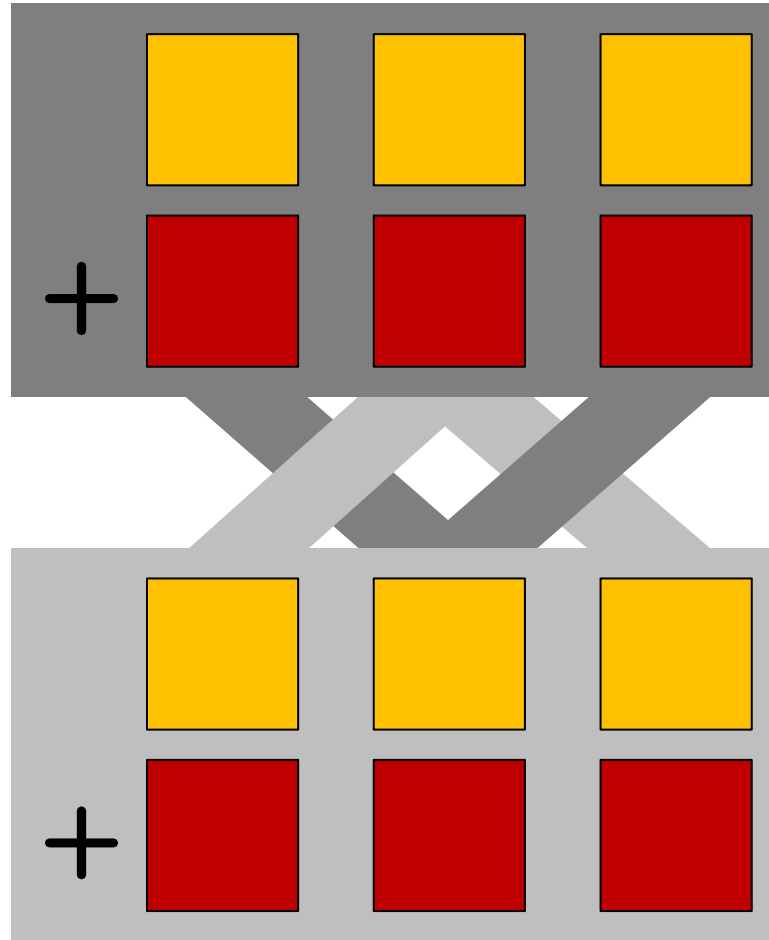
Locally comparing **strings**.



Locally comparing **product** states: SWAP.

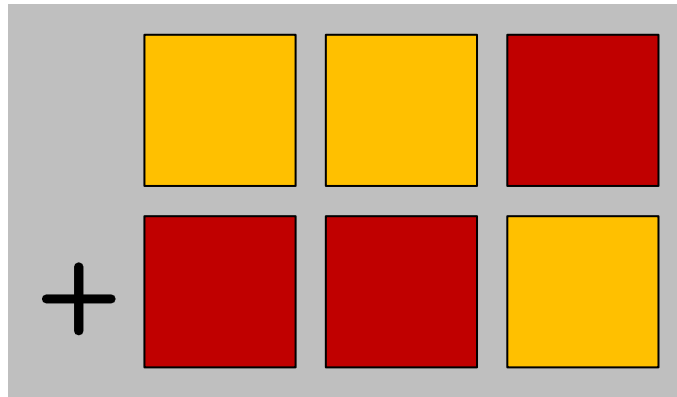


Locally comparing **entangled** states?



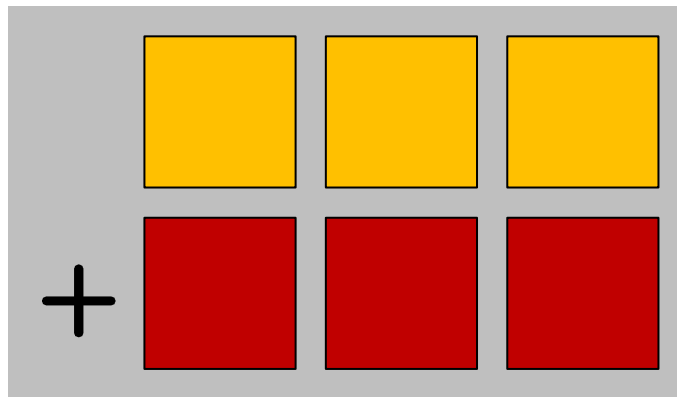
UGH!

2 The data & the clock



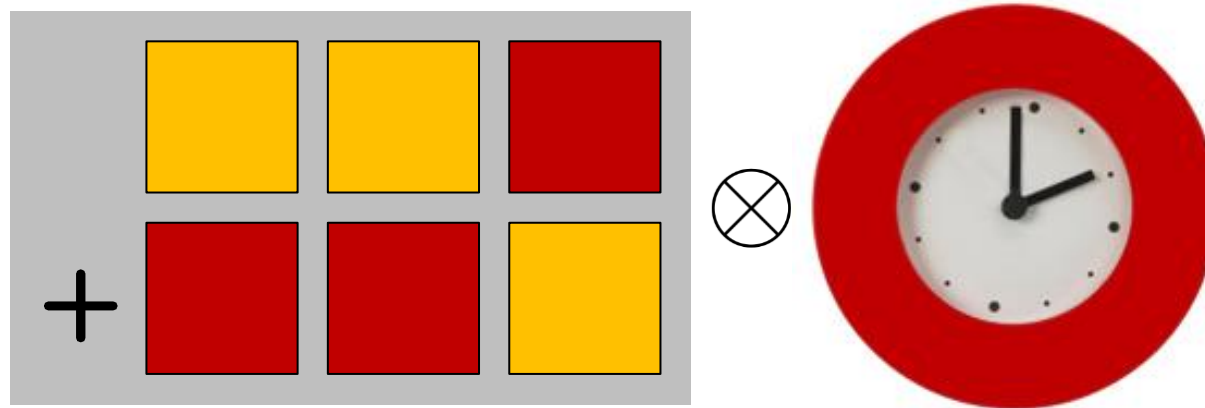
U^\dagger   U

Hard to compare directly (locally).

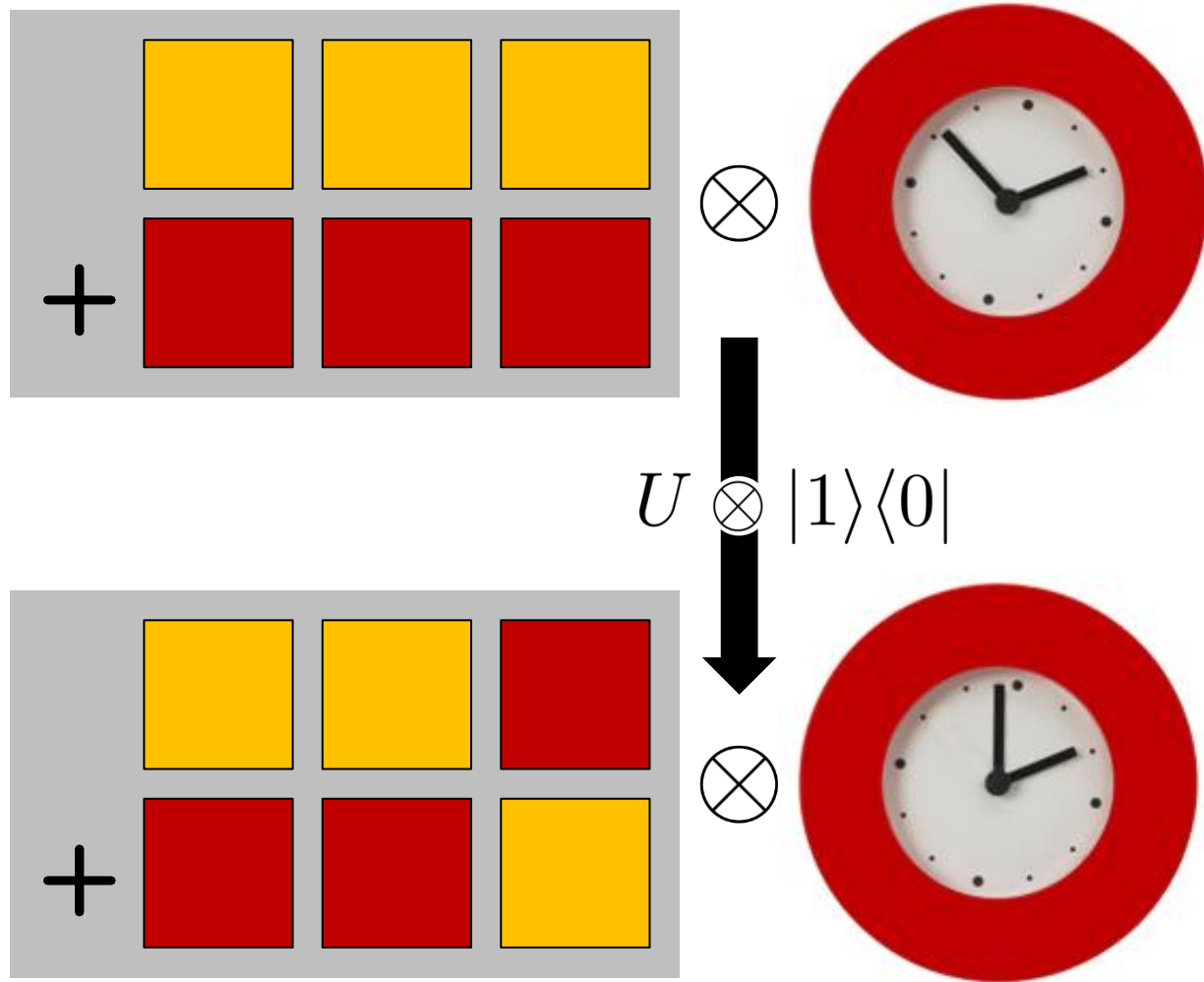


a clock

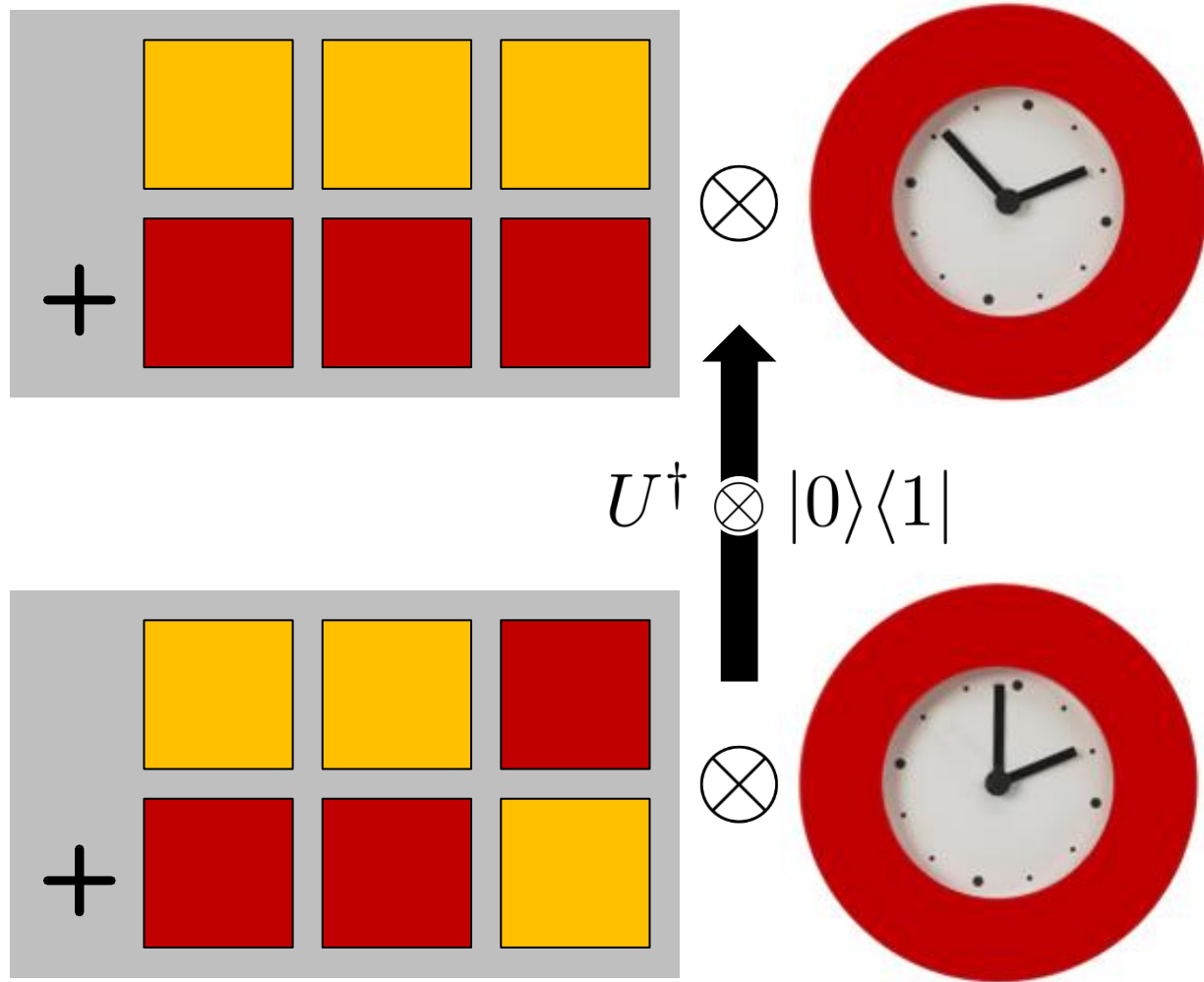
2 The data & the clock



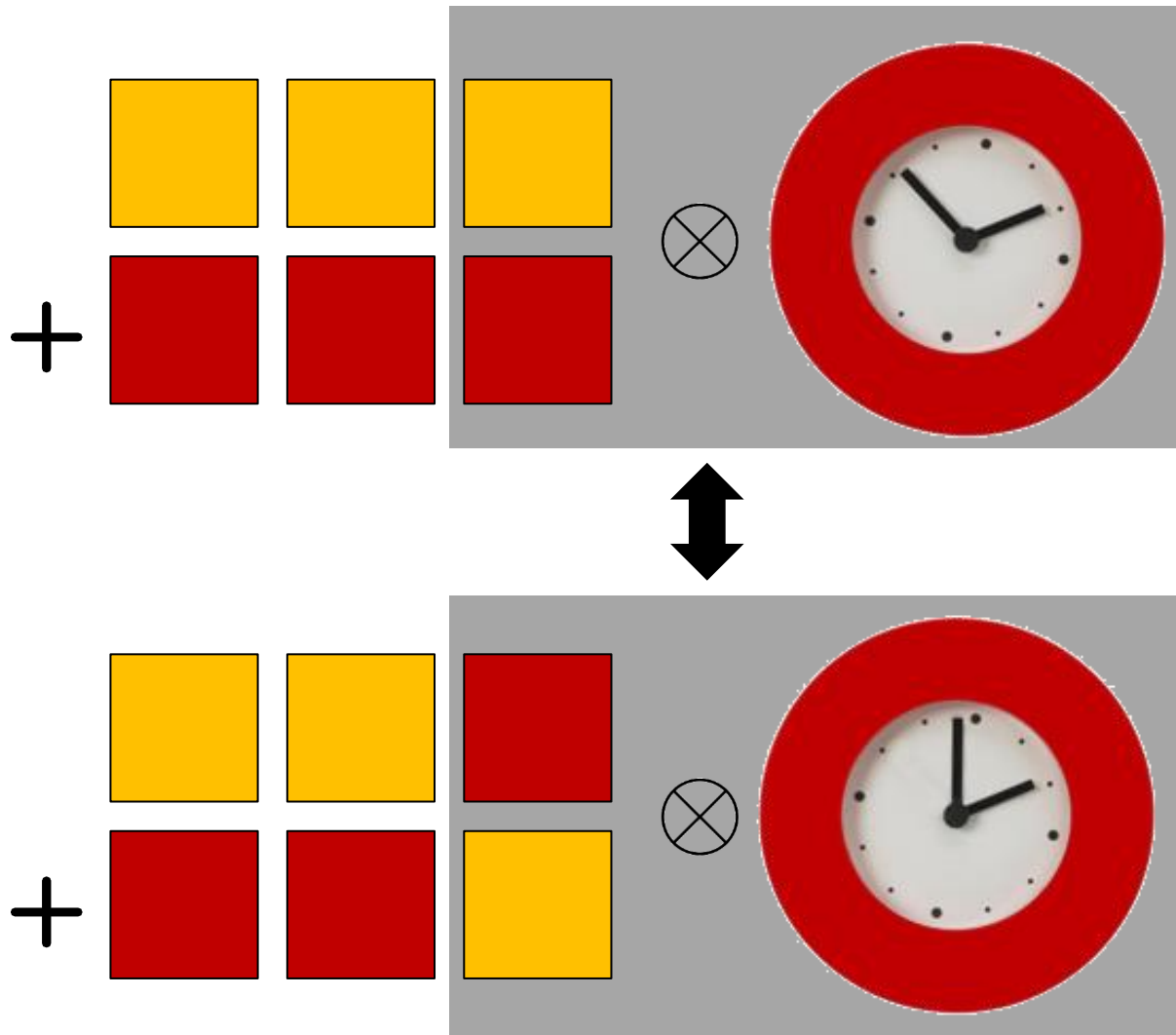
2 The data & the clock



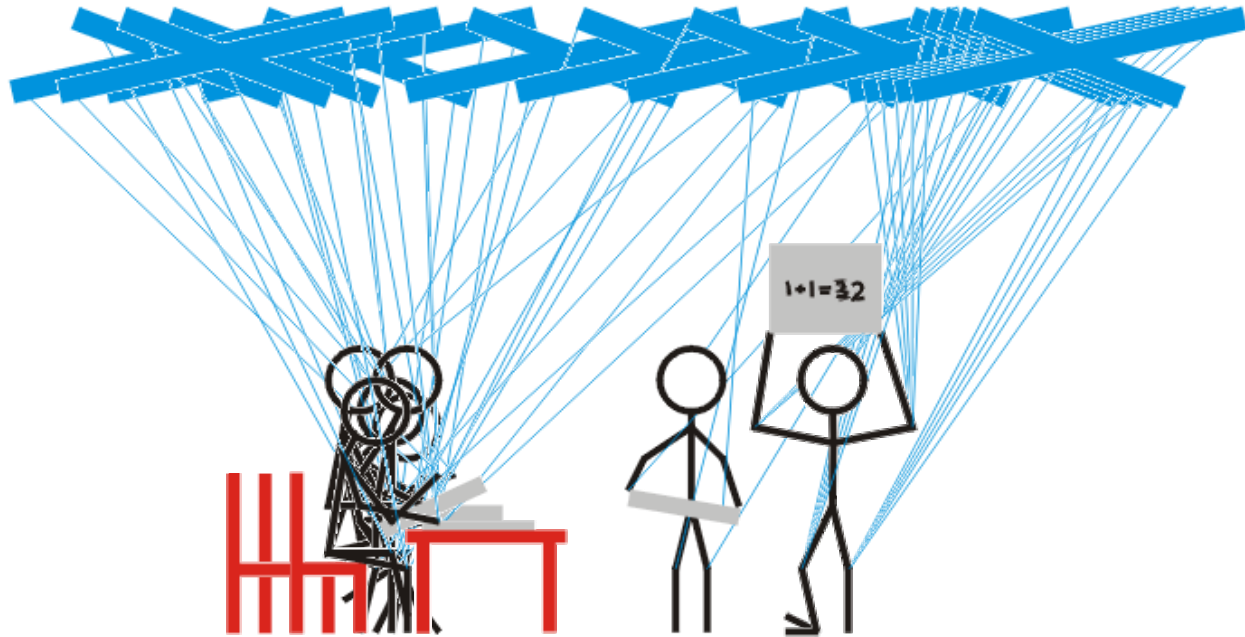
2 The data & the clock



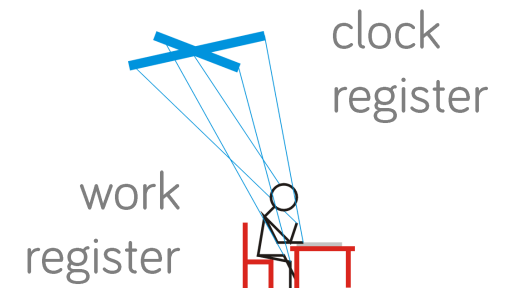
2 The data & the clock



2 Snapshots of a computation & a “clock”



2 Feynman's (Hamiltonian) computer

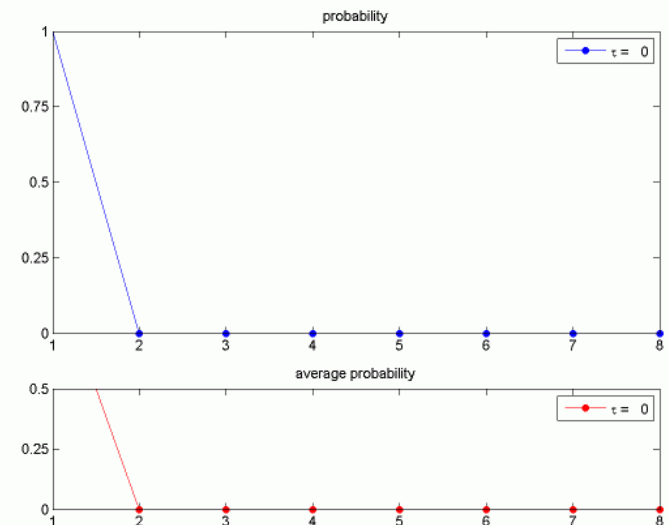


■ The Hamiltonian

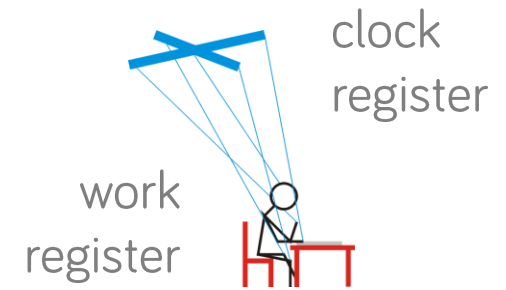
$$H_F = - \sum_{t=1}^L \left(U_t \otimes |t\rangle \langle t-1| + U_t^\dagger \otimes |t-1\rangle \langle t| \right)$$

■ A quantum walk on a “line”

$$\begin{aligned} & |\varphi_0\rangle \otimes |0\rangle_c \\ & U_1 |\varphi_0\rangle \otimes |1\rangle_c \\ & U_2 U_1 |\varphi_0\rangle \otimes |2\rangle_c \\ & U_3 U_2 U_1 |\varphi_0\rangle \otimes |3\rangle_c \end{aligned}$$



2 Feynman's (Hamiltonian) computer

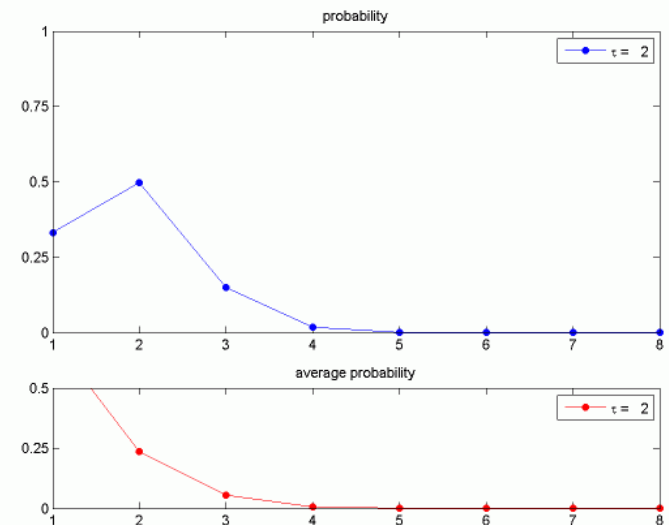


■ The Hamiltonian

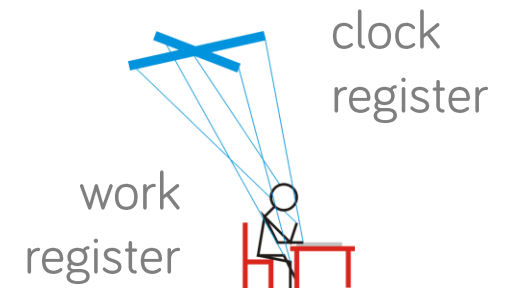
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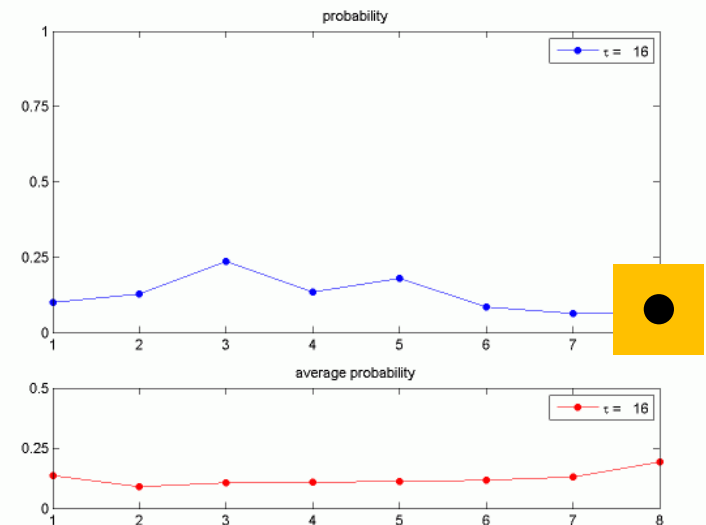
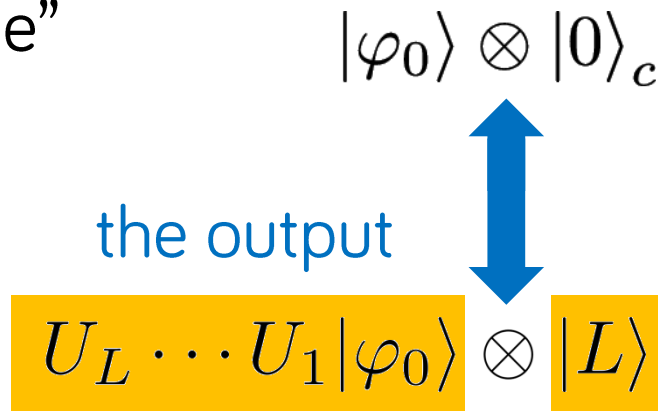
2 Feynman's (Hamiltonian) computer



■ The Hamiltonian

$$H_F = - \sum_{t=1}^L \left(U_t \otimes |t\rangle \langle t-1| + U_t^\dagger \otimes |t-1\rangle \langle t| \right)$$

■ A quantum walk on a "line"



2 Hamiltonian QC

It's a walk.

- Feynman's Hamiltonian

$$H_F = - \sum_{t=1}^L \left(U_t \otimes |t\rangle \langle t-1| + U_t^\dagger \otimes |t-1\rangle \langle t| \right)$$

- The “line” of states

$$\begin{array}{l} |\varphi_0\rangle \otimes |0\rangle_c \\ U_1 |\varphi_0\rangle \otimes |1\rangle_c \\ U_2 U_1 |\varphi_0\rangle \otimes |2\rangle_c \\ U_3 U_2 U_1 |\varphi_0\rangle \otimes |3\rangle_c \\ U_4 U_3 U_2 U_1 |\varphi_0\rangle \otimes |4\rangle_c \end{array} \quad H_F = - \begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

- The eigenvectors: combinations of plane waves

2 Hamiltonian QC

It's a walk.

- Feynman's Hamiltonian

$$H_F = - \sum_{t=1}^L \left(U_t \otimes |t\rangle \langle t-1| + U_t^\dagger \otimes |t-1\rangle \langle t| \right)$$

- The “line” of states

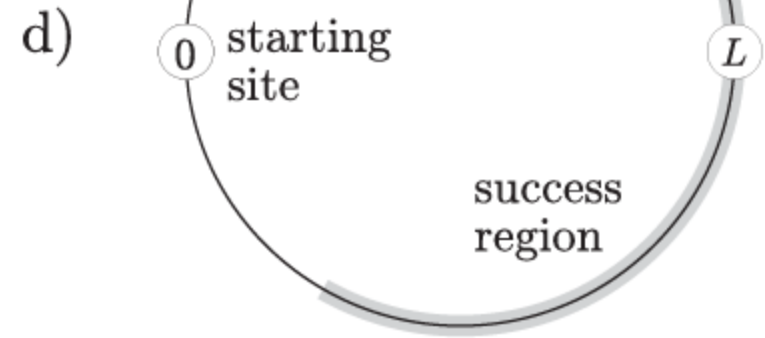
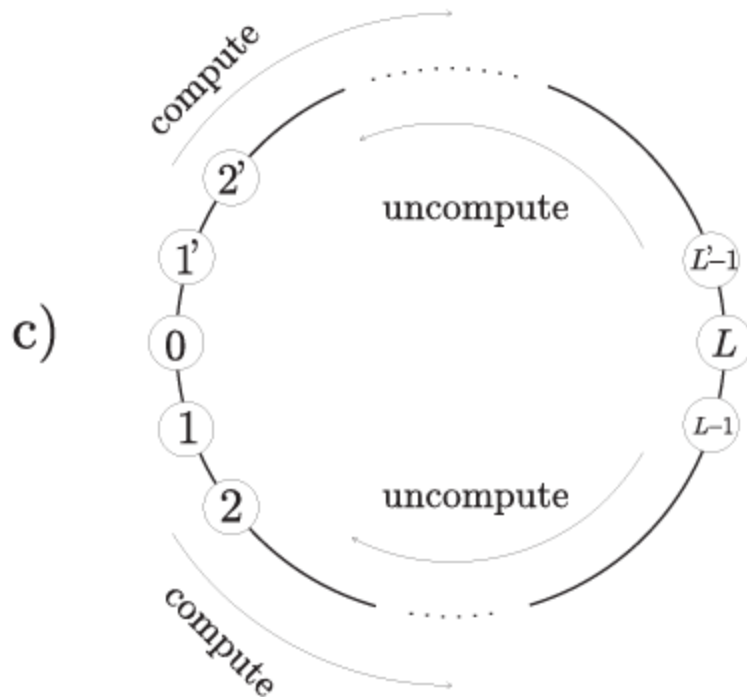
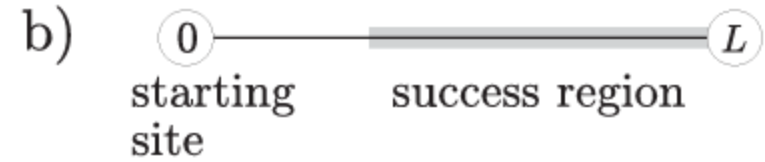
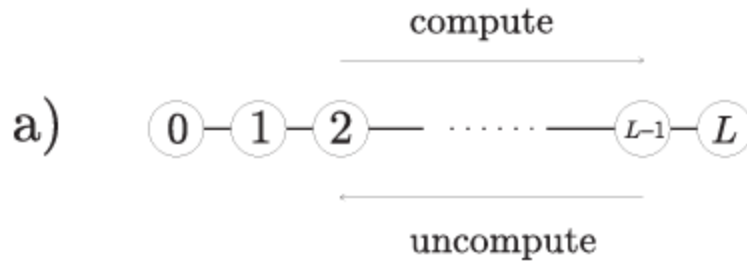
a possibility: wrap around a circle

$$\begin{array}{l} |\varphi_0\rangle \otimes |0\rangle_c \\ U_1 |\varphi_0\rangle \otimes |1\rangle_c \\ U_2 U_1 |\varphi_0\rangle \otimes |2\rangle_c \\ U_3 U_2 U_1 |\varphi_0\rangle \otimes |3\rangle_c \\ U_4 U_3 U_2 U_1 |\varphi_0\rangle \otimes |4\rangle_c \end{array} \quad H_F = - \begin{bmatrix} 0 & 1 & 0 & 0 & \mathbf{1} \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ \mathbf{1} & 0 & 0 & 1 & 0 \end{bmatrix}$$

- The eigenvectors: combinations of plane waves

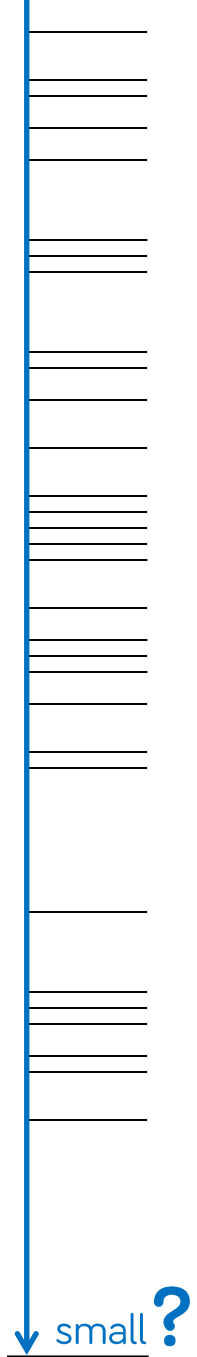
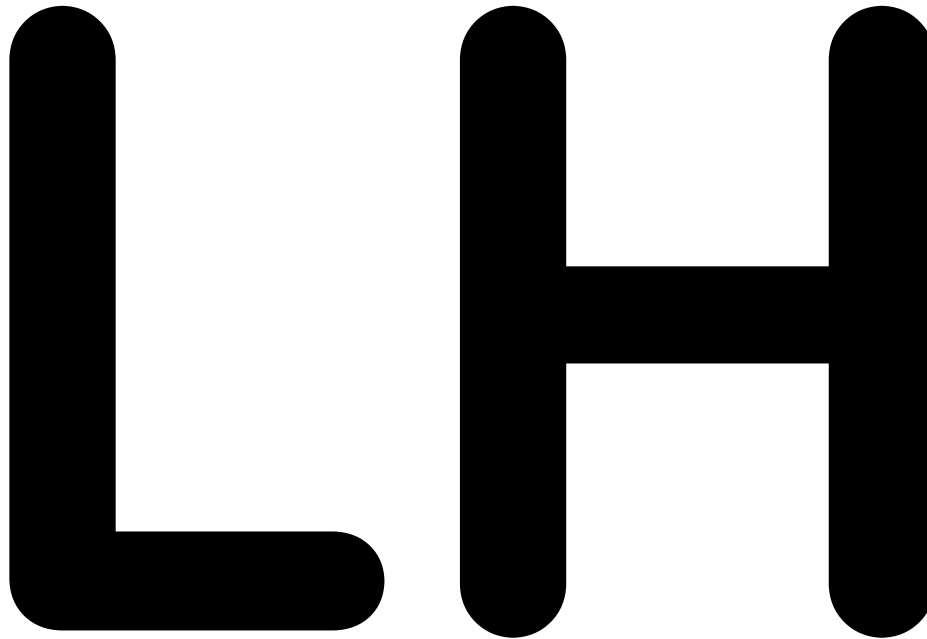
2 Boosting the success probability

Cruising at the end.

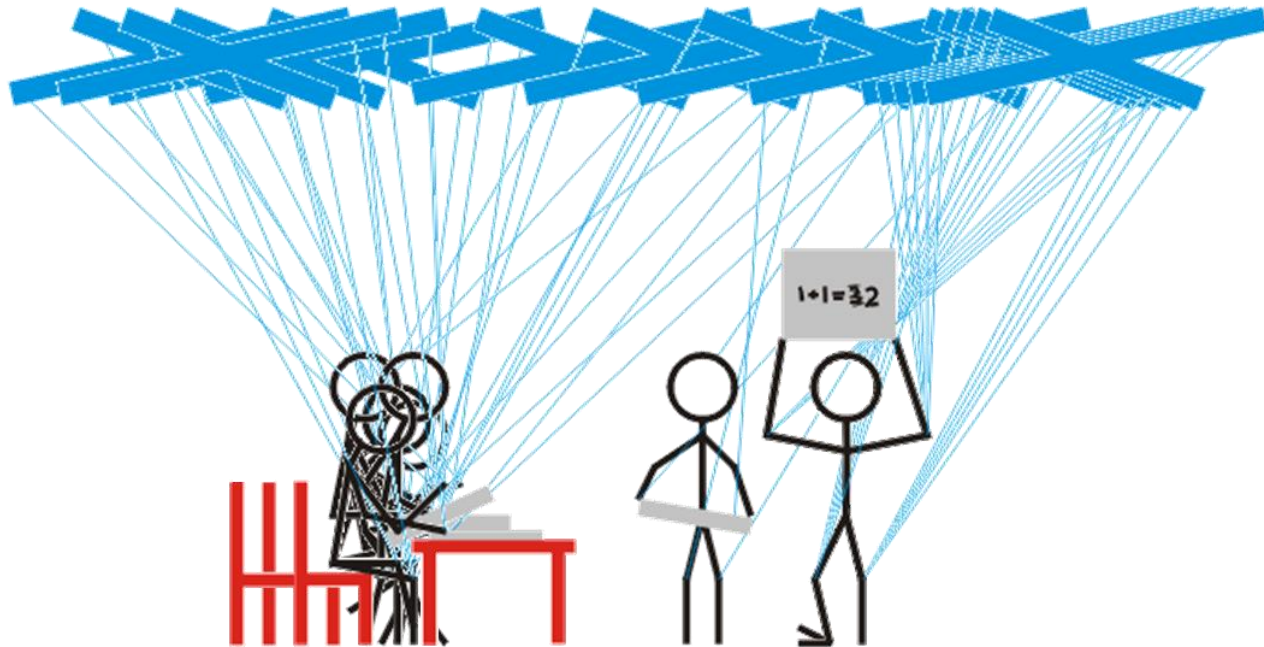


2 The Local Hamiltonian problem

Is
the
ground
state
energy
of a



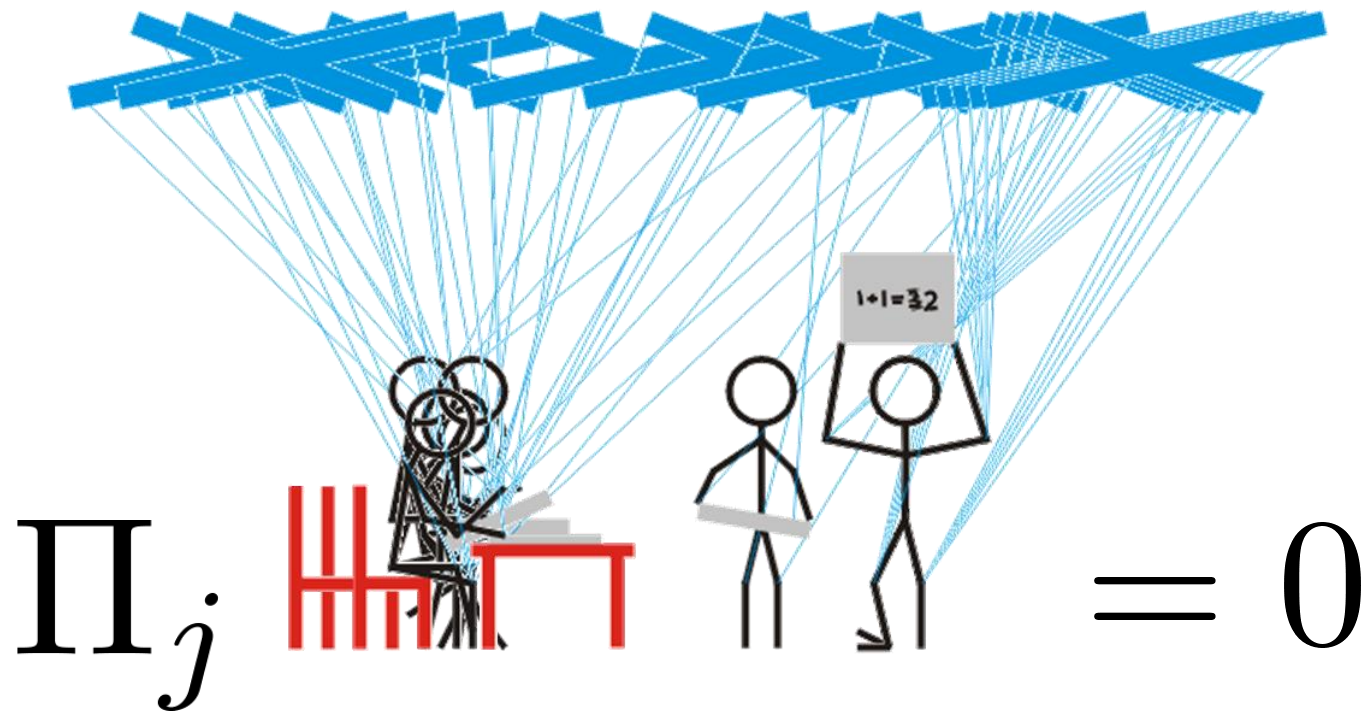
2 The history state



$$|\psi_{hist}\rangle = \frac{1}{\sqrt{T+1}} \sum_{t=0}^T |\varphi_t\rangle \otimes |t\rangle$$

$\underbrace{U_t \cdots U_1 |\varphi_0\rangle}_{|t\rangle}$

2 The history state: a ground state



$$|\psi_{hist}\rangle = \frac{1}{\sqrt{T+1}} \sum_{t=0}^T \underbrace{|\varphi_t\rangle \otimes |t\rangle}_{U_t \cdots U_1 |\varphi_0\rangle}$$

2 Do we have a history state?

k-local
c-o-n-d-i-t-i-o-n-s

clock encoding
state progression

$$\begin{aligned} &|\varphi_t\rangle \otimes |t\rangle \\ &|\varphi_{t+1}\rangle \otimes |t+1\rangle \end{aligned}$$

$$|\psi_{hist}\rangle = \frac{1}{\sqrt{T+1}} \sum_{t=0}^T |\varphi_t\rangle \otimes |t\rangle$$



2 Checking if a state is an accepted history

k-local
c-o-n-d-i-t-i-o-n-s

clock encoding
state progression
initialization

$$|\dots 0\rangle \otimes |0\rangle$$

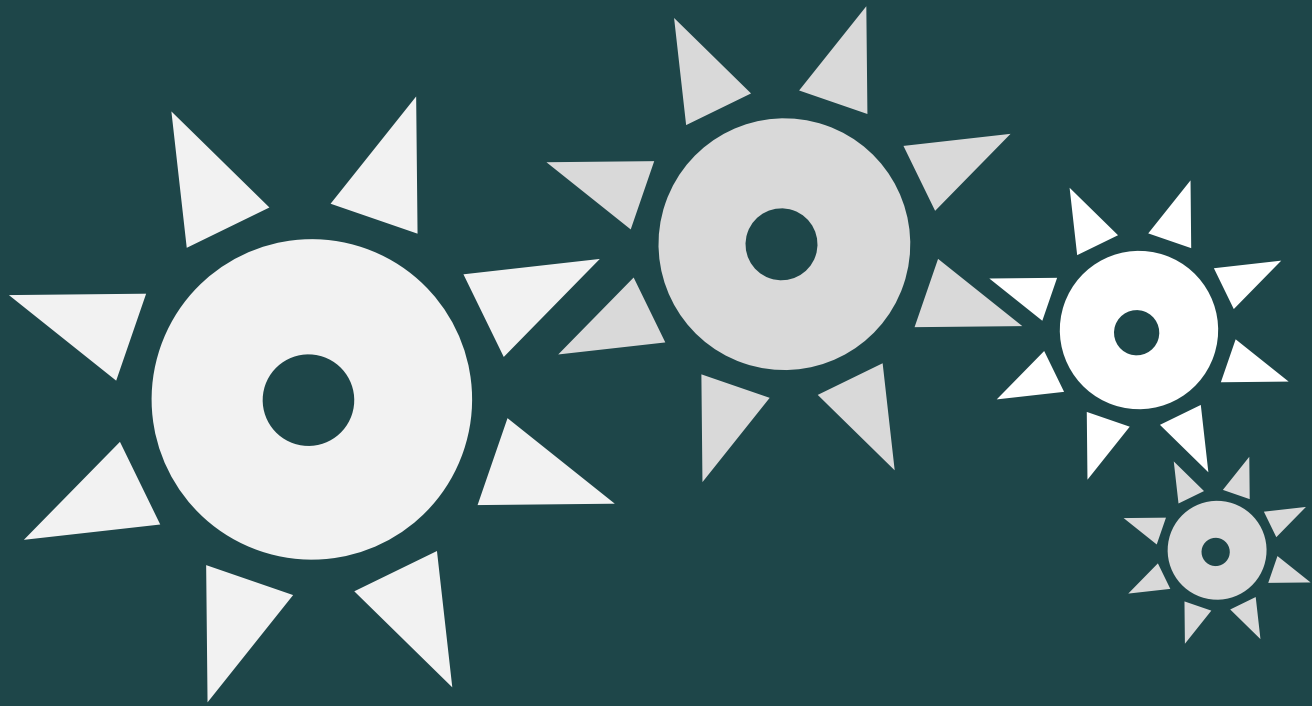
$$|\varphi_t\rangle \otimes |t\rangle$$

$$|\varphi_{t+1}\rangle \otimes |t+1\rangle$$

$$|\psi_{hist}\rangle = \frac{1}{\sqrt{T+1}} \sum_{t=0}^T |\varphi_t\rangle \otimes |t\rangle$$

output $|\dots 1\rangle \otimes |T\rangle$





a clock workshop

3 Run the clock, apply gates ...

$$|\varphi_{t-2}\rangle \otimes |t-2\rangle$$

$$|\varphi_{t-1}\rangle \otimes |t-1\rangle$$

$$|\varphi_t\rangle \otimes |t\rangle$$

$$|\varphi_{t+1}\rangle \otimes |t+1\rangle$$

$$|\varphi_{t+2}\rangle \otimes |t+2\rangle$$

$$|\varphi_{t+3}\rangle \otimes |t+3\rangle$$



3 Constructing local clocks

- the pulse



3 Constructing local clocks

- the pulse



transitions: 2-local

- joining the states
by projectors



3 Constructing local clocks

- the pulse



transitions: 2-local

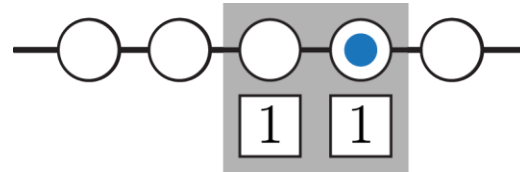
$$\begin{array}{c} 00100 \\ \downarrow \\ +00010 \end{array}$$

- joining the states
by projectors

$$|01 - 10\rangle\langle 01 - 10|$$

3 Constructing local clocks

- the pulse



transitions: 2-local
2-qubit gates: 4-local

interaction
with the data

- joining the states

by projectors $|01 - 10\rangle\langle 01 - 10|$

3 Constructing local clocks

- the pulse



transitions: 2-local
2-qubit gates: 4-local

00000

a “dead” state

Initialization!

- joining the states
by projectors

$$|01 - 10\rangle\langle 01 - 10|$$

3 Constructing local clocks

- the domain wall 

$$\begin{aligned} |t\rangle &= |4\rangle \\ &= |11110\rangle \end{aligned}$$

- 2-local terms
“compatible” with
11...1100...00



$$|01\rangle\langle 01|$$

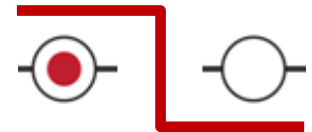
3 Constructing local clocks

- the domain wall  transitions: 3-local

$$\begin{aligned}
 |t\rangle &= |3\rangle \\
 &= |11100\rangle
 \end{aligned}$$

- joining states by transitions? $|100 - 110\rangle\langle 100 - 110|$

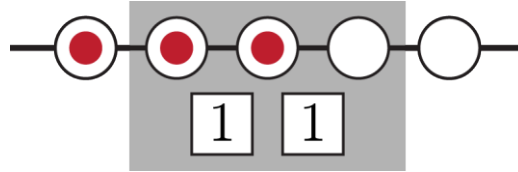
- enforce a domain wall: fix the ends



- the ground state $\cdots + |2\rangle + |3\rangle + \cdots$

3 Constructing local clocks

- the domain wall



transitions: 3-local
2-qubit gates: 5-local

- interacting with
work (data) qubits

$$H_t = \frac{1}{2} (|t+1\rangle\langle t+1| + |t\rangle\langle t|) - \frac{1}{2} \left(U_{t+1} \otimes |t+1\rangle\langle t| + U_{t+1}^\dagger \otimes |t\rangle\langle t+1| \right)$$

5-local

3 Local Hamiltonian: putting it all together

- punish bad ancilla initialization

$$\sum_{a=1}^{n_a} |1\rangle\langle 1|_a \otimes |10\rangle\langle 10|_{c_1, c_2} \quad |\varphi_0\rangle \otimes |0\rangle_c$$

- check the computation

$$H_t = \frac{1}{2} (|t+1\rangle\langle t+1| + |t\rangle\langle t|) \\ - \frac{1}{2} \left(U_{t+1} \otimes |t+1\rangle\langle t| + U_{t+1}^\dagger \otimes |t\rangle\langle t+1| \right)$$

- punish non-accepting computations

$$|0\rangle\langle 0|_{out} \otimes |1\rangle\langle 1|_{c_L} \quad |\varphi_L\rangle \otimes |L\rangle_c$$



● YES

ground state

• NO



lower bound on the
ground state energy

good
clock
states

... 01 ...
bad
clock
states

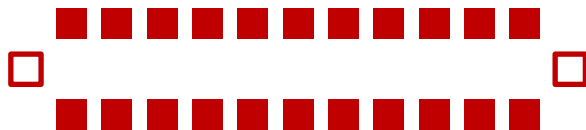
history states

non-uniform
superpositions

history states



a polynomially small gap



$$\Delta = O(L^{-2})$$



history states

well

badly

initialized history states

well

initialized histories

accepted
states

well

initialized histories

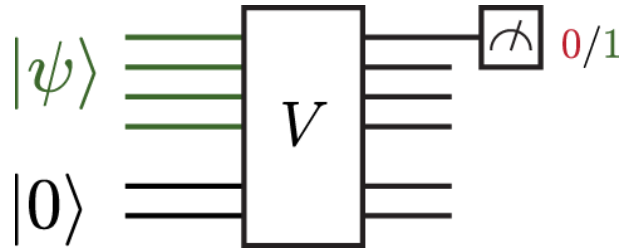
accepted
states

$$H_A + H_B$$

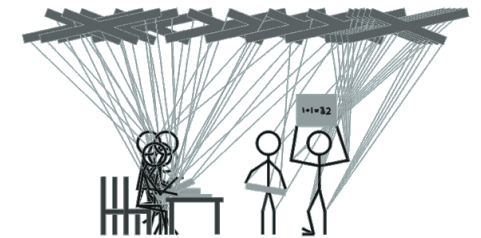
$$\lambda_0 \geq \sin^2 \frac{\vartheta}{2} \times \min(\Delta_A, \Delta_B)$$

\uparrow L^{-2} \uparrow L^{-1}

3 LH and QMA verification



$$H_{clock} + H_{init} + H_{prop} + H_{out}$$



NO no witness accepted by V more likely than ϵ

any state has energy

$$\langle \eta | H | \eta \rangle \geq \frac{c(1 - \sqrt{\epsilon})}{L^3}$$

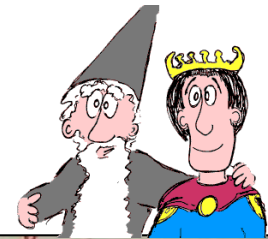
YES there is a proof accepted by V with probability $1 - \epsilon$

$$\langle \psi_{hist} | H | \psi_{hist} \rangle \leq 0 + 0 + \frac{\epsilon}{L + 1}$$

the history for the proof



7 A randomly checked proof

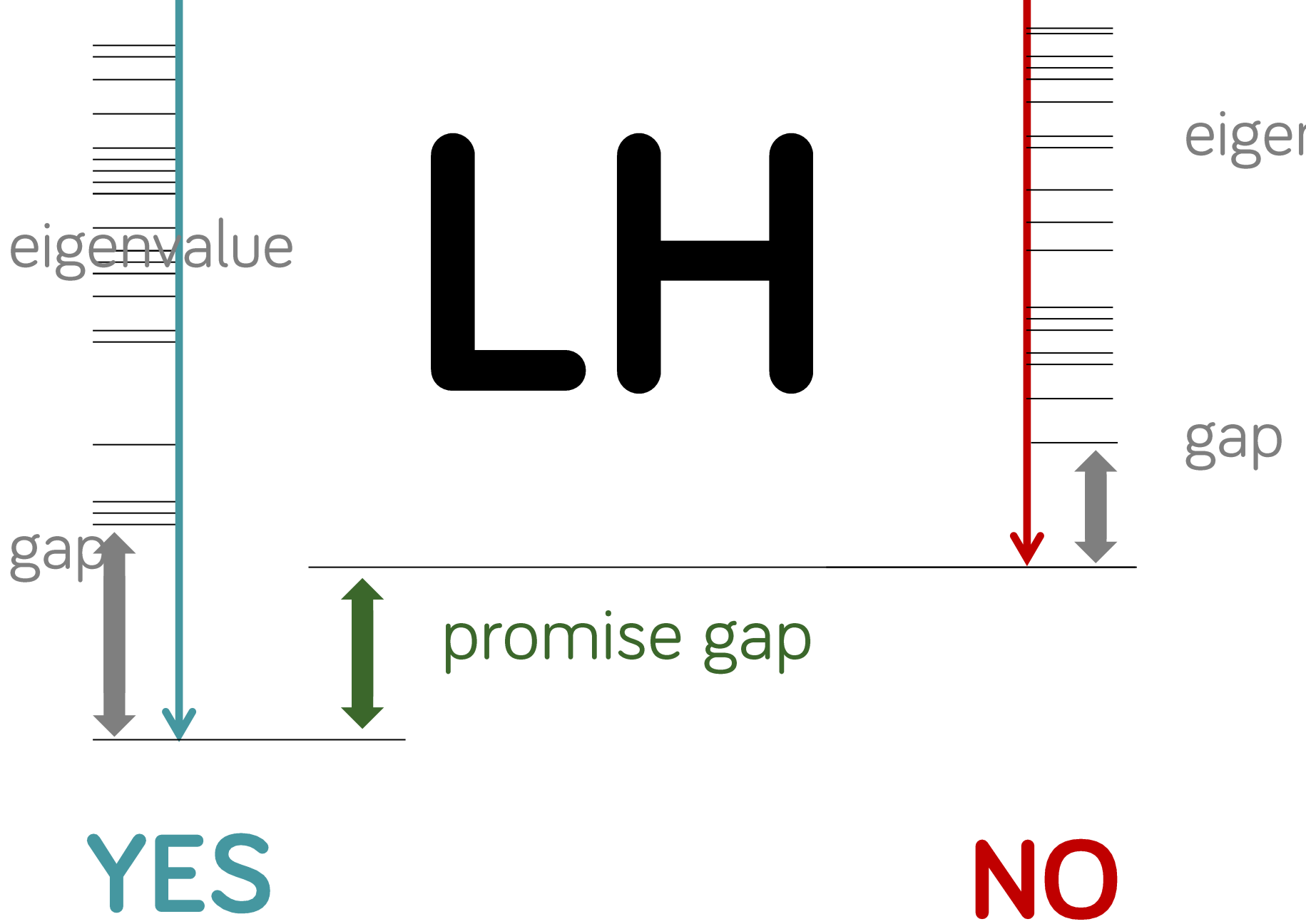


Handwritten notes on a whiteboard, featuring several orange arrows pointing to specific sections:

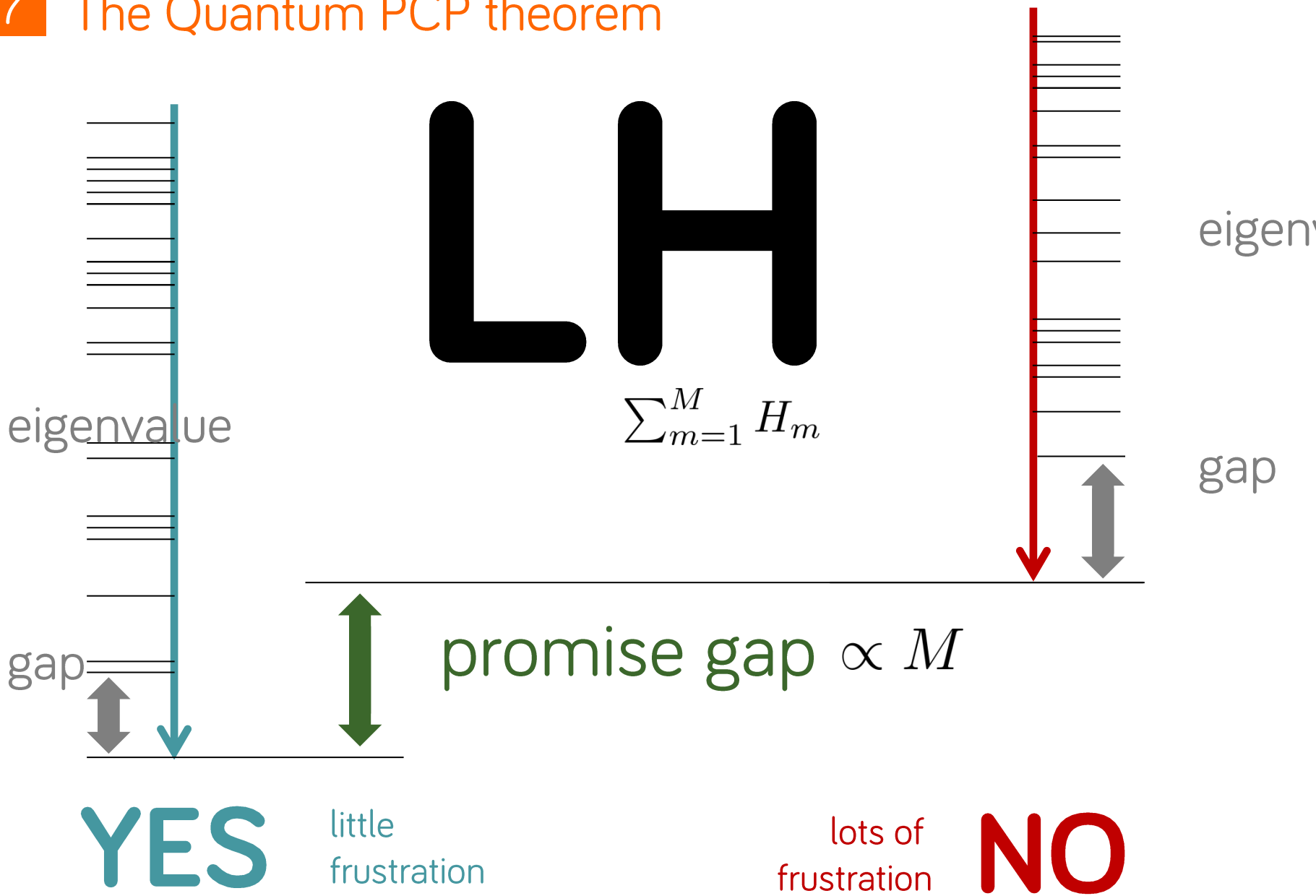
- Top Left:** A sequence of states $|10^{(1)}\rangle, |10^{(2)}\rangle, |10^{(3)}\rangle, |10^{(4)}\rangle, \dots$ with a label "states ruled out".
- Top Center:** A box containing "Qubit - comp bits (linear, linear, Gatterman...)" and "BQP universality (loggy)".
- Top Right:** A grid of bits:

1	1	1	1	1	1	0	0	0	0
1	1	1	1	1	1	1	0	0	0
1	1	1	1	1	1	0	1	1	0

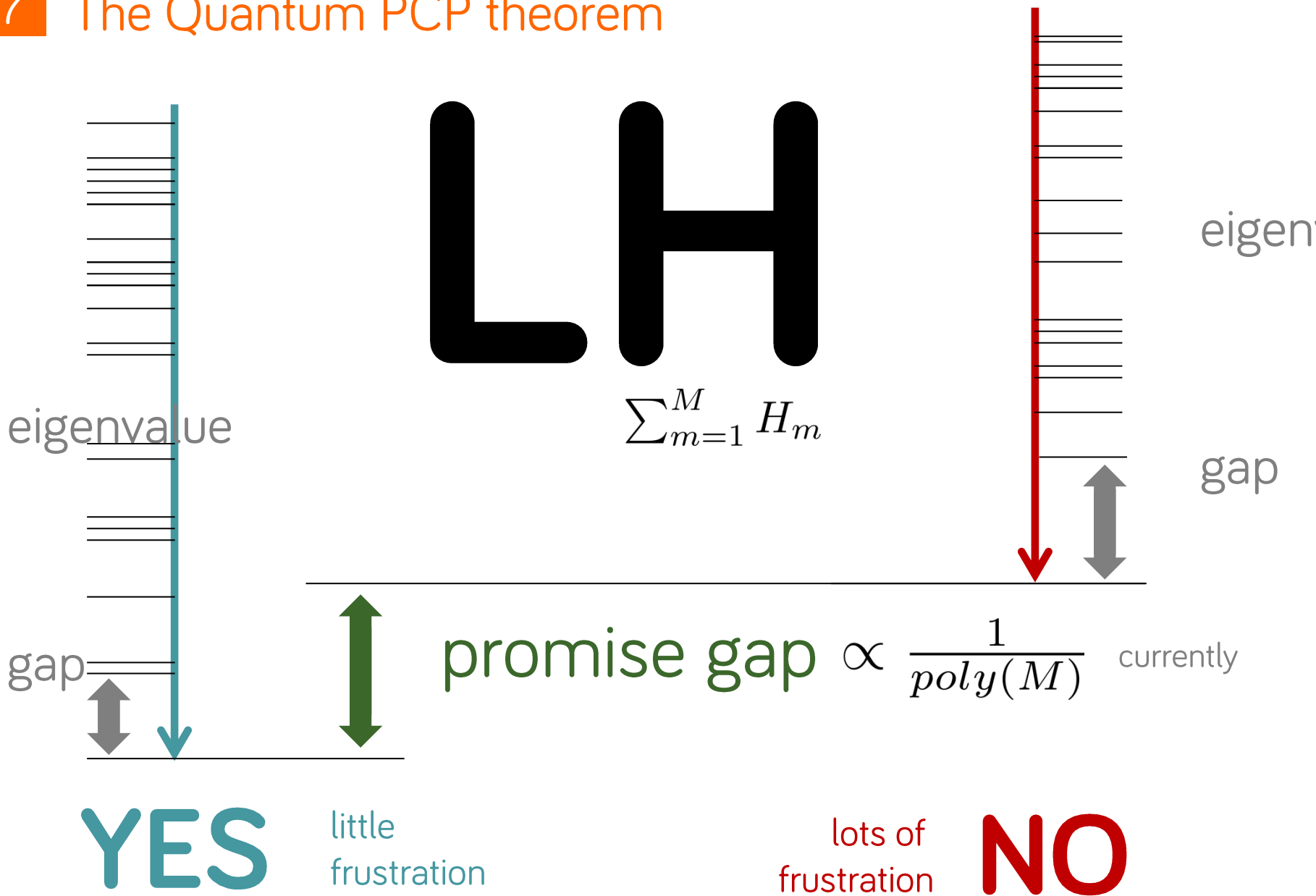
 Below it, a pulse diagram labeled "no pulse" and "domain wall".
- Middle Left:** A box labeled "d=2" with "solutions: $|D_2\rangle = \dots$ ".
- Middle Center:** A box labeled "d=3" with "conjecture: $A(50)$ have a $\square\%$ stab." and "very entangled local ground state".
- Middle Right:** A diagram of a chain with sites P_1, P_2, P_3 and a state $|D_2\rangle = \text{Sign}[\frac{1}{\sqrt{2}} \dots]$.
- Bottom Left:** A diagram of a lattice structure with a state $|D_0\rangle = |00000\rangle$.
- Bottom Center:** A state $|4\rangle = \sum_{k=0}^{n-3} |1111 2222 0000 0\rangle$ and a diagram of a lattice.
- Bottom Right:** A list of states:
 - $|2\rangle|1\rangle$
 - $|1\rangle|1\rangle|1\rangle$
 - $|10\rangle|2\rangle$
 - $|12\rangle|10\rangle$
 - $|11\rangle|11\rangle$
 - $|11\rangle|11\rangle$
 - $|10\rangle|2\rangle$



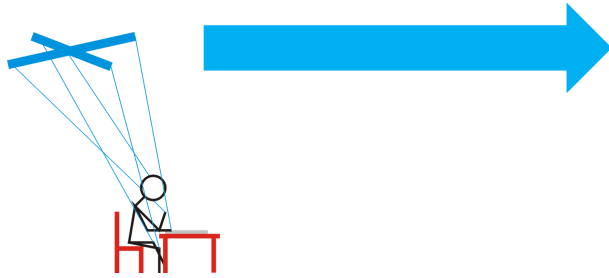
7 The Quantum PCP theorem



7 The Quantum PCP theorem

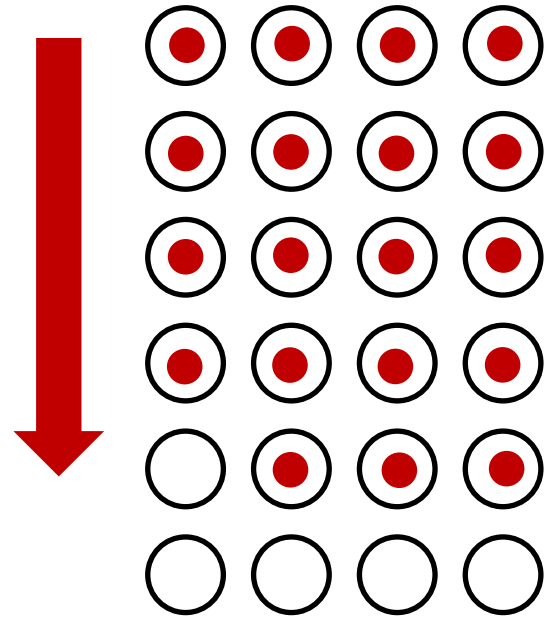


clock/work registers



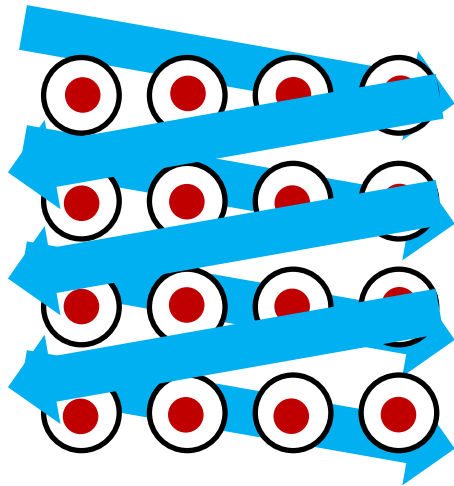
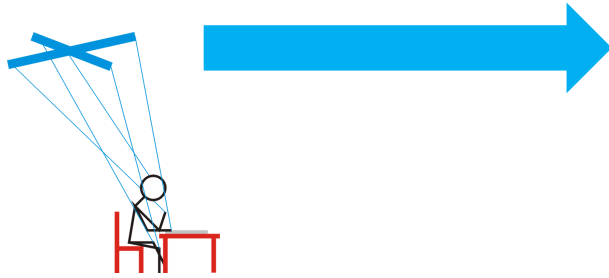
[Kempe, Kitaev, Regev]

a geometric clock



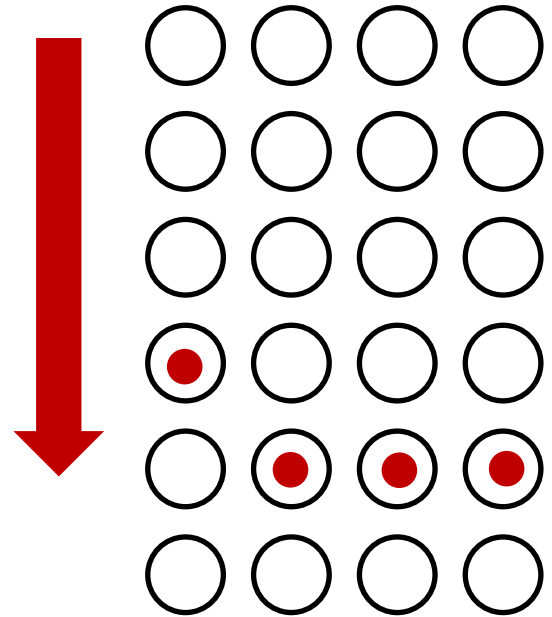
[Mizel] [Aharonov+]

clock/work registers

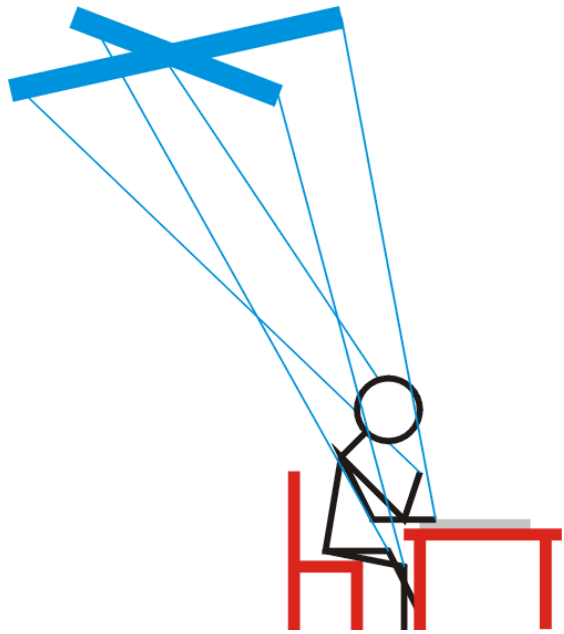


geometric locality

a geometric clock



moving data on a line



An introduction to

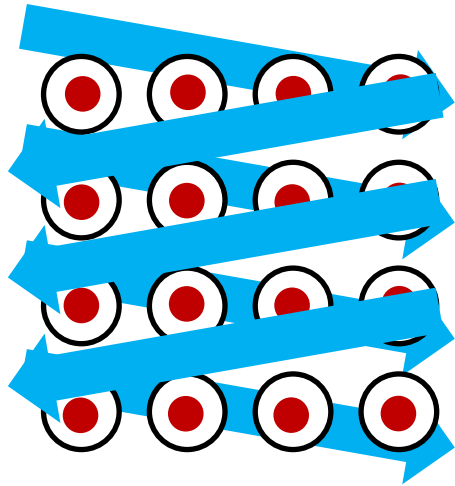
Local Hamiltonians & Quantum Complexity



QMA
complete

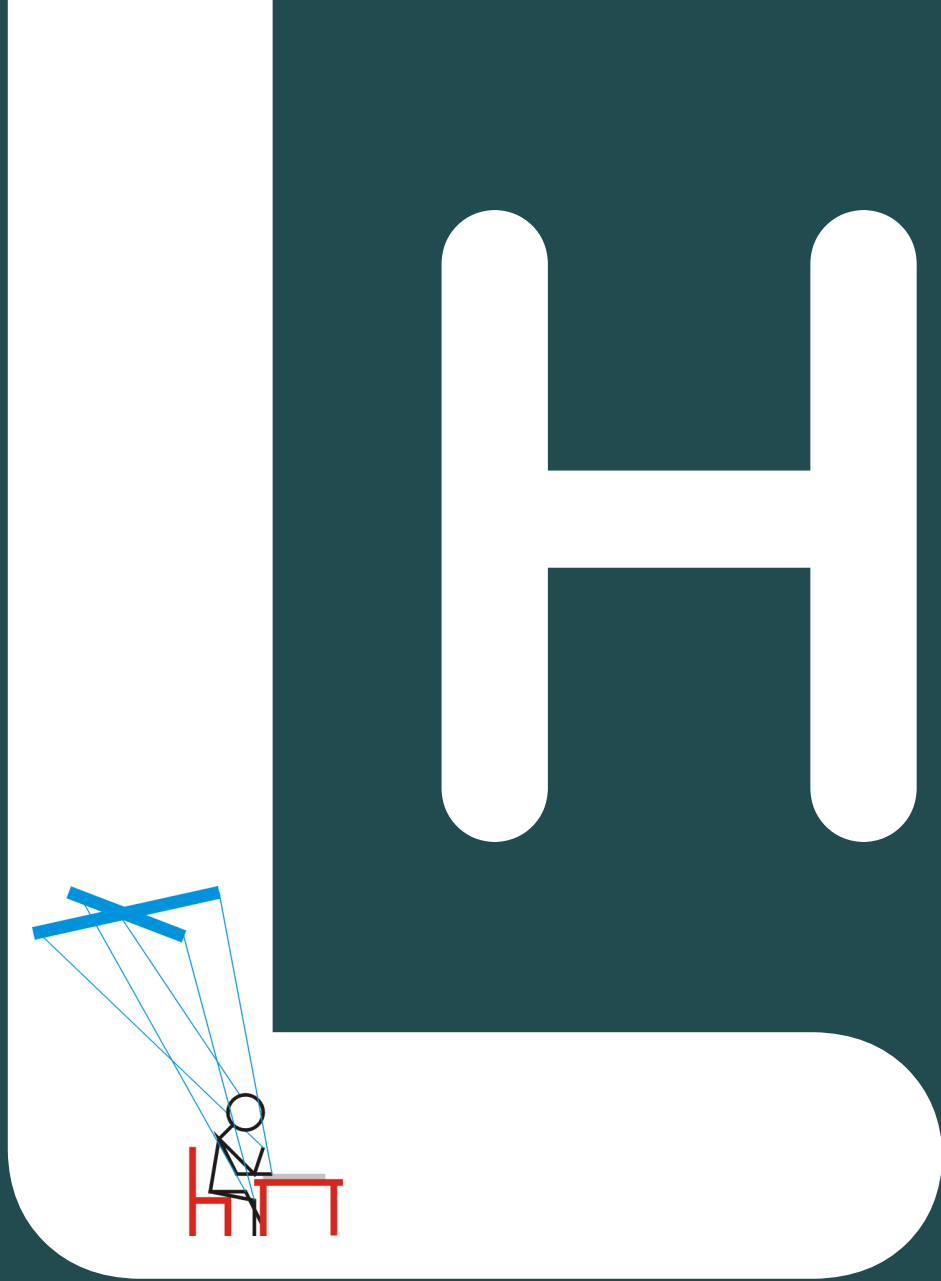


history state



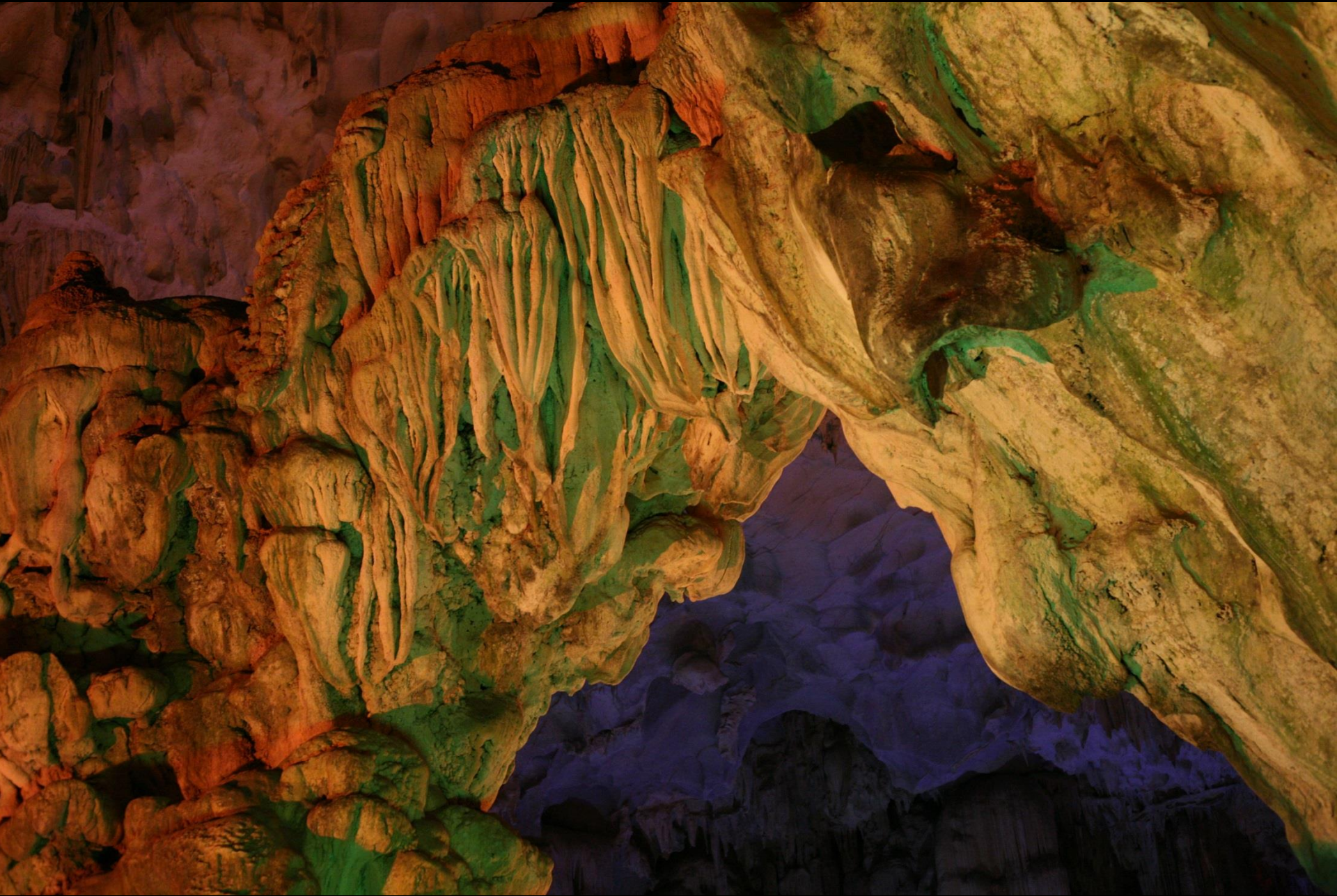
clock

constructions



Quantum SAT





VẬT LÝ QUANG VINH MUÔN NĂM!